Principles of AI Planning

13. Computational complexity of classical planning

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How hard is planning?

- We have seen that planning can be done in time polynomial in the size of the transition system.
- However, we have not seen algorithms which are polynomial in the input size (size of the task description).

$\rightsquigarrow$ What is the precise computational complexity of the planning problem?

Why computational complexity?

- understand the problem
- know what is not possible
- find interesting subproblems that are easier to solve
- distinguish essential features from syntactic sugar
  - Is STRIPS planning easier than general planning?
  - Is planning for FDR tasks harder than for propositional tasks?

Nondeterministic Turing machines

Definition (nondeterministic Turing machine)

A nondeterministic Turing machine (NTM) is a 6-tuple $\langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ with the following components:

- input alphabet Σ and blank symbol $\Box \notin \Sigma$
  - alphabets always nonempty and finite
  - tape alphabet $\Sigma_\Box = \Sigma \cup \{\Box\}$
- finite set Q of internal states with initial state $q_0 \in Q$ and accepting state $q_Y \in Q$
  - nonterminal states $Q' := Q \setminus \{q_Y\}$
- transition relation $\delta \subseteq (Q' \times \Sigma_\Box) \times (Q \times \Sigma_\Box \times \{-1, +1\})$

Deterministic Turing machines

Definition (deterministic Turing machine)

A deterministic Turing machine (DTM) is an NTM where the transition relation is functional, i. e., for all $\langle q, a \rangle \in Q' \times \Sigma_{\square}$, there is exactly one triple $\langle q', a', \Delta \rangle$ with $\langle \langle q, a \rangle, \langle q', a', \Delta \rangle \rangle \in \delta$.

Notation: We write $\delta(q, a)$ for the unique triple $\langle q', a', \Delta \rangle$ such that $\langle \langle q, a \rangle, \langle q', a', \Delta \rangle \rangle \in \delta$.

Turing machine configurations

Definition (Configuration)

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be an NTM.

A configuration of M is a triple $\langle w, q, x \rangle \in \Sigma_\Box^* \times Q \times \Sigma_\Box^+$.

- w: tape contents before tape head
- q: current state
- x: tape contents after and including tape head

Turing machine transitions

Definition (yields relation)

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be an NTM.

A configuration c of M **yields** a configuration c' of M, in symbols $c \vdash c'$, as defined by the following rules, where $a, a', b \in \Sigma_\square$, $w, x \in \Sigma_\square^*$, $q, q' \in Q$ and $\langle \langle q, a \rangle, \langle q', a', \Delta \rangle \rangle \in \delta$:

$$
\begin{aligned}
(w, q, ax) &\vdash (wa', q', x) && \text{if } \Delta = +1, |x| \geq 1 \\
(w, q, a) &\vdash (wa', q', \square) && \text{if } \Delta = +1 \\
(wb, q, ax) &\vdash (w, q', ba'x) && \text{if } \Delta = -1 \\
(\epsilon, q, ax) &\vdash (\epsilon, q', \square a'x) && \text{if } \Delta = -1
\end{aligned}
$$

Accepting configurations

Definition (accepting configuration, time)

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be an NTM,
let $c = \langle w, q, x \rangle$ be a configuration of M, and let $n \in \mathbb{N}_0$.

- If $q = q_Y$, M accepts c in time n.
- If $q \neq q_Y$ and M accepts some c' with $c \vdash c'$ in time n, then M accepts c in time $n + 1$.

Definition (accepting configuration, space)

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be an NTM,
let $c = \langle w, q, x \rangle$ be a configuration of M, and let $n \in \mathbb{N}_0$.

- If $q = q_Y$ and $|w| + |x| \leq n$, M accepts c in space n.
- If $q \neq q_Y$ and M accepts some c' with $c \vdash c'$ in space n, then M accepts c in space n.

Accepting words and languages

Definition (accepting words)

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be an NTM.

M accepts the word $w \in \Sigma^*$ in time (space) $n \in \mathbb{N}_0$
iff M accepts (ϵ, q_0, w) in time (space) n.

- Special case: M accepts ϵ in time (space) $n \in \mathbb{N}_0$
  iff M accepts $(\epsilon, q_0, \Box)$ in time (space) n.

Definition (accepting languages)

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be an NTM, and let $f : \mathbb{N}_0 \to \mathbb{N}_0$.

M accepts the language $L \subseteq \Sigma^*$ in time (space) f
iff M accepts each word $w \in L$ in time (space) $f(|w|)$,
and M does not accept any word $w \notin L$ (in any time/space).

Time and space complexity classes

Definition (DTIME, NTIME, DSPACE, NSPACE)

Let $f : \mathbb{N}_0 \to \mathbb{N}_0$.

Complexity class $\mathrm{DTIME}(f)$ contains all languages accepted in time f by some DTM.

Complexity class $\mathrm{NTIME}(f)$ contains all languages accepted in time f by some NTM.

Complexity class $\mathrm{DSPACE}(f)$ contains all languages accepted in space f by some DTM.

Complexity class $\mathrm{NSPACE}(f)$ contains all languages accepted in space f by some NTM.

Polynomial time and space classes

Let $\mathcal{P}$ be the set of polynomials $p : \mathbb{N}_0 \to \mathbb{N}_0$ whose coefficients are natural numbers.

Definition (P, NP, PSPACE, NPSPACE)

$$\begin{aligned}
\mathsf{P} &= \textstyle\bigcup_{p \in \mathcal{P}} \mathsf{DTIME}(p) \\
\mathsf{NP} &= \textstyle\bigcup_{p \in \mathcal{P}} \mathsf{NTIME}(p) \\
\mathsf{PSPACE} &= \textstyle\bigcup_{p \in \mathcal{P}} \mathsf{DSPACE}(p) \\
\mathsf{NPSPACE} &= \textstyle\bigcup_{p \in \mathcal{P}} \mathsf{NSPACE}(p)
\end{aligned}$$

Polynomial complexity class relationships

Theorem (complexity class hierarchy)

$\mathrm{P} \subseteq \mathrm{NP} \subseteq \mathrm{PSPACE} = \mathrm{NPSPACE}$

Proof.

$\mathrm{P} \subseteq \mathrm{NP}$ and $\mathrm{PSPACE} \subseteq \mathrm{NPSPACE}$ is obvious because deterministic Turing machines are a special case of nondeterministic ones.

$\mathrm{NP} \subseteq \mathrm{NPSPACE}$ holds because a Turing machine can only visit polynomially many tape cells within polynomial time.

$\mathrm{PSPACE} = \mathrm{NPSPACE}$ is a special case of a classical result known as Savitch's theorem (Savitch 1970). □

The propositional planning problem

Definition (plan existence)

The plan existence problem (PlanEx)
is the following decision problem:

Given: Planning task Π
Question: Is there a plan for Π?

$\rightsquigarrow$ decision problem analogue of satisficing planning

Definition (bounded plan existence)

The bounded plan existence problem (PlanLen)
is the following decision problem:

Given: Planning task Π, length bound $K \in \mathbb{N}_0$
Question: Is there a plan for Π of length at most K?

$\rightsquigarrow$ decision problem analogue of optimal planning

Plan existence vs. bounded plan existence

Theorem (reduction from PlanEx to PlanLen)

PlanEx $\leq_p$ PlanLen

Proof.

A propositional planning task with n state variables has a plan iff it has a plan of length at most $2^n - 1$.

$\rightsquigarrow$ map instance Π of PlanEx to instance $\langle \Pi, 2^n - 1 \rangle$ of PlanLen, where n is the number of n state variables of Π

$\rightsquigarrow$ polynomial reduction $\square$

Membership in PSPACE

Theorem (PSPACE membership for PlanLen)

PlanLen $\in$ PSPACE

Proof.

Show PlanLen $\in$ NPSPACE and use Savitch's theorem.

Nondeterministic algorithm:

def plan($\langle A, I, O, G \rangle$, K):
    $s := I$
    $k := K$
    while $s \not\models G$:
        guess $o \in O$
        fail if o not applicable in s **or** k $= 0$
        $s := app_o(s)$
        $k := k - 1$
    accept

□

Hardness for PSPACE

Idea: generic reduction

- For an arbitrary fixed DTM M with space bound polynomial p and input w, generate planning task which is solvable iff M accepts w in space $p(|w|)$.
- For simplicity, restrict to TMs which never move to the left of the initial head position (no loss of generality).

Reduction: state variables

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{1, \dots, p(n)\}$.

State variables

- state_q for all $q \in Q$
- head_i for all $i \in X \cup \{0, p(n) + 1\}$
- $\text{content}_{i,a}$ for all $i \in X$, $a \in \Sigma_\square$

$\rightsquigarrow$ allows encoding a Turing machine configuration

Reduction: initial state

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{1, \dots, p(n)\}$.

Initial state

Initially true:

- state_{q_0}
- head_1
- $\mathrm{content}_{i,w_i}$ for all $i \in \{1, \dots, n\}$
- $\mathrm{content}_{i,\square}$ for all $i \in X \setminus \{1, \dots, n\}$

Initially false:

- all others

Reduction: operators

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{1, \dots, p(n)\}$.

Operators

One operator for each transition rule $\delta(q, a) = \langle q', a', \Delta \rangle$ and each cell position $i \in X$:

- precondition: $\text{state}_q \land \text{head}_i \land \text{content}_{i,a}$
- effect: $\neg\text{state}_q \land \neg\text{head}_i \land \neg\text{content}_{i,a}$
  $\land\, \text{state}_{q'} \land \text{head}_{i+\Delta} \land \text{content}_{i,a'}$
  - If $q = q'$ and/or $a = a'$, omit the effects on state_q and/or $\text{content}_{i,a}$, to avoid consistency condition issues.

Reduction: goal

Let $M = \langle \Sigma, \square, Q, q_0, q_\mathrm{Y}, \delta \rangle$ be the fixed DTM and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{1, \dots, p(n)\}$.

Goal

$\mathrm{state}_{q_\mathrm{Y}}$

PSPACE-completeness for STRIPS plan existence

Theorem (PSPACE-completeness; Bylander, 1994)

PlanEx *and* PlanLen *are* PSPACE*-complete.*
This is true even when restricting to STRIPS tasks.

Proof.

Membership for PlanLen was already shown.

Hardness for PlanEx follows because we just presented a polynomial reduction from an arbitrary problem in PSPACE to PlanEx. (Note that the reduction only generates STRIPS tasks.)

Membership for PlanEx and hardness for PlanLen follows from the polynomial reduction from PlanEx to PlanLen. □

More complexity results

In addition to the basic complexity result presented in this chapter, there are many special cases, generalizations, variations and related problems studied in the literature:

- different planning formalisms
  - e. g., finite-domain representation, nondeterministic effects, partial observability, schematic operators, numerical state variables
- syntactic restrictions of planning tasks
  - e. g., without preconditions, without conjunctive effects, STRIPS without delete effects
- semantic restrictions of planning task
  - e. g., restricting to certain classes of causal graphs
- particular planning domains
  - e. g., Blocksworld, Logistics, FreeCell

Complexity results for different planning formalisms

Some results for different planning formalisms:

- FDR tasks:
  - same complexity as for propositional tasks ("folklore")
  - also true for the SAS^+ special case
- nondeterministic effects:
  - fully observable: EXP-complete (Littman, 1997)
  - unobservable: EXPSPACE-complete (Haslum & Jonsson, 1999)
  - partially observable: 2EXP-complete (Rintanen, 2004)
- schematic operators:
  - usually adds one exponential level to PlanEx complexity
  - e. g., classical case EXPSPACE-complete (Erol et al., 1995)
- numerical state variables:
  - undecidable in most variations (Helmert, 2002)