6. Planning as search: search algorithms

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# Introduction

#### Introduction

Nodes and states Search for plannir Common procedures

Uninformed search

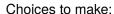
Heuristic search

Summary

# Introduction to search algorithms for planning

# Our plan for the next lectures





- search direction: progression/regression/both → previous chapter

- search control: heuristics, pruning techniques
   → next chapters

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- Search algorithms are used to find solutions (plans) for transition systems in general, not just for planning tasks.
- Planning is one application of search among many.
- In this chapter, we describe some popular and/or representative search algorithms, and (the basics of) how they apply to planning.
- Most of this is review of material that should be known (details: Russell and Norvig's textbook).

#### Introduction

#### Nodes and states

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## Search states vs. search nodes





#### In search, one distinguishes:

- search states s → states (vertices) of the transition system
- search nodes σ → search states plus information on where/when/how they are encountered during search

#### What is in a search node?

Different search algorithms store different information in a search node  $\sigma$ , but typical information includes:

- $\blacksquare$  *state*( $\sigma$ ): associated search state
- $\blacksquare$  *parent*( $\sigma$ ): pointer to search node from which  $\sigma$  is reached
- **action**( $\sigma$ ): action leading from  $state(parent(\sigma))$  to  $state(\sigma)$
- $g(\sigma)$ : cost of  $\sigma$  (length of path from the root node)

For the root node,  $parent(\sigma)$  and  $action(\sigma)$  are undefined.

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# Search states vs. planning states



#### Search states ≠ (planning) states:

- Search states don't have to correspond to states in the planning sense.
  - progression: search states ≈ (planning) states
  - lacktriangledown regression: search states pprox sets of states (formulae)
- Search algorithms for planning where search states are planning states are called state-space search algorithms.
- Strictly speaking, regression is not an example of state-space search, although the term is often used loosely.
- However, we will put the emphasis on progression, which is almost always state-space search.

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# Required ingredients for search



A general search algorithm can be applied to any transition system for which we can define the following three operations:

- init(): generate the initial state
- is-goal(s): test if a given state is a goal state
- succ(s): generate the set of successor states of state s, along with the operators through which they are reached (represented as pairs  $\langle o, s' \rangle$  of operators and states)

Together, these three functions form a search space (a very similar notion to a transition system).

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# Search for planning: progression





Let  $\Pi = \langle A, I, O, \gamma \rangle$  be a planning task.

## Search space for progression search

states: all states of  $\Pi$  (assignments to A)

■ is-goal(s) = 
$$\begin{cases} \text{true} & \text{if } s \models \gamma \\ \text{false} & \text{otherwise} \end{cases}$$

$$\blacksquare$$
  $succ(s) = \{\langle o, s' \rangle \mid applicable \ o \in O, s' = app_o(s)\}$ 

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# Search for planning: regression



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Let  $\Pi = \langle A, I, O, \gamma \rangle$  be a planning task.

## Search space for regression search

states: all formulae over A (how many?)

■ init() = 
$$\gamma$$

■ is-goal(
$$\varphi$$
) = 
$$\begin{cases} \text{true} & \text{if } I \models \varphi \\ \text{false} & \text{otherwise} \end{cases}$$

succ(
$$\varphi$$
) = { $\langle o, \varphi' \rangle \mid o \in O$ ,  
 $\varphi' = regr_o(\varphi), \varphi'$  is satisfiable}  
(modified if splitting is used)

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# Classification of search algorithms



#### uninformed search vs. heuristic search:

- uninformed search algorithms only use the basic ingredients for general search algorithms
- heuristic search algorithms additionally use heuristic functions which estimate how close a node is to the goal

#### systematic search vs. local search:

- systematic algorithms consider a large number of search nodes simultaneously
- local search algorithms work with one (or a few) candidate solutions (search nodes) at a time
- not a black-and-white distinction; there are crossbreeds (e.g., enforced hill-climbing)

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- For satisficing planning, heuristic search vastly outperforms uninformed algorithms on most domains.
- For optimal planning, the difference is less pronounced.

systematic search vs. local search:

- For satisficing planning, the most successful algorithms are somewhere between the two extremes.
- For optimal planning, systematic algorithms are required.

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# Common procedures for search algorithms



Before we describe the different search algorithms, we introduce three procedures used by all of them:

- make-root-node: Create a search node without parent.
- make-node: Create a search node for a state generated as the successor of another state.
- extract-solution: Extract a solution from a search node representing a goal state.

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#### Procedure make-root-node



make-root-node: Create a search node without parent.

#### Procedure make-root-node

```
def make-root-node(s):
       \sigma := \mathbf{new} \text{ node}
      state(\sigma) := s
```

 $parent(\sigma) := undefined$ 

 $action(\sigma) := undefined$ 

 $g(\sigma) := 0$ return  $\sigma$ 

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#### Procedure make-node





make-node: Create a search node for a state generated as the successor of another state.

#### Procedure make-node

```
def make-node(\sigma, o, s):

\sigma' := \mathbf{new} node

state(\sigma') := s

parent(\sigma') := \sigma

action(\sigma') := o

g(\sigma') := g(\sigma) + 1

return \sigma'
```

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#### Procedure extract-solution



extract-solution: Extract a solution from a search node representing a goal state.

#### Procedure extract-solution

```
def extract-solution(\sigma):

solution := new list

while parent(\sigma) is defined:

solution.push-front(action(\sigma))

\sigma := parent(\sigma)

return solution
```

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# Uninformed search algorithms

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#### Uninformed search

duplicate detect Breadth-first wit duplicate detect Random walk

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# Uninformed search algorithms

- UNI
- Uninformed algorithms are less relevant for planning than heuristic ones, so we keep their discussion brief.
- Uninformed algorithms are mostly interesting to us because we can compare and contrast them to related heuristic search algorithms.

Popular uninformed systematic search algorithms:

- breadth-first search
- depth-first search
- iterated depth-first search

Popular uninformed local search algorithms:

random walk

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duplicate detection Breadth-first with duplicate detection

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```
queue := new 	ext{ fifo-queue} queue.push-back(make-root-node(init())) while not queue.empty(): \sigma = queue.pop-front() if is-goal(state(\sigma)): return extract-solution(\sigma) for each <math>\langle o,s \rangle \in succ(state(\sigma)): \sigma' := make-node(\sigma,o,s) queue.push-back(\sigma') return unsolvable
```

- Possible improvement: duplicate detection (see next slide).
- Another possible improvement: test if  $\sigma'$  is a goal node; if so, terminate immediately. (We don't do this because it obscures the similarity to some of the later algorithms.)

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searcn





## Breadth-first search with duplicate detection

```
queue := new fifo-queue
queue.push-back(make-root-node(init()))
closed := \emptyset
while not queue.empty():
      \sigma = queue.pop-front()
     if state(\sigma) \notin closed:
            closed := closed \cup \{state(\sigma)\}\
            if is-goal(state(\sigma)):
                 return extract-solution(\sigma)
           for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                 \sigma' := \mathsf{make-node}(\sigma, o, s)
                 queue.push-back(\sigma')
return unsolvable
```

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## Breadth-first search with duplicate detection

```
queue := new fifo-queue
queue.push-back(make-root-node(init()))
closed := 0
while not queue.empty():
      \sigma = queue.pop-front()
     if state(\sigma) \notin closed:
           closed := closed \cup \{state(\sigma)\}\
           if is-goal(state(\sigma)):
                 return extract-solution(\sigma)
           for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                 \sigma' := \mathsf{make-node}(\sigma, o, s)
                 queue.push-back(\sigma')
return unsolvable
```

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## Random walk





#### Random walk

```
\begin{split} \sigma &:= \mathsf{make}\text{-root-node}(\mathsf{init}()) \\ \textbf{forever} &: \\ &\quad \textbf{if is-goal}(\mathsf{state}(\sigma)) \text{:} \\ &\quad \textbf{return } \mathsf{extract}\text{-solution}(\sigma) \\ &\quad \mathsf{Choose } \mathsf{a} \mathsf{ random } \mathsf{element } \langle o, s \rangle \mathsf{ from } \mathsf{succ}(\mathsf{state}(\sigma)). \\ &\quad \sigma &:= \mathsf{make}\text{-node}(\sigma, o, s) \end{split}
```

- The algorithm usually does not find any solutions, unless almost every sequence of actions is a plan.
- Often, it runs indefinitely without making progress.
- It can also fail by reaching a dead end, a state with no successors. This is a weakness of many local search approaches.

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# Heuristic search algorithms

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# Heuristic search algorithms: systematic



 Heuristic search algorithms are the most common and overall most successful algorithms for classical planning.

#### Popular systematic heuristic search algorithms:

- greedy best-first search
- A\*
- weighted A\*
- IDA\*
- depth-first branch-and-bound search
- ...

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# Heuristic search algorithms: local



 Heuristic search algorithms are the most common and overall most successful algorithms for classical planning.

#### Popular heuristic local search algorithms:

- hill-climbing
- enforced hill-climbing
- beam search
- tabu search
- genetic algorithms
- simulated annealing
- **.** . . .

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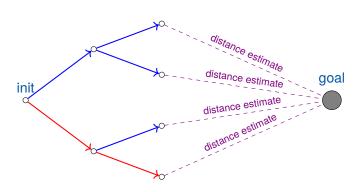
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#### Heuristic search: idea







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# Required ingredients for heuristic search



A heuristic search algorithm requires one more operation in addition to the definition of a search space.

## Definition (heuristic function)

Let  $\Sigma$  be the set of nodes of a given search space. A heuristic function or heuristic (for that search space) is a

function  $h: \Sigma \to \mathbb{N}_0 \cup \{\infty\}$ .

The value  $h(\sigma)$  is called the heuristic estimate or heuristic value of heuristic h for node  $\sigma$ . It is supposed to estimate the distance from  $\sigma$  to the nearest goal node.

Heuristics

# What exactly is a heuristic estimate?



What does it mean that *h* "estimates the goal distance"?

- For most heuristic search algorithms, *h* does not need to have any strong properties for the algorithm to work (= be correct and complete).
- However, the efficiency of the algorithm closely relates to how accurately *h* reflects the actual goal distance.
- For some algorithms, like A\*, we can prove strong formal relationships between properties of *h* and properties of the algorithm (optimality, dominance, run-time for bounded error, ...)
- For other search algorithms, "it works well in practice" is often as good an analysis as one gets.

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# Heuristics applied to nodes or states?



- Most texts apply heuristic functions to states, not nodes.
- This is slightly less general than our definition:
  - Given a state heuristic h, we can define an equivalent node heuristic as  $h'(\sigma) := h(state(\sigma))$ .
  - The opposite is not possible. (Why not?)
- There is good justification for only allowing state-defined heuristics: why should the estimated distance to the goal depend on how we ended up in a given state s?
- We call heuristics which don't just depend on  $state(\sigma)$  pseudo-heuristics.
- In practice there are sometimes good reasons to have the heuristic value depend on the generating path of σ (e.g., landmark pseudo-heuristic, Richter et al. 2008).

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#### Perfect heuristic



) **E** 

Let  $\Sigma$  be the set of nodes of a given search space.

## Definition (optimal/perfect heuristic)

The optimal or perfect heuristic of a search space is the heuristic  $h^*$  which maps each search node  $\sigma$  to the length of a shortest path from  $state(\sigma)$  to any goal state.

Note:  $h^*(\sigma) = \infty$  iff no goal state is reachable from  $\sigma$ .

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# Properties of heuristics



#### A heuristic *h* is called

- safe if  $h^*(\sigma) = \infty$  for all  $\sigma \in \Sigma$  with  $h(\sigma) = \infty$
- **goal-aware** if  $h(\sigma) = 0$  for all goal nodes  $\sigma \in \Sigma$
- **admissible** if  $h(\sigma) \le h^*(\sigma)$  for all nodes  $\sigma \in \Sigma$
- consistent if  $h(\sigma) \le h(\sigma') + 1$  for all nodes  $\sigma, \sigma' \in \Sigma$  such that  $\sigma'$  is a successor of  $\sigma$ .<sup>1</sup>

## Relationships?

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<sup>&</sup>lt;sup>1</sup> or:  $h(\sigma) \le h(\sigma') + cost(\sigma, \sigma')$  for non-unit costs, where  $cost(\sigma, \sigma')$  is the cost of the transition from  $\sigma$  to  $\sigma'$ .

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Summary

# Greedy best-first search (with duplicate detection)

```
open := new min-heap ordered by (\sigma \mapsto h(\sigma))
open.insert(make-root-node(init()))
closed := 0
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed:
            closed := closed \cup \{state(\sigma)\}\
            if is-goal(state(\sigma)):
                  return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                  \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < \infty:
                         open.insert(\sigma')
```

return unsolvable

# Properties of greedy best-first search





- one of the three most commonly used algorithms for satisficing planning
- complete for safe heuristics (due to duplicate detection)
- suboptimal unless h satisfies some very strong assumptions (similar to being perfect)
- invariant under all strictly monotonic transformations of h
   (e.g., scaling with a positive constant or adding a
   constant)

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## A\* (with duplicate detection and reopening)

```
open := new min-heap ordered by (\sigma \mapsto g(\sigma) + h(\sigma))
open.insert(make-root-node(init()))
closed := \emptyset
distance := \emptyset
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed or q(\sigma) < distance(state(\sigma)):
            closed := closed \cup \{state(\sigma)\}
            distance(state(\sigma)) := q(\sigma)
            if is-goal(state(\sigma)):
                  return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                  \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < ∞: open.insert(\sigma')
return unsolvable
```

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Example



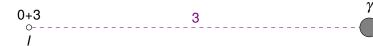


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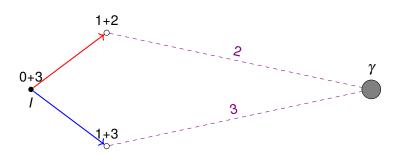
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Example







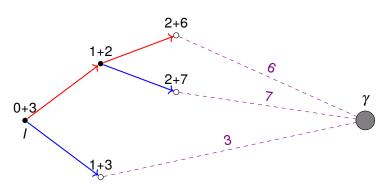
search

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Example





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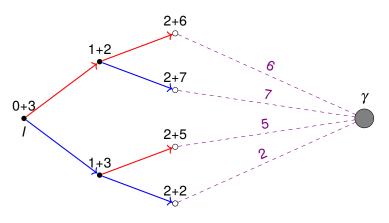
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Example







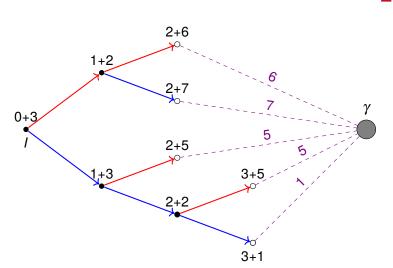
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Example





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# Terminology for A\*



- f value of a node: defined by  $f(\sigma) := g(\sigma) + h(\sigma)$
- generated nodes: nodes inserted into open at some point
- expanded nodes: nodes σ popped from *open* for which the test against *closed* and *distance* succeeds
- reexpanded nodes: expanded nodes for which  $state(\sigma) \in closed$  upon expansion (also called reopened nodes)

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# Properties of A\*

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- the most commonly used algorithm for optimal planning
- rarely used for satisficing planning
- complete for safe heuristics (even without duplicate detection)
- optimal if h is admissible (even without duplicate detection)
- never reopens nodes if h is consistent

#### Implementation notes:

- in the heap-ordering procedure, it is considered a good idea to break ties in favour of lower h values
- can simplify algorithm if we know that we only have to deal with consistent heuristics
- common, hard to spot bug: test membership in closed at the wrong time

## Weighted A\* (with duplicate detection and reopening)

```
open := new min-heap ordered by (\sigma \mapsto g(\sigma) + W \cdot h(\sigma))
open.insert(make-root-node(init()))
closed := \emptyset
distance := \emptyset
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed or g(\sigma) < distance(state(\sigma)):
            closed := closed \cup \{state(\sigma)\}
            distance(\sigma) := g(\sigma)
            if is-goal(state(\sigma)):
                   return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                   \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < ∞: open.insert(\sigma')
return unsolvable
```

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# Properties of weighted A\*



The weight  $W \in \mathbb{R}_0^+$  is a parameter of the algorithm.

- $\blacksquare$  for W = 0, behaves like breadth-first search
- $\blacksquare$  for W = 1, behaves like  $A^*$
- lacksquare for  $W o \infty$ , behaves like greedy best-first search

#### Properties:

- one of the most commonly used algorithms for satisficing planning
- for W > 1, can prove similar properties to A\*, replacing optimal with bounded suboptimal: generated solutions are at most a factor W as long as optimal ones

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# Hill-climbing





## Hill-climbing

```
\sigma := make-root-node(init())

forever:

if is-goal(state(\sigma)):

return extract-solution(\sigma)
```

```
\Sigma' := \{ \text{ make-node}(\sigma, o, s) \mid \langle o, s \rangle \in \text{succ}(\text{state}(\sigma)) \}
```

 $\sigma$  := an element of  $\Sigma'$  minimizing h (random tie breaking)

- can easily get stuck in local minima where immediate improvements of  $h(\sigma)$  are not possible
- many variations: tie-breaking strategies, restarts

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# Enforced hill-climbing: procedure improve

```
def improve(\sigma_0):
      queue := new fifo-queue
      queue.push-back(\sigma_0)
      closed := \emptyset
      while not queue.empty():
            \sigma = queue.pop-front()
            if state(\sigma) \notin closed:
                   closed := closed \cup \{state(\sigma)\}\
                   if h(\sigma) < h(\sigma_0):
                         return \sigma
                  for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                         \sigma' := \mathsf{make-node}(\sigma, o, s)
                         queue.push-back(\sigma')
      fail
```

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Summary

 $\rightsquigarrow$  breadth-first search for more promising node than  $\sigma_0$ 

# Enforced hill-climbing (ctd.)





## Enforced hill-climbing

```
\sigma := make-root-node(init())

while not is-goal(state(\sigma)):

\sigma := improve(\sigma)

return extract-solution(\sigma)
```

- one of the three most commonly used algorithms for satisficing planning
- $\blacksquare$  can fail if procedure improve fails (when the goal is unreachable from  $\sigma_0$ )
- complete for undirected search spaces (where the successor relation is symmetric) if  $h(\sigma) = 0$  for all goal nodes and only for goal nodes

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Summoru





- distinguish: planning states, search states, search nodes
  - planning state: situation in the world modelled by the task
  - search state: subproblem remaining to be solved
    - In state-space search (usually progression search), planning states and search states are identical.
      - In regression search, search states usually describe sets of states ("subgoals").
  - search node: search state + info on "how we got there"
- search algorithms mainly differ in order of node expansion
  - uninformed vs. informed (heuristic) search
  - local vs. systematic search

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# Summary (ctd.)



- heuristics: estimators for "distance to goal node"
  - usually: the more accurate, the better performance
  - desiderata: safe, goal-aware, admissible, consistent
  - the ideal: perfect heuristic h\*
- most common algorithms for satisficing planning:
  - greedy best-first search
  - weighted A\*
  - enforced hill-climbing
- most common algorithm for optimal planning:
  - A\*

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