

# Qualitative Spatial and Temporal Reasoning based on decomposition

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# Qualitative Spatial and Temporal Reasoning

- Continuous space and time: undecidable;
- Discretization of space: decidable;
- Representation of relations between regions and/or timeframe.

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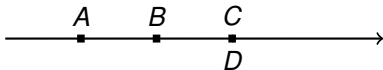


Figure: Points on a line

## Definition

The set of possible relations is:  $\mathcal{B} = \{<, =, >\}$

## Example

Relations in Figure 1:  $A < B$ ,  $C = D$ ,  $C > D$ .

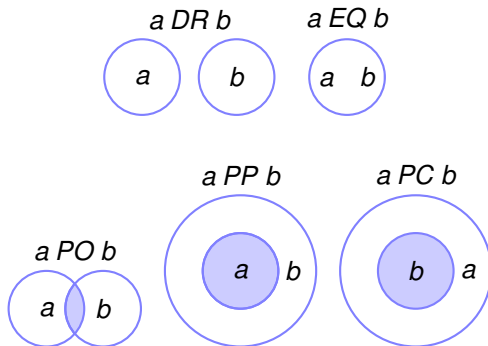


Figure: Open regions over a regular space

### Definition

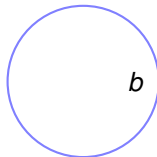
The set of possible relations is:  $\mathcal{B} = \{DR, EQ, PO, PP, PC\}$

# The Problem - Qualitative Constraint Networks (1)

## Definition

Lets consider a set  $\mathcal{C}$  of known relation. Is there any real world instantiation than satisfy all the relations in  $\mathcal{C}$ ?

$a DR b$



# The Problem - Qualitative Constraint Networks (2)

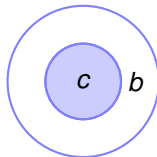
## Definition

Lets consider a set  $\mathcal{C}$  of known relation. Is there any real world instantiation than satisfy all the relations in  $\mathcal{C}$ ?

$a DR b$

$a DR c$

$b PC c$



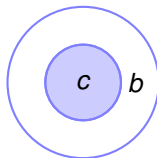
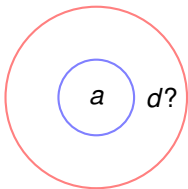


# The Problem - Qualitative Constraint Networks (3a)

## Definition

Lets consider a set  $\mathcal{C}$  of known relation. Is there any real world instantiation than satisfy all the relations in  $\mathcal{C}$ ?

$a DR b$   
 $a DR c$   
 $b PC c$   
 $a PP, PO d$

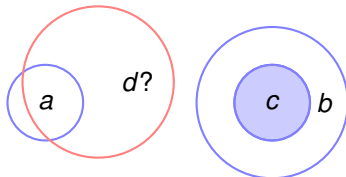


# The Problem - Qualitative Constraint Networks (3b)

## Definition

Lets consider a set  $\mathcal{C}$  of known relation. Is there any real world instantiation than satisfy all the relations in  $\mathcal{C}$ ?

$a DR b$   
 $a DR c$   
 $b PC c$   
 $a PP, PO d$

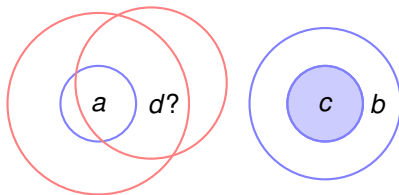


# The Problem - Qualitative Constraint Networks (4)

## Definition

Lets consider a set  $\mathcal{C}$  of known relation. Is there any real world instantiation than satisfy all the relations in  $\mathcal{C}$ ?

$a DR b$   
 $a DR c$   
 $b PC c$   
 $a PP, PO d$   
 $c EQ d$

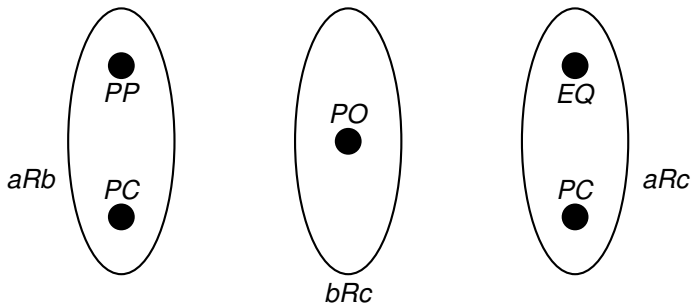


# The Composition Table

$\circ$	$DR_{cd}$	$PO_{cd}$	$EQ_{cd}$	$PP_{cd}$	$PC_{cd}$
$DR_{ac}$	*	$DR\ PO$ $PP$	$DR$	$DR\ PO$ $PP$	$DR$
$PO_{ac}$	$DR\ PO$ $PC$	*	$PO$	$PO\ PP$	$DR\ PO$ $PC$
$EQ_{ac}$	$DR$	$PO$	$EQ$	$PP$	$PC$
$PP_{ac}$	$DR$	$DR\ PO$ $PP$	$PP$	$PP$	*
$PC_{ac}$	$DR$ $PO\ PC$	$PO\ PC$	$PC$	$PO\ EQ$ $PP\ PC$	$PC$

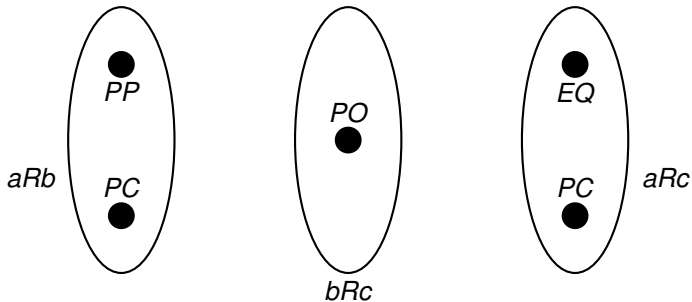
# What is usually done

Solving QCNs as CSPs with the composition table as global constraint (1)



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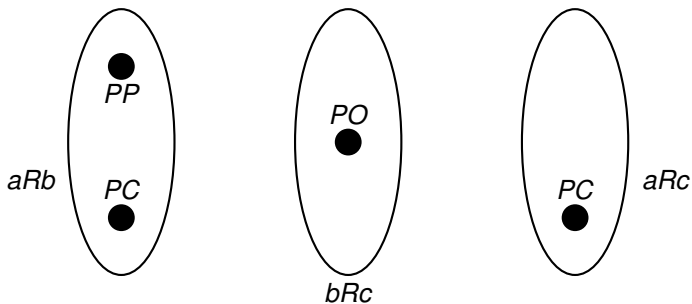
Solving QCNs as CSPs with the composition table as global constraint (1)



$$PP \circ PO = \{DR, PO, PP\} \quad PC \circ PO = \{PO, PC\}$$

# What is usually done

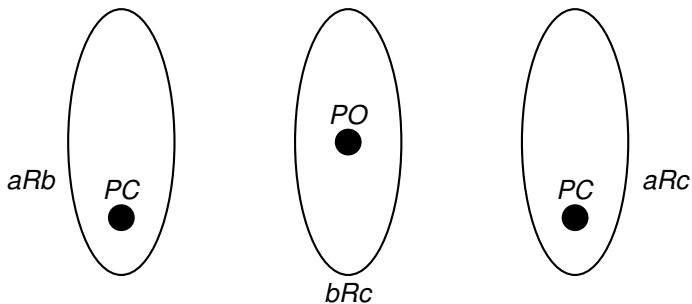
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# What is usually done

Solving QCNs as CSPs with the composition table as global constraint (1)



$$PP \circ PO = \{DR, PO, PP\} \quad PC \circ PO = \{PO, PC\}$$



# What is usually done

Solving QCNs as CSPs with the composition table as global constraint (2)

## **Backtracking**( $\mathcal{C}, a$ ):

---

*Input:* a constraint network  $\mathcal{C} = V, D, C$  and  
a partial assignment  $a$  of  $\mathcal{C}$

*Output:* a solution of  $\mathcal{C}$  or “inconsistent”

**if**  $a$  is not consistent with  $\mathcal{C}$ : **return** “inconsistent”

**if**  $a$  is defined for all variables in  $V$ : **return**  $a$

select some variable  $v_i$  for which  $a$  is not defined

**for** each value  $x$  from  $D_i$ :

$a' := a \cup \{v_i \mapsto x\}$

$a'' := \text{ApplyCompositionTable}(\mathcal{C}, a')$

$a''' \leftarrow \text{Backtracking}(\mathcal{C}, a'')$

**if**  $a'''$  is not “inconsistent”: **return**  $a'''$

**return** “inconsistent”

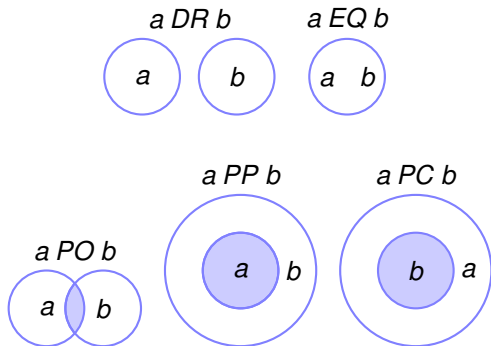
# The axiomatic point of view (1)

Instead of considering the composition aspect, we deal with the semantic and divide relations in pieces:

- $a C b$  :=  $a$  and  $b$  share at least a common point;
- $a P b$  := all points of  $a$  are points of  $b$ .

Then

	$a C b$	$a P b$	$b P a$
$DR$	×	×	×
$PO$	○	×	×
$EQ$	○	○	○
$PP$	○	○	×
$PC$	○	×	○



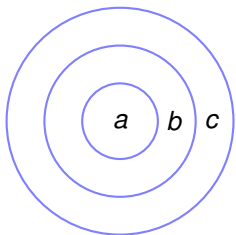
## The axiomatic point of view (2)

Instead of considering the composition aspect, we deal with the semantic and divide relations in pieces:

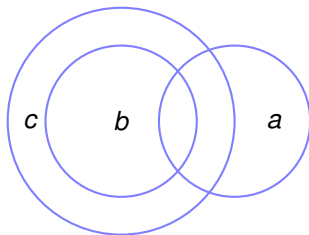
- $a C b := a$  and  $b$  share at least a common point;
- $a P b :=$  all points of  $a$  are points of  $b$ .

Then

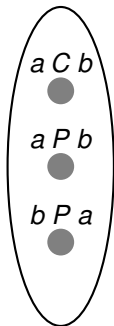
$$\frac{a P b \quad b P c}{a P c}$$



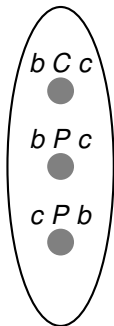
$$\frac{a C b \quad b P c}{a C c}$$



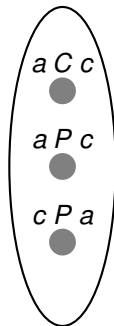
# The axiomatic point of view (3)



$a\{PP, PC\}b$

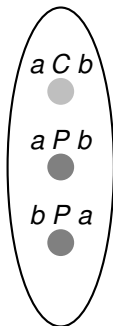


$b\{PO\}c$



$a\{EQ, PC\}c$

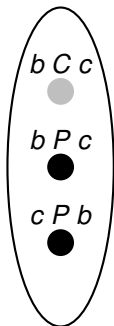
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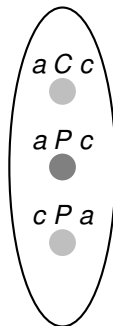
$PP = \circ \circ \times$

$PC = \circ \times \circ$



$b\{PO\}c$

$PO = \circ \times \times$

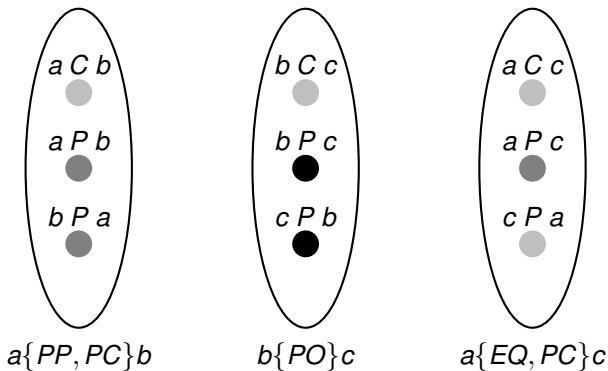


$a\{EQ, PC\}c$

$EQ = \circ \circ \circ$

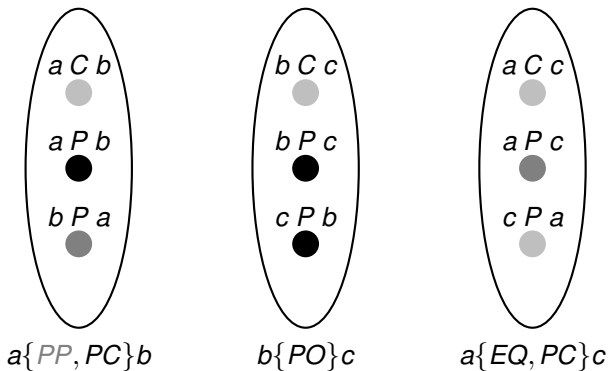
$PC = \circ \times \circ$

# The axiomatic point of view (3)



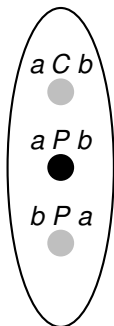
$$cPa + \neg cPb \rightarrow \neg aPb$$

# The axiomatic point of view (3)

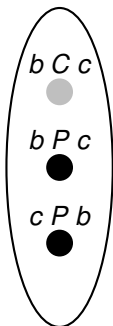


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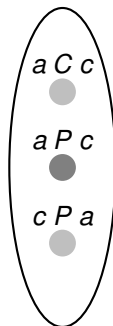
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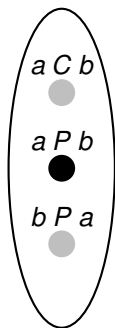
$b\{PO\}c$



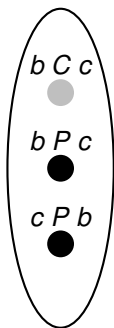
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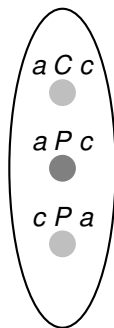
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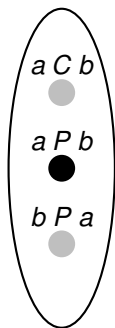
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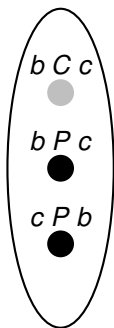
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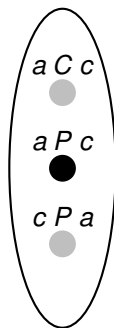
# The axiomatic point of view (3)



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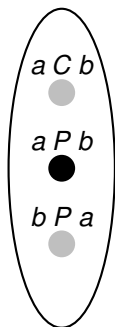
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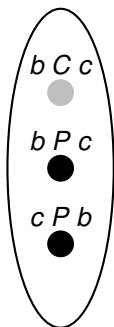
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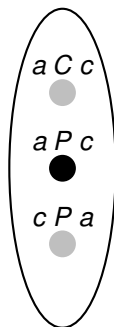
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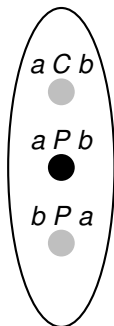
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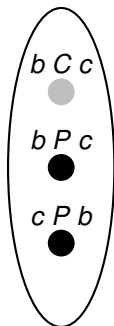
$a\{EQ,PC\}c$

$$bPa + \neg bPc \rightarrow \neg aPc$$

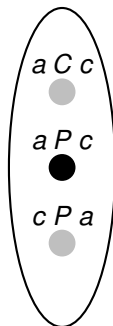
# The axiomatic point of view (3)



$a\{PC\}b$



$b\{PO\}c$



$a\{PC\}c$

# What we are doing

Solving QCNs as SATs with axioms as global constraints

**Backtracking**( $\mathcal{C}, a$ ):

---

$a'' := \text{ApplyAxioms}(\mathcal{C}, a')$   
**if**  $a$  is not consistent with  $\mathcal{C}$ : **return** “inconsistent”  
**if**  $a$  is defined for all variables in  $V$ : **return**  $a$   
select some variable  $v_i$  which is not defined  
 $a' := a \cup \{v_i \mapsto \top\}$   
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