

Principles of AI Planning

16. Strong cyclic planning

Albert-Ludwigs-Universität Freiburg



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January 29th, 2014

1 Strong cyclic plans



- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm

Strong cyclic plans

Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm

Maintenance

Summary

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3 / 70

Planning objectives

Strong plans



Strong cyclic plans

Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm

Maintenance

Summary

- The simplest objective for nondeterministic planning is the one we have considered in the previous lecture: reach a goal state with certainty.
- With this objective the nondeterminism can also be understood as **an opponent** like in 2-player games. The plan guarantees reaching a goal state no matter what the opponent does: plans are **winning strategies**.

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4 / 70

Planning objectives

Limitations of strong plans



Strong cyclic plans

Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm

Maintenance

Summary

- In strong plans, goal states can be reached without visiting any state twice.
- This property guarantees that the length of executions is bounded by some constant (which is smaller than the number of states.)
- Some solvable problems are not solvable this way.
 - 1 Action may fail to have any effect.
Hit a coconut to break it.
 - 2 Action may fail and take us away from the goals.
Build a house of cards.

Consequences:

- 1 It is impossible to avoid visiting some states several times.
- 2 There is no finite upper bound on execution length.

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5 / 70

Planning objectives

When strong cyclic plans make sense



Fairness assumption

For any nondeterministic operator $\langle \chi, \{e_1, \dots, e_n\} \rangle$, the “probability” of every effect $e_i, i = 1, \dots, n$, is greater than 0.

Alternatively: For each $s' \in \text{img}_o(s)$ the “probability” of reaching s' from s by o is greater than 0.

This assumption guarantees that a strong cyclic plan reaches the goal **almost certainly** (with probability 1).

This is **not compatible** with viewing nondeterminism as an opponent in a 2-player game: the opponent’s strategy might rule out some of the choices e_1, \dots, e_n .

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

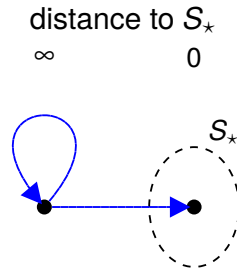
Need for strong cyclic plans

Example



Example (Breaking a coconut)

- Initial state: coconut is intact.
- Goal state: coconut is broken.
- On every hit the coconut may or may not break.
- There is no finite upper bound on the number of hits.



This is equivalent to coin tossing.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

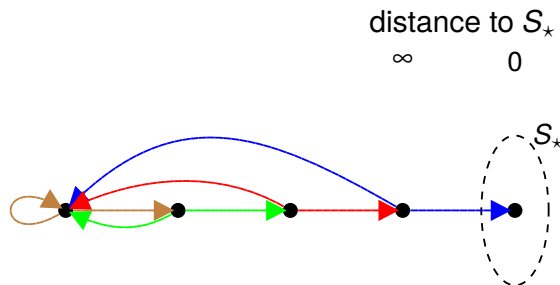
Need for strong cyclic plans

Example



Example (Build a house of cards)

- Initial state: all cards lie on the table.
- Goal state: house of cards is complete.
- At every construction step the house may collapse.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Algorithms for strong cyclic planning

We present two algorithms for strong cyclic planning:

- The **nested fixpoint algorithm** is conceptually simpler, but typically very costly, especially if not implemented symbolically.
 - Historically older
 - **Uninformed**
 - **Considers entire state space**
- The **determinization-based incremental planning algorithm** is a bit more complicated, but typically more efficient.
 - Historically newer, **state of the art**
 - Can use **informed** classical planner as sub-procedure
 - Often only **considers small portion of state space**

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Nested Fixpoint Algorithm

Idea

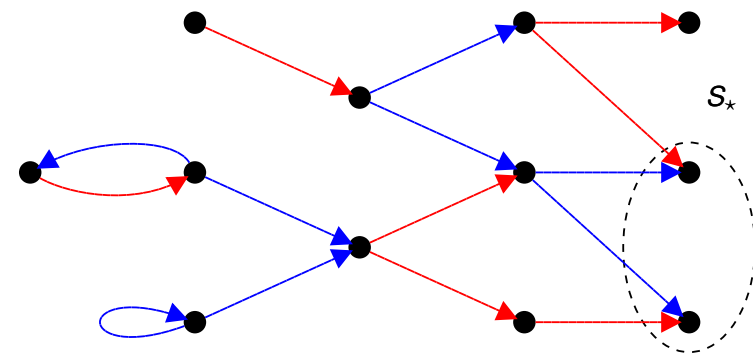


- Finds plans that may loop (strong cyclic plans).
- The algorithm is rather tricky in comparison to the algorithm for strong plans.
- Every state covered by a plan satisfies two properties:
 - 1 The state is **good**: there is at least one execution (= path in the graph defined by the plan) leading to a goal state.
 - 2 Every successor state is either a goal state or good.
- The algorithm repeatedly eliminates states that are not good.

Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

Example



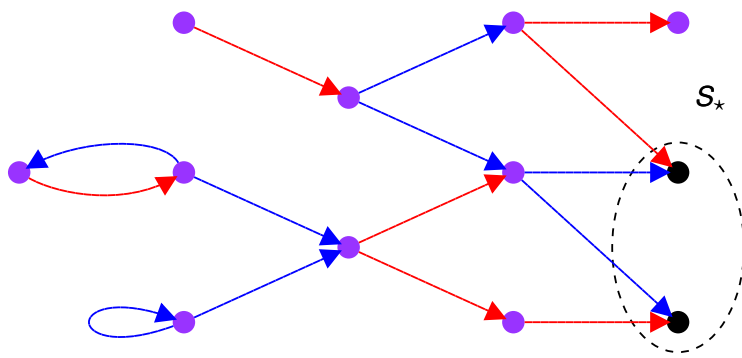
Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

Example



All states are candidates for being **good**.



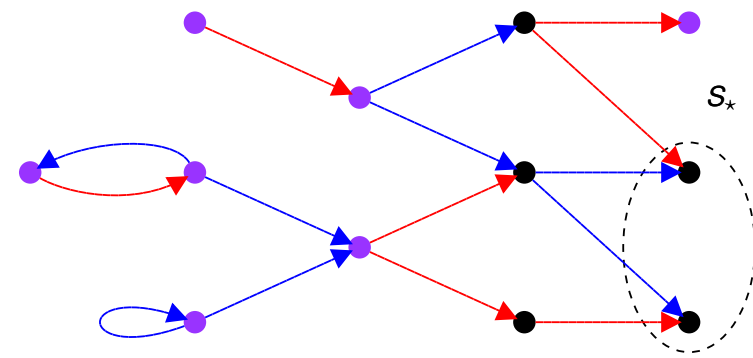
Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

Example



States from which goals are reachable in ≤ 1 steps so that all immediate successors are possibly good.



Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

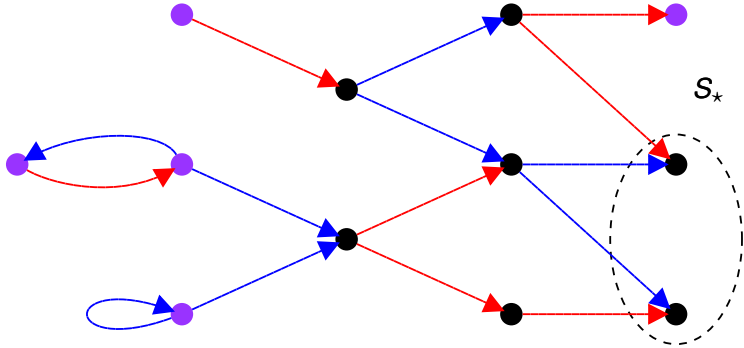
Example



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States from which goals are reachable in ≤ 2 steps so that all immediate successors are possibly good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



Nested Fixpoint Algorithm

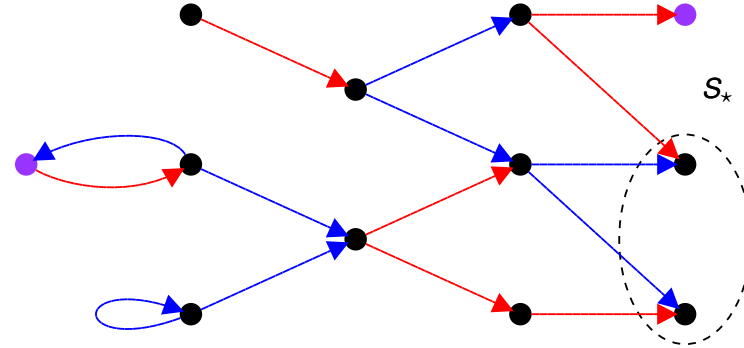
Example



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States from which goals are reachable in ≤ 3 steps so that all immediate successors are possibly good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



Nested Fixpoint Algorithm

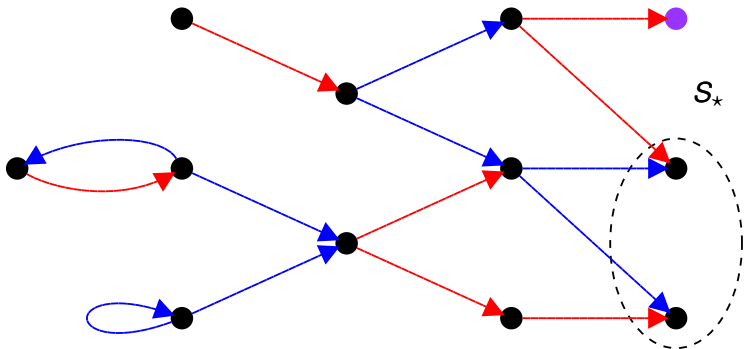
Example



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States from which goals are reachable in ≤ 4 steps so that all immediate successors are possibly good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



Nested Fixpoint Algorithm

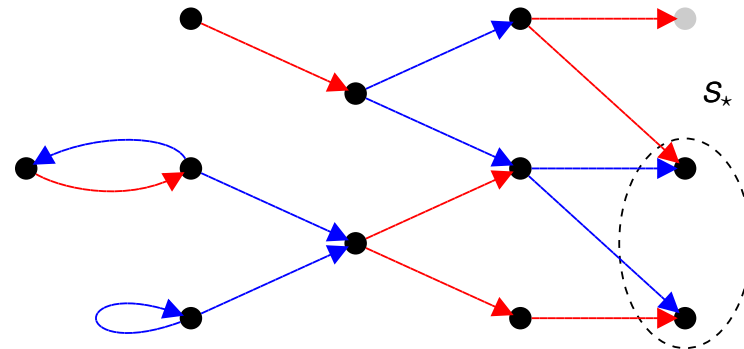
Example



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Eliminate states that turned out not to be good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



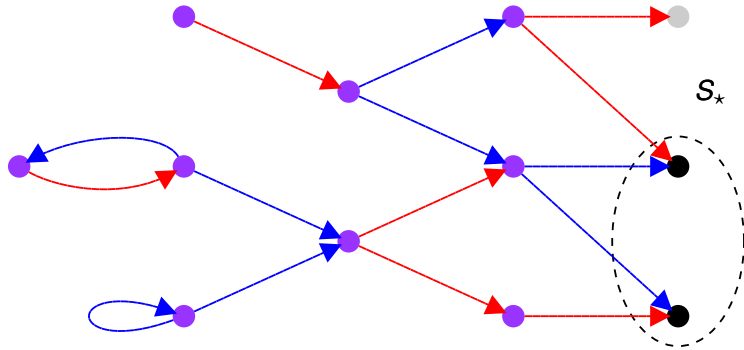
Nested Fixpoint Algorithm

Example



The set of possibly good states is now smaller.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



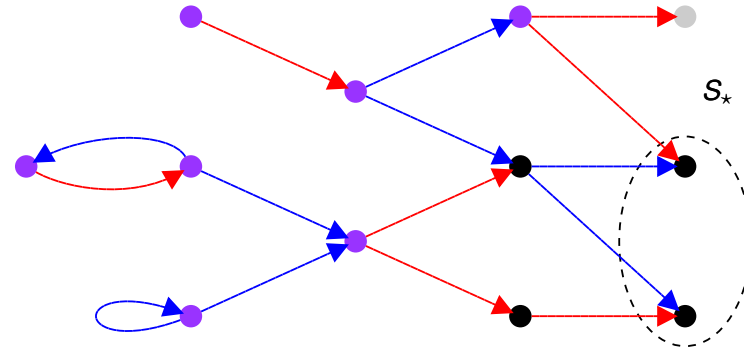
Nested Fixpoint Algorithm

Example



States from which goals are reachable in ≤ 1 steps so that all immediate successors are possibly good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



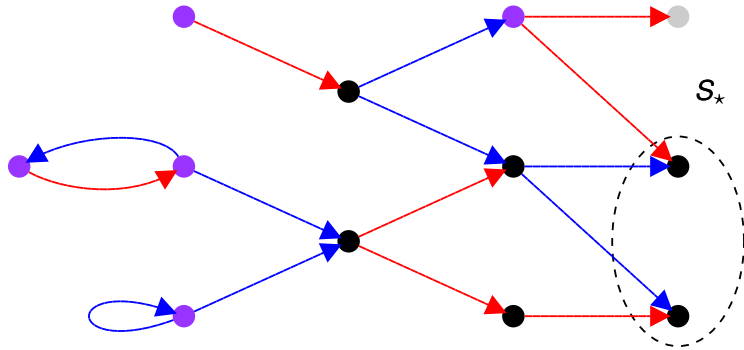
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



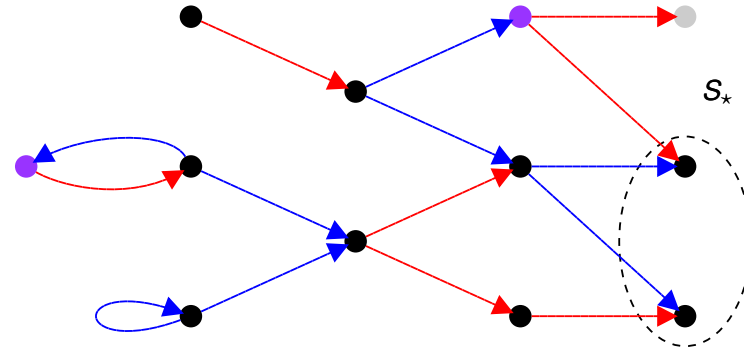
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



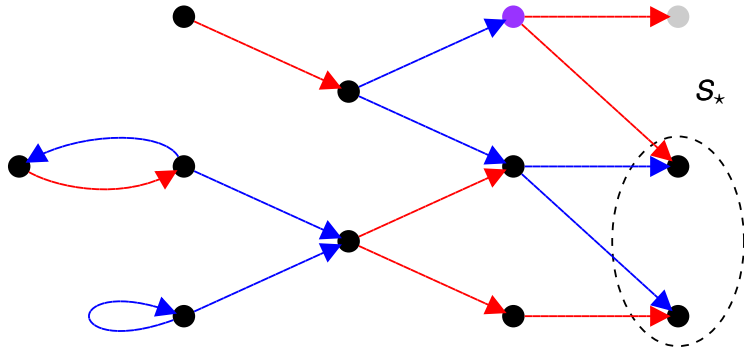
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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



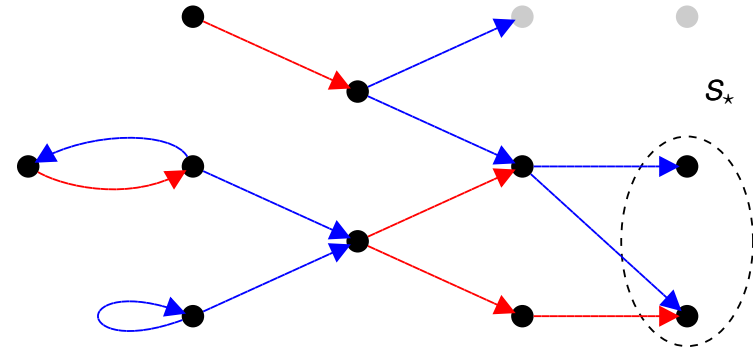
Nested Fixpoint Algorithm

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Eliminate states that turned out not to be good.

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



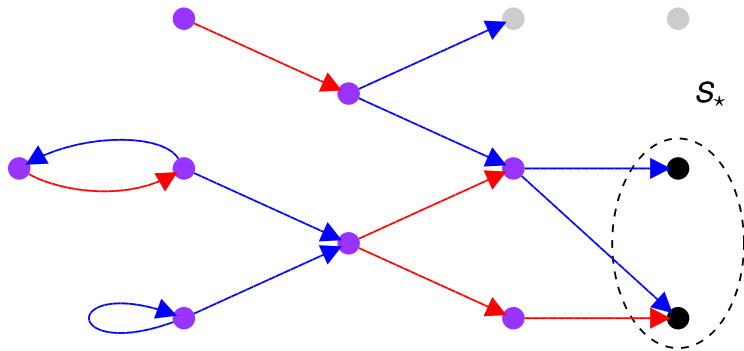
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



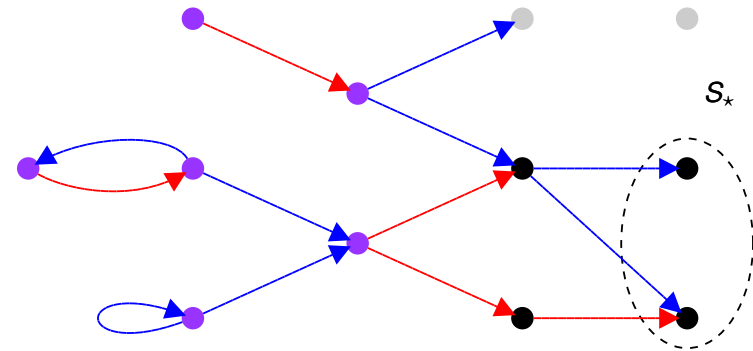
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



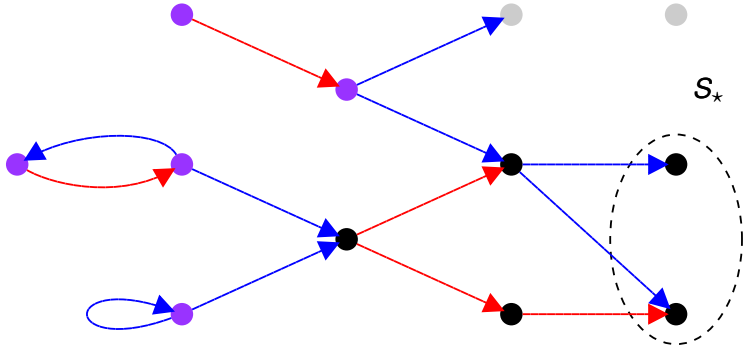
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



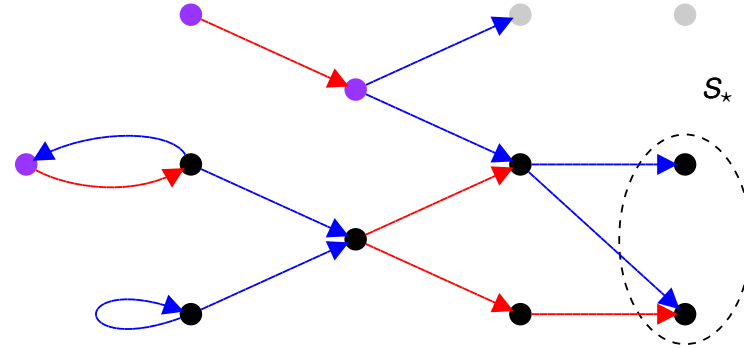
Nested Fixpoint Algorithm

Example



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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



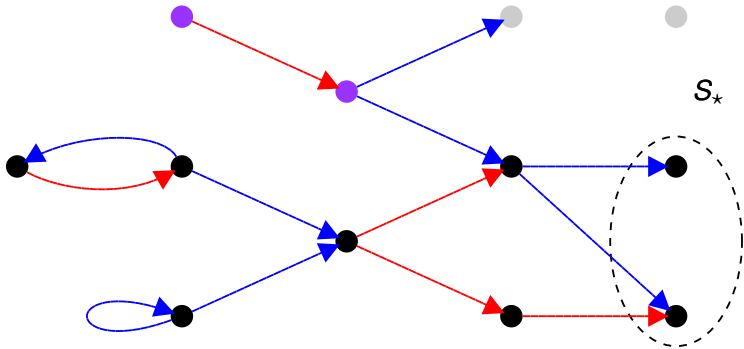
Nested Fixpoint Algorithm

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- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



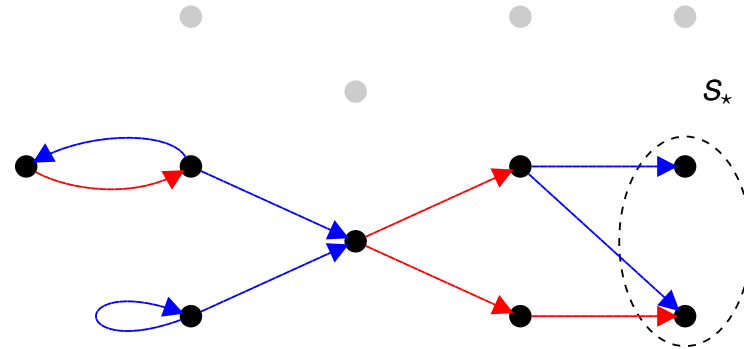
Nested Fixpoint Algorithm

Example



Remaining states are all good. A further iteration would not eliminate more states.

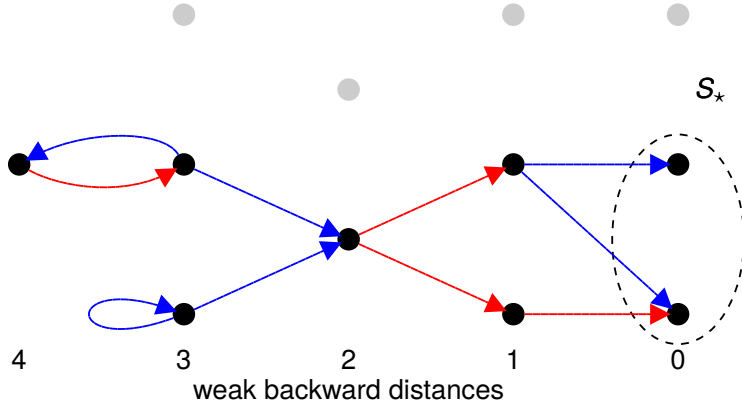
- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary



Nested Fixpoint Algorithm

Example

Assign each state an operator so that the successor states are goal states or good, and some of them are closer to goal states. Use **weak distances** computed with **weak preimages**. For this example this is trivial.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Strong cyclic plans

Recall the definition of cyclic strong plans:

Definition (strong cyclic plan)

Let S be the set of states of a planning task Π . Then a **strong cyclic plan** for Π is a function $\pi : S_\pi \rightarrow O$ for some subset $S_\pi \subseteq S$ such that

- $\pi(s)$ is applicable in s for all $s \in S_\pi$,
- $S_\pi(s_0) \subseteq S_\pi \cup S_*$ (π is closed), and
- $S_\pi(s') \cap S_* \neq \emptyset$ for all $s' \in S_\pi(s_0)$ (π is proper).

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Procedure *prune*

- The procedure **prune** finds a maximal set of states for which reaching goals with looping is possible.
- It consists of two nested loops:
 - 1 The outer loop iterates through $i = 0, 1, 2, \dots$ and produces a **shrinking** sequence of candidate good state sets C_0, C_1, \dots, C_n until $C_i = C_{i+1}$.
 - 2 The inner loop identifies **growing** sets W_j of states from which a goal state can be reached with j steps without leaving the current set of candidate good states C_i . The union of all W_0, W_1, \dots will be C_{i+1} .

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Procedure *prune*

Definition

Procedure *prune*

```

def prune(S, O, S*):
    C0 := S
    for each i in N1:
        W0 := S*
        for each j in N1:
            Wj := Wj-1 union U o in O (wpreimg_o(Wj-1) intersection spreimg_o(Ci-1))
            if Wj = Wj-1:
                break
        Ci := Wj
        if Ci = Ci-1:
            return (Ci, (W0, ..., Wj-1))
    
```

- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Procedure *prune*

Correctness



Lemma (Procedure *prune*)

Let S and $S_* \subseteq S$ be sets of states and O a set of operators. Then $\text{prune}(S, O, S_*)$ terminates after a finite number of steps and returns $C \subseteq S$ such that there is a strategy $\pi : C \setminus S_* \rightarrow O$ that is a strong cyclic plan (for the states for which it is defined) and maximal in the sense that there is no set $C' \supseteq C$ and a strong cyclic plan $\pi' : C' \setminus S_* \rightarrow O$.

- The sets W_j also returned by *prune* encode weak distances and can be used to define the strong cyclic plan π .

Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

Main algorithm



The planning algorithm

def strong-cyclic-plan($\langle V, I, O, \gamma \rangle$):

$S :=$ set of states over V

$S_* := \{s \in S \mid s \models \gamma\}$

$\langle C, (W_j)_{j=0,1,2,\dots} \rangle = \text{prune}(S, O, S_*)$

if $I \notin C$:

return no solution

for each $s \in C$:

$\delta(s) := \min\{j \in \mathbb{N}_0 \mid s \in W_j\}$

for each $s \in C \setminus S_*$:

$\pi(s) :=$ some operator $o \in O$ with $\text{img}_o(s) \subseteq C$
 and $\min\{\delta(s') \mid s' \in \text{img}_o(s)\} < \delta(s)$

return π

Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Nested Fixpoint Algorithm

Complexity



- The procedure *prune* runs in polynomial time in the number of states because the number of iterations of each loop is at most n – hence there are $O(n^2)$ iterations – and computation on each iteration takes polynomial time in the number of states.
- Finding strong cyclic plans for full observability is in the complexity class EXPTIME.
- The problem is also EXPTIME-hard.
- Similar to strong planning, we can speed up the algorithm in many practical cases by using a symbolic implementation (e. g. with BDDs).

Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Determinization-based Incremental Alg.



Idea [Kuter/Nau/Reisner/Goldman, 2008; Fu/Ng/Bastiani/Yen, 2011]:

1. Pretend the planning task was deterministic: Turn each action $o = \langle \chi, E \rangle$ with $E = \{e_1, \dots, e_n\}$ into n actions $o_i = \langle \chi, e_i \rangle$ for $i = 1, \dots, n$. Obtain classical problem Π' .
2. Find classical plan P in Π' . Add state-action mapping corresponding to P to π .
3. For each operator o_i used in P (in state s), identify original nondeterministic operator o and states $S' = \text{img}_o(s)$.
4. For each “open” state $s' \in S'$, go to 2.

Remark: May require backtracking, if some state used in a classical plan turns out not to admit a strong cyclic plan.

Strong cyclic plans
Motivation
Nested Fixpoint Algorithm
Incremental Planning Algorithm
Maintenance
Summary

Determinization-based Incremental Alg.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Definition (all-outcomes determinization)

Let $\Pi = \langle V, I, O, \gamma \rangle$ be a nondeterministic planning task. The **all-outcomes determinization** of Π is the deterministic planning task $\Pi_{det} = \langle V, I, O_{det}, \gamma \rangle$, where $O_{det} = \bigcup_{o \in O} o_{det}$, and $\langle \chi, E \rangle_{det} = \{ \langle \chi, e \rangle \mid e \in E \}$.

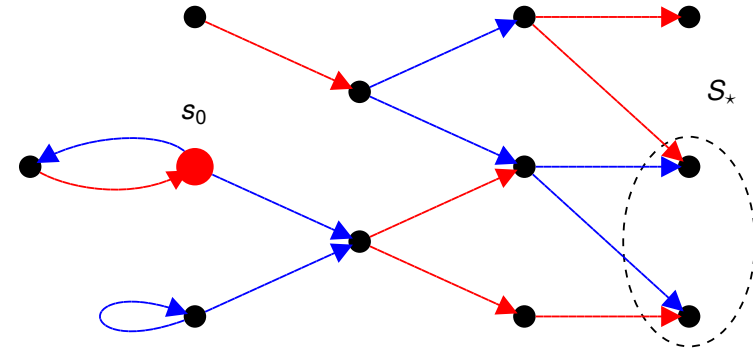
Determinization-based Incremental Alg.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Example

List of states to solve: $\{s_0\}$



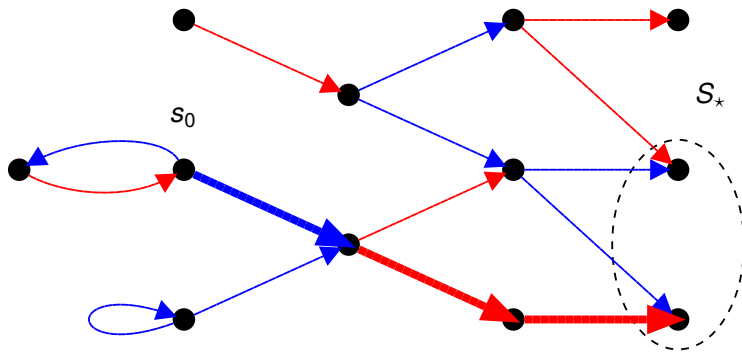
Determinization-based Incremental Alg.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Example

Plan for s_0 in determinization: *blue₂, red₂, red*



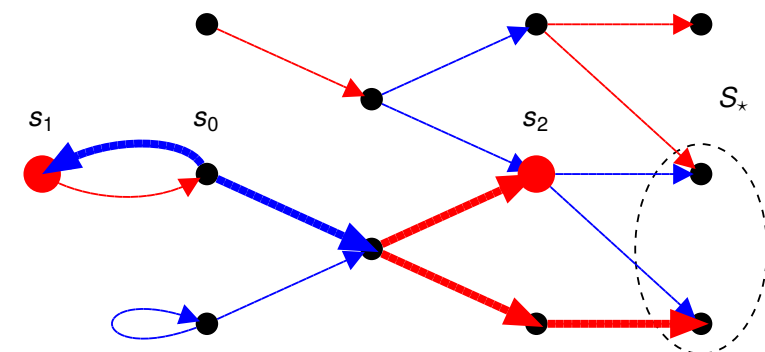
Determinization-based Incremental Alg.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Example

“Undesired” outcomes of *blue₁* and *red₁* lead to new list of states to solve: $\{s_1, s_2\}$

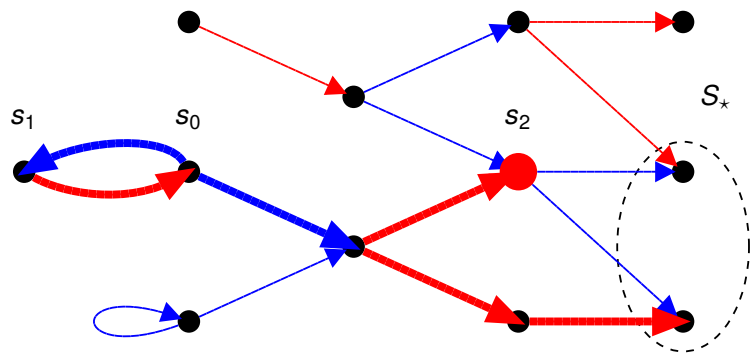


Determinization-based Incremental Alg.

Example



Plan for s_1 in determinization: *red, blue₂, red₂, red*



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

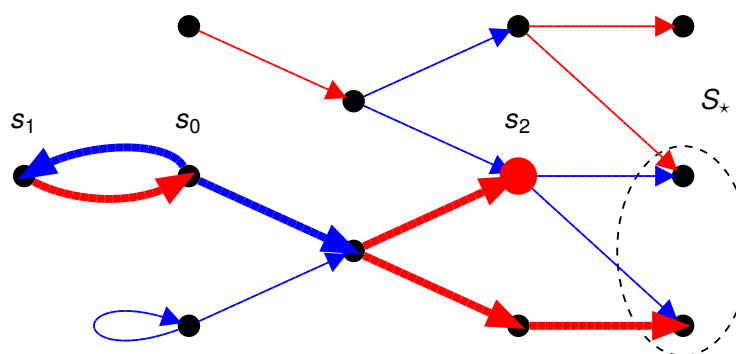
Determinization-based Incremental Alg.

Example



No new “undesired” outcomes.

List of states to solve: $\{s_2\}$



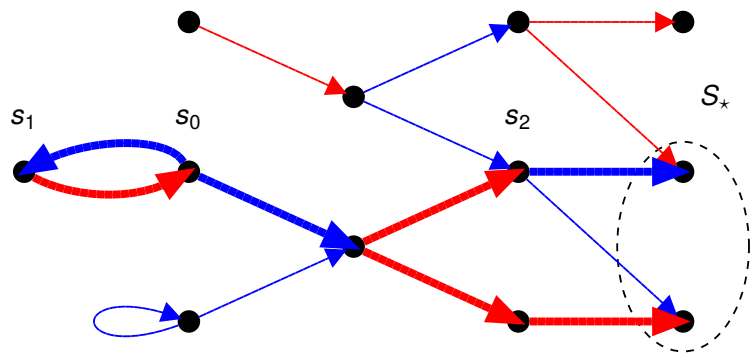
- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Determinization-based Incremental Alg.

Example



Plan for s_2 in determinization: *blue₁*



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

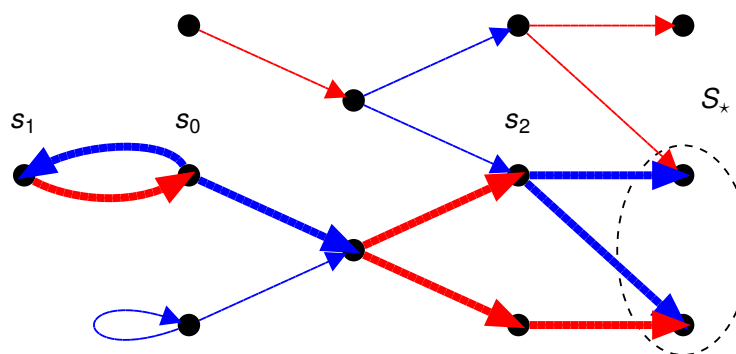
Determinization-based Incremental Alg.

Example



“Undesired” outcome of *blue₂* in s_2 leads to goal state, too.

List of states to solve: \emptyset . Strong cyclic plan found.



- Strong cyclic plans
- Motivation
- Nested Fixpoint Algorithm
- Incremental Planning Algorithm
- Maintenance
- Summary

Procedure incremental-strong-cyclic-plan

def incremental-strong-cyclic-plan($\langle V, I, O, \gamma \rangle$):

$\pi \leftarrow \emptyset$; $fail \leftarrow \{I\}$

while $fail \neq \emptyset$:

$s \leftarrow \text{SELECTANDREMOVEFROM}(fail)$

$\pi' \leftarrow \text{DETSEARCH}(\langle V, s, O_{det}, \gamma \rangle)$

if $\pi' = \text{FAILURE}$:

if $s = I$: **return** FAILURE

else: $\text{BACKTRACK}(s, \pi, \langle V, I, O, \gamma \rangle)$

else:

$\pi \leftarrow \pi \cup \pi'$

$fail \leftarrow \{s \in S \mid s \text{ nongoal state reachable from } I \text{ following } \pi, \text{ but } \pi(s) \text{ undefined}\}$

return π

If a deterministic search fails, the state s from which it started cannot be part of a strong cyclic plan.

- If $s = I$, the whole given planning problem is unsolvable and the algorithm returns FAILURE.
- Otherwise, state s , which has already been added to the constructed policy π , has to be removed from π , and the algorithm has to ensure that s will never be reconsidered again. This is accomplished by the procedure BACKTRACK.

Procedure backtrack

def backtrack($s, \pi, \langle V, I, O, \gamma \rangle$):

update π by deleting all entries that would immediately lead to s , i.e. $\pi \leftarrow \pi \setminus \{(s', \pi(s')) \mid s \in \text{img}_{\pi(s')}(s')\}$

add all states s' removed from π to the set of fail-states $fail$

permanently mark all formerly assigned actions $\pi(s')$ removed from π at s' as inapplicable in s' to avoid running into the same dead end again.

- Iteratively solves all-outcomes determinizations of Π with “fail-states” as initial states.
- Planner can choose desired outcome of each action.
- Deterministic plans are added to policy under construction.
- Corresponding undesired outcomes have to be added to the set of “fail-states” $fail$.
- Deterministic plans for “fail-states” are constructed until no more “fail-states” remain.
- Eventually, the algorithm either returns a strong cyclic plan or FAILURE if no such plan exists.

Theorem

Procedure incremental-strong-cyclic-plan, called with task Π , returns a strong cyclic plan for Π iff such a plan exists, and FAILURE, otherwise.

- Can use any classical planner for deterministic searches.
- Can benefit from heuristics etc. used there.
- Classical planner can be configured to prefer short solutions or solutions using deterministic actions induced by nondeterministic actions with few different outcomes (likely fewer new “fail-states”).

- When to terminate a deterministic sub-search?
 - At goal states?
 - At states currently part of the partial solution?
 - At parent of currently solved “fail-state”?

This can make a huge difference.

- Similarly: Where should the heuristic guide the classical planner? Goals, partial solution, parent node?
- Additional marking of nodes as definitely solved if this can be detected.
- State reuse between subsequent classical planner calls.
- Generalization of solved states by regression search from goal along weak (deterministic) plan (cf. [Muisse/McIlraith/Beck, 2012]).

- Definition
- Example
- Algorithm

Maintenance goals



- In this lecture, we usually limit ourselves to the problem of finding plans that **reach a goal state**.
- In practice, planning is often about more general goals, where execution cannot be terminated.
 - 1 An animal: find food, eat, sleep, find food, eat, sleep, ...
 - 2 Cleaning robot: keep the building clean.
- These problems cannot be directly formalized in terms of reachability because infinite (unbounded) plan execution is needed.
- We do not discuss this topic in full detail. However, to give at least a little impression of **planning for temporally extended goals**, we will discuss the simplest objective with infinite plan executions: **maintenance**.

Strong cyclic plans
Maintenance
Definition
Example
Algorithm
Summary

Plan objectives

Maintenance



Definition

Let $\mathcal{T} = \langle V, I, O, \gamma \rangle$ be a planning task with state set S and set of goal states $S_* = \{s \in S \mid s \models \gamma\}$.

A strategy π for \mathcal{T} is called a **plan for maintenance** for \mathcal{T} iff

- $\pi(s)$ is applicable in s for all $s \in S_\pi$,
- $S_\pi(s_0) \subseteq S_\pi$, and
- $S_\pi(s_0) \subseteq S_*$.

Strong cyclic plans
Maintenance
Definition
Example
Algorithm
Summary

Maintenance goals

Example

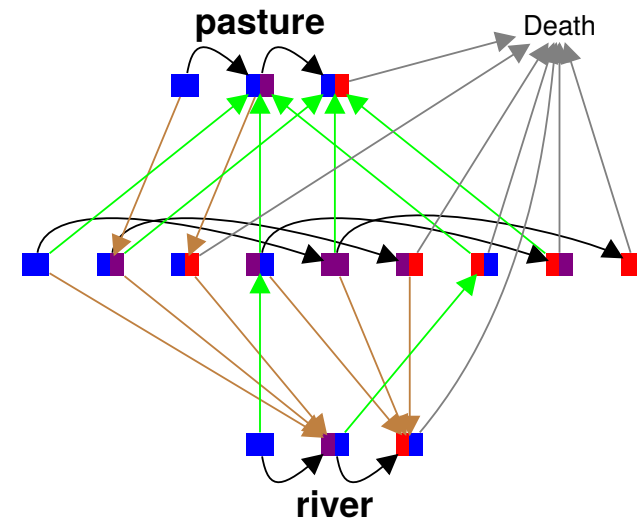


- The state of an animal is determined by three state values: hunger (0, 1, 2), thirst (0, 1, 2) and location (river, pasture, desert). There is also a special state called **death**.
- Thirst grows when not at river; at river it is 0.
- Hunger grows when not on pasture; on pasture it is 0.
- If hunger or thirst exceeds 2, the animal dies.
- The goal of the animal is to avoid death.

Strong cyclic plans
Maintenance
Definition
Example
Algorithm
Summary

Maintenance goals

Transition system for the example 0-safe states 1-safe states i -safe states for all $i \geq 2$



Strong cyclic plans
Maintenance
Definition
Example
Algorithm
Summary

Maintenance goals

Plan for the example



- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

We can infer rules backwards starting from the death condition.

- 1 If in desert and **thirst = 2**, must go to river.
- 2 If in desert and **hunger = 2**, must go to pasture.
- 3 If on pasture and **thirst = 1**, must go to desert.
- 4 If at river and **hunger = 1**, must go to desert.

If the above rules conflict, the animal will die.

Algorithm for maintenance goals

Idea



- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

Summary of the algorithm idea

Repeatedly eliminate from consideration those states that in one or more steps unavoidably lead to a non-goal state.

- A state is ***i*-safe** iff there is a plan that guarantees “survival” for the next *i* actions.
- A state is **safe** (or **∞ -safe**) iff it is *i*-safe for all $i \in \mathbb{N}_0$.
- The **0-safe** states are exactly the goal states: maintenance objective is satisfied for the current state.
- Given all *i*-safe states, compute all *i* + 1-safe states by using strong preimages.
- For some $i \in \mathbb{N}_0$, *i*-safe states equal *i* + 1-safe states because there are only finitely many states and at each step and *i* + 1-safe states are a subset of *i*-safe states. Then *i*-safe states are also ∞ -safe.

Algorithm for maintenance goals

Algorithm



- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

Planning for maintenance goals

def maintenance-plan((V, I, O, γ)):

$S :=$ set of states over V

$Safe_0 := \{s \in S \mid s \models \gamma\}$

for each $i \in \mathbb{N}_1$:

$Safe_i := Safe_{i-1} \cap \bigcup_{o \in O} spreimg_o(Safe_{i-1})$

if $Safe_i = Safe_{i-1}$:

break

if $I \notin Safe_i$:

return no solution

for each $s \in Safe_i$:

$\pi(s) :=$ some operator $o \in O$ with $img_o(s) \subseteq Safe_i$

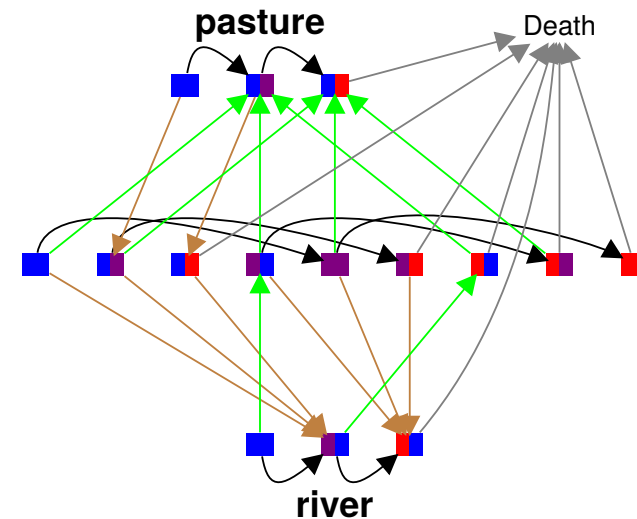
return π

Maintenance goals

Transition system for the example 0-safe states 1-safe states *i*-safe states for all $i \geq 2$



- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

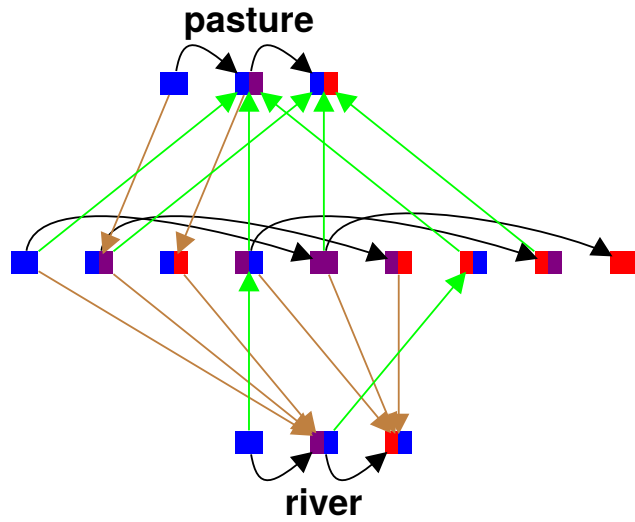


Maintenance goals

Transition system for the example 0-safe states 1-safe states i -safe states for all $i \geq 2$



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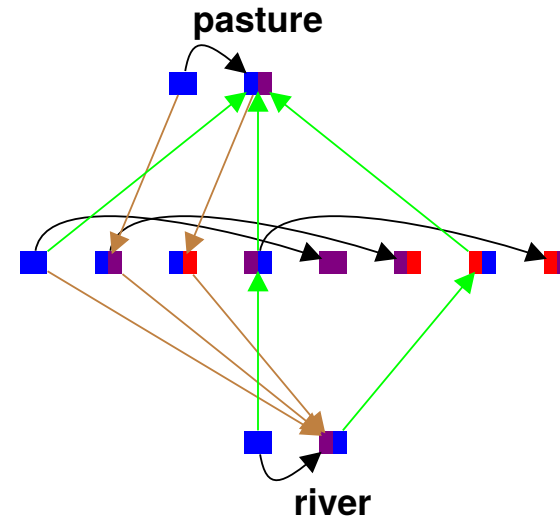
- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

Maintenance goals

Transition system for the example 0-safe states 1-safe states i -safe states for all $i \geq 2$



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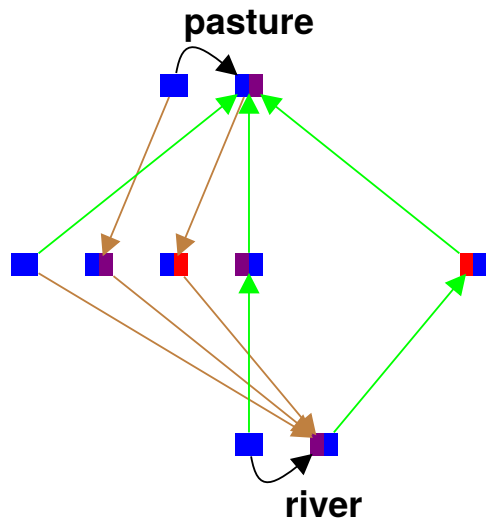
- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

Maintenance goals

Transition system for the example 0-safe states 1-safe states i -safe states for all $i \geq 2$



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- Strong cyclic plans
- Maintenance
- Definition
- Example
- Algorithm
- Summary

3 Summary



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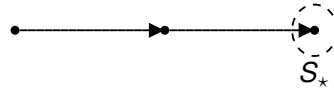
- Strong cyclic plans
- Maintenance
- Summary

Different planning objectives



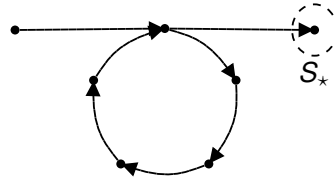
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Strong planning



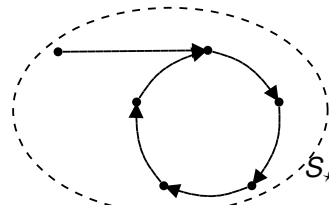
Strong cyclic plans

Strong cyclic planning



Maintenance

Maintenance



Summary

Outlook: Computational tree logic



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- We have considered different classes of solutions for planning tasks by defining **different planning problems**.
 - strong planning problem: find a strong plan
 - strong cyclic planning problem: find a strong cyclic plan
 - ...
- Alternatively, we could allow specifying goals in a **modal logic** like **computational tree logic** to directly express the type of plan we are interested in using **modalities** such as A (all), E (exists), G (globally), and F (finally).
 - Weak planning: $EF\varphi$
 - Strong planning: $AF\varphi$
 - Strong cyclic planning: $AGEF\varphi$
 - Maintenance: $AG\varphi$

Strong cyclic plans

Maintenance

Summary

Summary



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- We have extended our earlier planning algorithm from **strong plans** to **strong cyclic plans**.
- The story does not end there: When considering infinitely executing plans, many more types of goals are feasible.
- We considered **maintenance** as a simple example of a **temporally extended goal**.
- In general, temporally extended goals be expressed in **modal logics** such as computational tree logic (CTL).
- We presented dynamic programming (backward search) algorithms for strong cyclic and maintenance planning.
- In practice, one might implement both algorithms by using binary decision diagrams (BDDs) as a data structure for state sets.

Strong cyclic plans

Maintenance

Summary