

Motivation	BURG	
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	Motivation	
	Effect normal form	
Similarly to normal forms in propositional logic (DNF, CNF	normal form	
NNF,) we can define normal forms for effects, operato	rs _{strips}	
and planning tasks.	Summary	
This is useful because algorithms (and proofs) then only need		
to deal with effects (resp. operators or tasks) in normal for	rm.	
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Equivalence of operators and effects	22	
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Definition (equivalent effects)	Motivation	
Two effects e and e' over state variables A are equivalent	Effect normal form	
written $e \equiv e'$, if for all states <i>s</i> over <i>A</i> , $[e]_s = [e']_s$.	Equivalences	
	Example	

Definition (equivalent operators)

Two operators *o* and *o'* over state variables *A* are equivalent, written $o \equiv o'$, if they are applicable in the same states, and for all states *s* where they are applicable, $app_o(s) = app_{o'}(s)$.

Theorem

Let $o = \langle \chi, e \rangle$ and $o' = \langle \chi', e' \rangle$ be operators with $\chi \equiv \chi'$ and $e \equiv e'$. Then $o \equiv o'$.

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Note: The converse is not true. (Why not?)

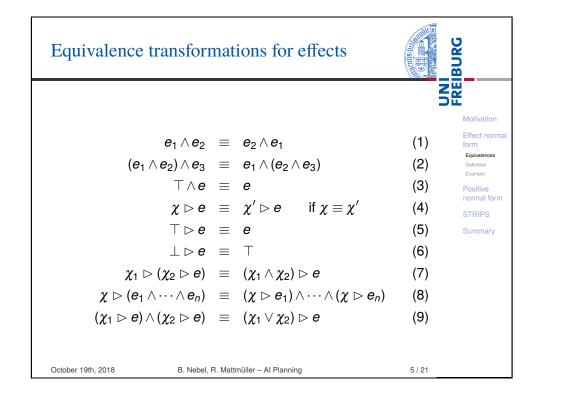
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Positive

STRIPS

Summary

normal form





Definition

An operator $\langle \chi, e \rangle$ is in effect normal form (ENF) if for all occurrences of $\chi' \triangleright e'$ in *e* the effect *e'* is either *a* or $\neg a$ for some $a \in A$, and there is at most one occurrence of any atomic effect in *e*.

normal form STRIPS

Motivation

Effect norma

form Equivalence

Definition

Example

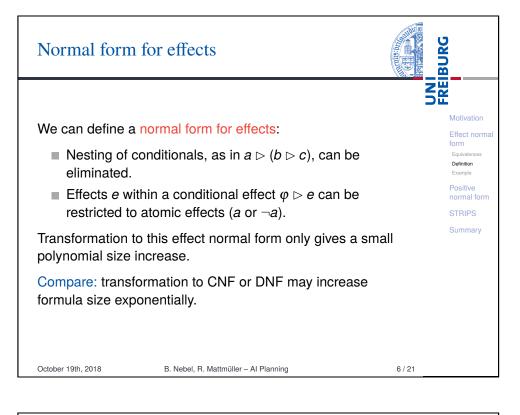
Positive

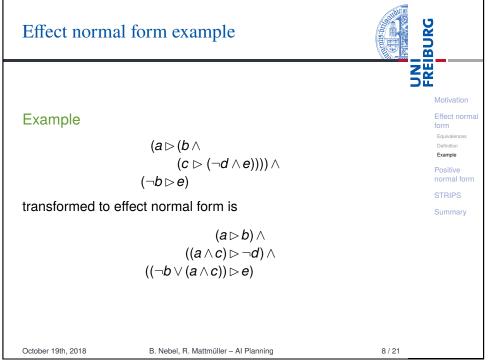
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Theorem

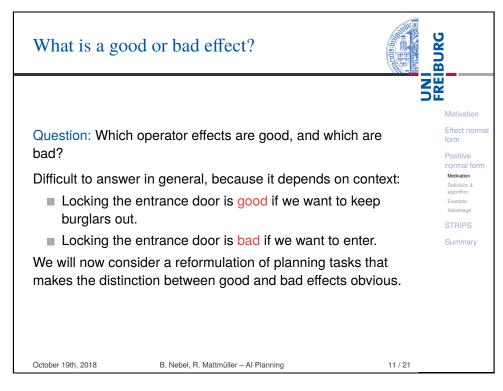
For every operator there is an equivalent one in effect normal form.

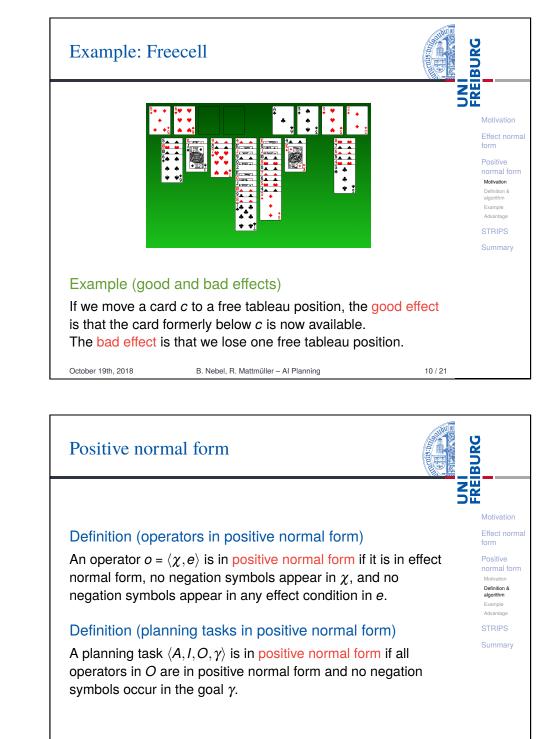
Proof is constructive: we can transform any operator into effect normal form using the equivalence transformations for effects.











Positive normal form: existence



Motivation

Effect norr

Positive

normal forn Motivation

Definition & algorithm Example

Advantage

Summary

Theorem (positive normal form)

Every planning task Π has an equivalent planning task Π' in positive normal form.

Moreover, Π' can be computed from Π in polynomial time.

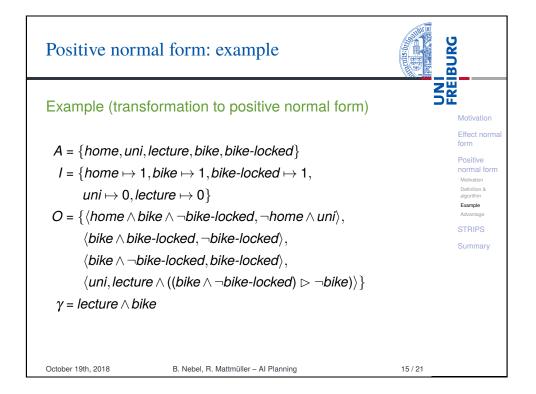
Note: Equivalence here means that the represented transition systems of Π and Π' , limited to the states that can be reached from the initial state, are isomorphic.

We prove the theorem by describing a suitable algorithm. (However, we do not prove its correctness or complexity.)

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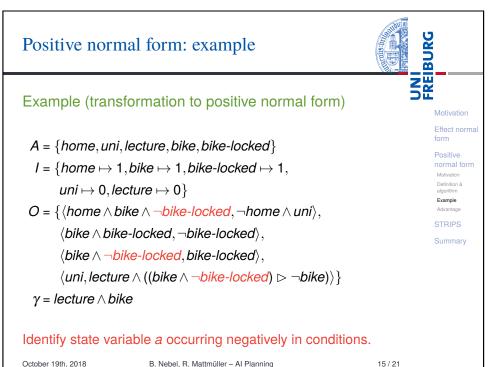
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Positive normal form: algorithm	BURG
Transformation of $\langle A, I, O, \gamma \rangle$ to positive normal form Convert all operators $o \in O$ to effect normal form. Convert all conditions to negation normal form (NNF). while any condition contains a negative literal $\neg a$: Let <i>a</i> be a variable which occurs negatively in a condition. $A := A \cup \{\hat{a}\}$ for some new state variable \hat{a} $I(\hat{a}) := 1 - I(a)$ Replace the effect <i>a</i> by $(a \land \neg \hat{a})$ in all operators $o \in O$. Replace the effect $\neg a$ by $(\neg a \land \hat{a})$ in all operators $o \in O$. Replace $\neg a$ by \hat{a} in all conditions. Convert all operators $o \in O$ to effect normal form (again). Here, <i>all conditions</i> refers to all operator preconditions, operator effect conditions and the goal.	Motivation Effect norr form Positive normal forn Motivation Definition & algorithm Example Advantage STRIPS Summary

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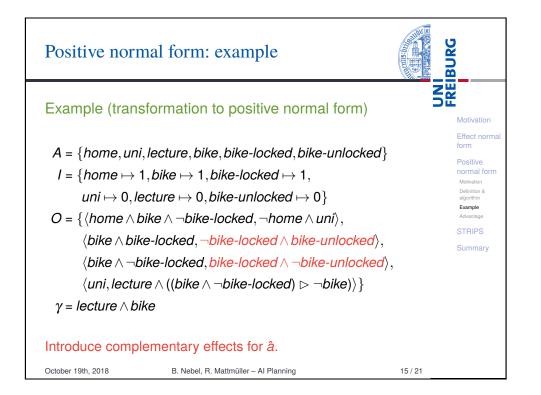
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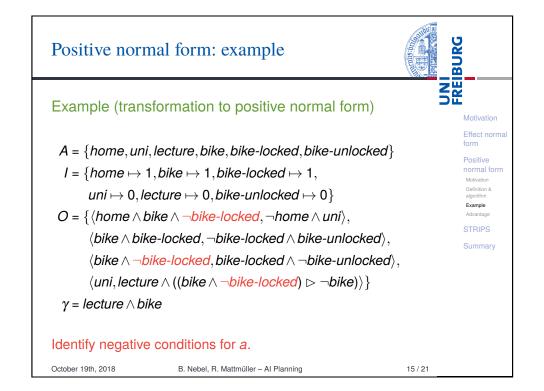
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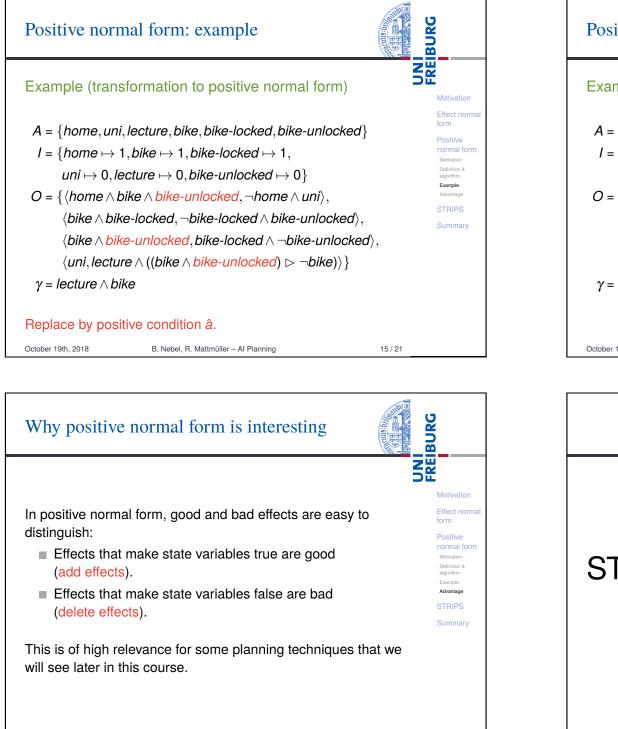
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Positive normal	form: example	BURG	
Example (transforr	nation to positive normal form	1) Motivation	
$I = \{home \mapsto 1, bik$ $uni \mapsto 0, lectur$ $O = \{\langle home \land bike \land \\ \langle bike \land bike-loc \land \\ \langle bike \land \neg bike-loc \land \\ \langle uni, lecture \land \\ \rangle \gamma = lecture \land bike$	ture, bike, bike-locked, bike-unlocked, bike-unlocked \mapsto 1, bike-locked \mapsto 1, $e \mapsto 0$, bike-unlocked $\mapsto 0$ } $\land \neg$ bike-locked, \neg home \land uni \rangle , cked, \neg bike-locked \rangle , ((bike $\land \neg$ bike-locked) $\triangleright \neg$ bike) \rangle	Positive normal form Definition & algorithm Example Advantage STRIPS Summary }	
Introduce new variable \hat{a} with complementary initial value.			
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UNI FREIBURG Positive normal form: example Example (transformation to positive normal form) Motivation A = {home, uni, lecture, bike, bike-locked, bike-unlocked} $I = \{home \mapsto 1, bike \mapsto 1, bike - locked \mapsto 1, \}$ normal for Motivation Definition & $uni \mapsto 0$, lecture $\mapsto 0$, bike-unlocked $\mapsto 0$ } algorithm Example $O = \{ \langle home \land bike \land \neg bike \text{-locked}, \neg home \land uni \rangle, \}$ Advantage $\langle bike \land bike - locked, \neg bike - locked \rangle$, Summary $\langle bike \land \neg bike - locked, bike - locked \rangle$, $\langle uni, lecture \land ((bike \land \neg bike-locked) \triangleright \neg bike) \rangle \}$ γ = lecture \wedge bike Identify effects on variable a. October 19th, 2018 B. Nebel, R. Mattmüller - Al Planning 15/21





UNI FREIBURG Positive normal form: example Example (transformation to positive normal form) Motivation A = {home, uni, lecture, bike, bike-locked, bike-unlocked} $I = \{home \mapsto 1, bike \mapsto 1, bike - locked \mapsto 1, \}$ normal for Motivation Definition & $uni \mapsto 0$, lecture $\mapsto 0$, bike-unlocked $\mapsto 0$ } algorithm Example $O = \{ \langle home \land bike \land bike-unlocked, \neg home \land uni \rangle, \}$ Advantage $\langle bike \land bike-locked, \neg bike-locked \land bike-unlocked \rangle$, Summary $\langle bike \land bike-unlocked, bike-locked \land \neg bike-unlocked \rangle$, $\langle uni, lecture \land ((bike \land bike-unlocked) \triangleright \neg bike) \rangle \}$ γ = lecture \wedge bike 15/21 October 19th, 2018 B. Nebel, R. Mattmüller - Al Planning



STRIPS operators

Definition

An operator $\langle \chi, e \rangle$ is a STRIPS operator if

- χ is a conjunction of atoms, and
- e is a conjunction of atomic effects.

Hence every STRIPS operator is of the form

 $\langle a_1 \wedge \cdots \wedge a_n, I_1 \wedge \cdots \wedge I_m \rangle$

where a_i are atoms and l_j are atomic effects.

Note: Sometimes we allow conjunctions of literals as preconditions. We denote this as STRIPS with negative preconditions.

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Transformation to STRIPS

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Effect norr

Positive normal form

Properties

Summarv

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Motivation

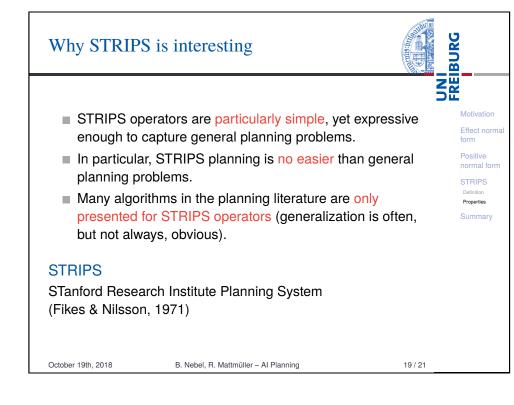
Positive

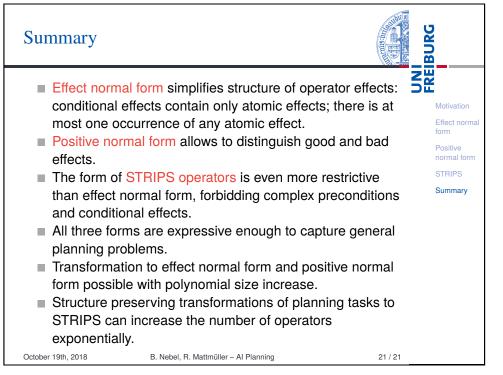
Definition

normal forn

Not every operator is equivalent to a STRIPS operator.

- However, each operator can be transformed into a set of STRIPS operators whose "combination" is equivalent to the original operator. (How?)
- However, this transformation may exponentially increase the number of required operators. There are planning tasks for which such a blow-up is unavoidable.
- There are polynomial transformations of planning tasks to STRIPS, but these do not preserve the structure of the transition system (e.g., length of shortest plans may change).





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