Semantic Networks and Description Logics

Description Logics – Terminology and Notation

Knowledge Representation and Reasoning

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Description Logics – Terminology and Notation

Motivation

- Main problem with semantic networks and frames
- The lack of formal semantics!
- Disadvantage of simple inheritance networks
- Concepts are atomic and do not have any structure

$\rightsquigarrow$ Brachman's **structural inheritance networks** (1977)

Structural Inheritance Networks

- Concepts are *defined/described* using a small set of well-defined operators
- Distinction between *conceptual* and *object-related* knowledge
- Computation of *subconcept relation* and of *instance relation*
- *Strict inheritance* (of the entire structure of a concept)

Systems and Applications

- **Systems**:
  - KL-ONE: First implementation of the ideas (1978)
  - . . . then NIKL, KL-TWO, KRYPTON, KANDOR, CLASSIC, BACK, KRIS, YAK, CRACK . . .
  - . . . currently **FaCT**, **DLP**, **RACER** 1998
- **Applications**:
  - First, natural language understanding systems
  - . . . then configuration systems,
  - . . . information systems,
  - . . . currently, it is one tool for the *semantic web*
  - **DAML+OIL**, now **OWL**

Description Logics

- Previously also *KL-ONE-alike languages*, *frame-based languages*, *terminological logics*, *concept languages*
- **Description Logics (DL)** allow us
  - to describe concepts using *complex descriptions*,
  - to introduce the terminology of an application and to structure it (**TBox**),
  - to introduce objects (**ABox**) and relate them to the introduced terminology,
  - and to *reason* about the terminology and the objects.

Informal Example

Male is:	the opposite of female
A human is a kind of:	living entity
A woman is:	a human and a female
A man is:	a human and a male
A mother is:	a woman with at least one child that is a human
A father is:	a man with at least one child that is a human
A parent is:	a mother or a father
A grandmother is:	a woman, with at least one child that is a parent
A mother-wod is:	a mother with only male children

Elizabeth is a woman
Elizabeth has the child Charles
Charles is a man
Diana is a mother-wod
Diana has the child William

Possible Questions:
Is a grandmother a parent?
Is Diana a parent?
Is William a man?
Is Elizabeth a mother-wod?

Atomic Concepts and Roles

- **Concept names**:
  - E.g., `Grandmother`, `Male`, ... (in the following usually *capitalized*)
  - We will use **symbols** such as $A, A_1, \ldots$
  - **Semantics**: Monadic predicates $A(\cdot)$ or set-theoretically a subset of the universe $A^I \subseteq \mathcal{D}$.
- **Role names**:
  - In our example, e.g., `child`. Often we will use names such as `has-child` or something similar (in the following usually *lowercase*).
  - Role names are *disjoint* from concept names
  - **Symbolically**: $t, t_1, \ldots$
  - **Semantics**: Dyadic predicates $t(\cdot,\cdot)$ or set-theoretically $t^I \subseteq \mathcal{D} \times \mathcal{D}$.

Concept and Role Description

- Out of *concept* and *role names*, complex **descriptions** can be created
- In our example, e.g. "a `Human` and `Female`."
- **Symbolically**: C for concept descriptions and r for role descriptions
- Which particular constructs are available depends on the chosen description logic
- **Predicate logic semantics**: A concept descriptions C corresponds to a formula $C(x)$ with the free variable x. Similarly with r: It corresponds to formula $r(x,y)$ with free variables x,y.
- **Set semantics**:

$$
\begin{aligned}
C^{I} &= \{d \,|\, C(d) \text{ "is true in" } I\} \\
r^{I} &= \{(d,e) \,|\, r(d,e) \text{ "is true in" } I\}
\end{aligned}
$$

Boolean Operators

- **Syntax**: let C and D be concept descriptions, then the following are also concept descriptions:
  - $C \sqcap D$ (**Concept conjunction**)
  - $C \sqcup D$ (**Concept disjunction**)
  - $\neg C$ (**Concept negation**)
- **Examples**:
  - $\texttt{Human} \sqcap \texttt{Female}$
  - $\texttt{Father} \sqcup \texttt{Mother}$
  - $\neg\, \texttt{Female}$
- **Predicate logic semantics**: $C(x) \wedge D(x)$, $C(x) \vee D(x)$, $\neg C(x)$
- **Set semantics**: $C^{I} \cap D^{I}$, $C^{I} \cup D^{I}$, $\mathcal{D} - C^{I}$

Role Restrictions

- **Motivation**:
  - Often we want to describe something by *restricting* the possible “fillers” of a role, e.g. `Mother-wod`.
  - Sometimes we want to say that there is at least a filler of a particular type, e.g. `Grandmother`
- **Idea**: Use quantifiers that range over the role-fillers
  - `Mother` $\sqcap \forall$ `has-child.Man`
  - `Woman` $\sqcap \exists$ `has-child.Parent`
- **Predicate logic semantics**:

$$(\exists r.C)(x) = \exists y : (r(x,y) \wedge C(y))$$
$$(\forall r.C)(x) = \forall y : (r(x,y) \rightarrow C(y))$$

Set semantics:

$$(\exists r.C)^I = \{d | \exists e : (d,e) \in r^I \wedge e \in C^I\}$$
$$(\forall r.C)^I = \{d | \forall e : (d,e) \in r^I \rightarrow e \in C^I\}$$

Cardinality Restriction

- **Motivation**:
  - Often we want to describe something by *restricting the number* of possible "fillers" of a role, e.g., a `Mother` with at least 3 `children` or at most 2 `children`.
- **Idea**: We restrict the cardinality of the role filler sets:
  - $\texttt{Mother} \sqcap (\geq 3\, \texttt{has-child})$
  - $\texttt{Mother} \sqcap (\leq 2\, \texttt{has-child})$
- **Predicate logic semantics**:

$$
\begin{aligned}
(\geq n\ r)(x) &= \exists y_1 \ldots y_n : \big(r(x,y_1) \wedge \ldots \wedge r(x,y_n) \wedge \\
&\qquad y_1 \neq y_2 \wedge \ldots \wedge y_{n-1} \neq y_n\big) \\
(\leq n\ r)(x) &= \neg(\geq n+1\ r)(x)
\end{aligned}
$$

- **Set semantics**:

$$
\begin{aligned}
(\geq n\ r)^{\mathcal{I}} &= \{d \mid |\{e | r^{\mathcal{I}}(d,e)\}| \geq n\} \\
(\leq n\ r)^{\mathcal{I}} &= \mathcal{D} - (\geq n+1\ r)^{\mathcal{I}}
\end{aligned}
$$

Inverse Roles

- **Motivation**:
  - How can we describe the concept "*children of rich parents*"?
- **Idea**: Define the "inverse" role for a given role (the converse relation)
  - $\texttt{has-child}^{-1}$
- **Application**: $\exists\texttt{has-child}^{-1}.\texttt{Rich}$
- **Predicate logic semantics**:

$$r^{-1}(x,y) \quad = \quad r(y,x)$$

- **Set semantics**:

$$(r^{-1})^{I} \quad = \quad \{(d,e) \mid (e,d) \in r^{I}\}$$

Role Composition

- **Motivation**:
  - How can we define the role `has-grandchild` given the role `has-child`?
- **Idea**: Compose roles (as one can compose binary relations)
  - `has-child` $\circ$ `has-child`
- **Predicate logic semantics**:

$$(r \circ s)(x,y) \quad = \quad \exists z : (r(x,z) \wedge s(z,y))$$

- **Set semantics**:

$$(r \circ s)^I \quad = \quad \{(d,e) \,|\, \exists f : (d,f) \in r^I \wedge (f,e) \in s^I\}$$

Role Value Maps

- **Motivation**:
  - How do we express the concept "*women who know all the friends of their children*"
- **Idea**: Relate role filler sets to each other
  - `Woman` $\sqcap$ `(has-child` $\circ$ `has-friend` $\sqsubseteq$ `knows)`
- **Predicate logic semantics**:

$$(r \sqsubseteq s)(x) \quad = \quad \forall y : \big(r(x,y) \to s(x,y)\big)$$

- **Set semantics**: Let $r^I(d) = \{e \mid r^I(d,e)\}$.

$$(r \sqsubseteq s)^I \quad = \quad \{d | r^I(d) \subseteq s^I(d)\}$$

- **Note**: Role value maps lead to undecidability of satisfiability of concept descriptions!

Terminology Box

- In order to *introduce* new terms, we use two kinds of **terminological axioms**:
  - $A \doteq C$
  - $A \sqsubseteq C$

  where A is a *concept name* and C is a *concept description*.
- A **terminology** or **TBox** is a finite set of such axioms with the following additional restrictions:
  - no multiple definitions of the same symbol such as $A \doteq C$, $A \sqsubseteq D$
  - no cyclic definitions (even not indirectly), such as $A \doteq \forall r.B$, $B \doteq \exists s.A$

TBoxes: Semantics

- TBoxes restrict the set of possible interpretations.
- **Predicate logic semantics**:
  - $A \doteq C$ corresponds to $\forall x : \big(A(x) \leftrightarrow C(x)\big)$
  - $A \sqsubseteq C$ corresponds to $\forall x : \big(A(x) \rightarrow C(x)\big)$
- **Set semantics**:
  - $A \doteq C$ corresponds to $A^I = C^I$
  - $A \sqsubseteq C$ corresponds to $A^I \subseteq C^I$
- Non-empty interpretations which satisfy all terminological axioms are called **models** of the TBox.

Assertional Box

- In order to state something about objects in the world, we use two forms of **assertions**:
  - $a : C$
  - $(a, b) : r$

  where a and b are **individual names** (e.g., ELIZABETH, PHILIP), C is a concept description, and r is a role description.
- An **ABox** is a finite set of assertions.

ABoxes: Semantics

- **Individual names** are interpreted as elements of the universe under the **unique-name-assumption**, i.e., different names refer to different objects.
- **Assertions** express that an object is an instance of a concept or that two objects are related by a role.
- **Predicate logic semantics**:
  - $a : C$ corresponds to $C(a)$
  - $(a,b) : r$ corresponds to $r(a,b)$
- **Set semantics**:
  - $a^I \in D$
  - $a : C$ corresponds to $a^I \in C^I$
  - $(a,b) : r$ corresponds to $(a^I, b^I) \in r^I$
- **Models** of an ABox and of ABox$+$TBox can be defined analogously to models of a TBox.

Example TBox

$$
\begin{aligned}
\texttt{Male} &\doteq \neg\texttt{Female} \\
\texttt{Human} &\sqsubseteq \texttt{Living\_entity} \\
\texttt{Woman} &\doteq \texttt{Human} \sqcap \texttt{Female} \\
\texttt{Man} &\doteq \texttt{Human} \sqcap \texttt{Male} \\
\texttt{Mother} &\doteq \texttt{Woman} \sqcap \exists\texttt{has-child.Human} \\
\texttt{Father} &\doteq \texttt{Man} \sqcap \exists\texttt{has-child.Human} \\
\texttt{Parent} &\doteq \texttt{Father} \sqcup \texttt{Mother} \\
\texttt{Grandmother} &\doteq \texttt{Woman} \sqcap \exists\texttt{has-child.Parent} \\
\texttt{Mother-without-daughter} &\doteq \texttt{Mother} \sqcap \forall\texttt{has-child.Male} \\
\texttt{Mother-with-many-children} &\doteq \texttt{Mother} \sqcap (\geq 3\,\texttt{has-child})
\end{aligned}
$$

Example ABox

```
CHARLES:   Man                              DIANA: Woman
EDWARD:    Man                              ELIZABETH: Woman
ANDREW:    Man
DIANA:     Mother-without-daughter
(ELIZABETH, CHARLES):   has-child
(ELIZABETH, EDWARD):    has-child
(ELIZABETH, ANDREW):    has-child
(DIANA, WILLIAM):       has-child
(CHARLES, WILLIAM):     has-child
```

Some Reasoning Services

- Does a description C make sense at all, i.e., is it **satisfiable**?
- A concept description C is satisfiable iff there exists an interpretation I such that $C^I \neq \emptyset$.
- Is one concept a specialization of another one, is it **subsumed**?
- C is **subsumed by** D, in symbols $C \sqsubseteq D$ iff we have for all interpretations $C^I \subseteq D^I$.
- Is a an **instance** of a concept C?
- a is an instance of C iff for all interpretations, we have $a^I \in C^I$.
- **Note**: These questions can be posed with or without a TBox that restricts the possible interpretations.

Outlook

- Can we **reduce** the reasoning services to perhaps just one problem?
- What could be **reasoning algorithms**?
- What about **complexity** and **decidability**?
- What has all that to do with **modal logics**?
- How can one build **efficient systems**?

Literature

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Summary: Concept Descriptions

Abstract	Concrete	Interpretation
A	A	A^I
$C \sqcap D$	(and C D)	$C^I \cap D^I$
$C \sqcup D$	(or C D)	$C^I \cup D^I$
$\neg C$	(not C)	$\mathcal{D} - C^I$
$\forall r.C$	(all r C)	$\{d \in \mathcal{D} \mid r^I(d) \subseteq C^I\}$
$\exists r$	(some r)	$\{d \in \mathcal{D} \mid r^I(d) \neq \emptyset\}$
$\geq n\ r$	(atleast n r)	$\{d \in \mathcal{D} \mid \lvert r^I(d)\rvert \geq n\}$
$\leq n\ r$	(atmost n r)	$\{d \in \mathcal{D} \mid \lvert r^I(d)\rvert \leq n\}$
$\exists r.C$	(some r C)	$\{d \in \mathcal{D} \mid r^I(d) \cap C^I \neq \emptyset\}$
$\geq n\ r.C$	(atleast n r C)	$\{d \in \mathcal{D} \mid \lvert r^I(d) \cap C^I\rvert \geq n\}$
$\leq n\ r.C$	(atmost n r C)	$\{d \in \mathcal{D} \mid \lvert r^I(d) \cap C^I\rvert \leq n\}$
$r \doteq s$	(eq r s)	$\{d \in \mathcal{D} \mid r^I(d) = s^I(d)\}$
$r \neq s$	(neq r s)	$\{d \in \mathcal{D} \mid r^I(d) \neq s^I(d)\}$
$r \sqsubseteq s$	(subset r s)	$\{d \in \mathcal{D} \mid r^I(d) \subseteq s^I(d)\}$
$g \doteq h$	(eq g h)	$\{d \in \mathcal{D} \mid g^I(d) = h^I(d) \neq \emptyset\}$
$g \neq h$	(neq g h)	$\{d \in \mathcal{D} \mid \emptyset \neq g^I(d) \neq h^I(d) \neq \emptyset\}$
$\{i_1, i_2, \ldots, i_n\}$	(oneof $i_1 \ldots i_n$)	$\{i_1^I, i_2^I, \ldots, i_n^I\}$

Summary: Role Descriptions

Abstract	Concrete	Interpretation
t	t	$t^{\mathcal{I}}$
f	f	$f^{\mathcal{I}}$, *(functional role)*
$r \sqcap s$	(and r s)	$r^{\mathcal{I}} \cap s^{\mathcal{I}}$
$r \sqcup s$	(or r s)	$r^{\mathcal{I}} \cup s^{\mathcal{I}}$
$\neg r$	(not r)	$\mathcal{D} \times \mathcal{D} - r^{\mathcal{I}}$
r^{-1}	(inverse r)	$\{(d,d') \mid (d',d) \in r^{\mathcal{I}}\}$
$r\|_C$	(restr r C)	$\{(d,d') \in r^{\mathcal{I}} \mid d' \in C^{\mathcal{I}}\}$
r^+	(trans r)	$(r^{\mathcal{I}})^+$
$r \circ s$	(compose r s)	$r^{\mathcal{I}} \circ s^{\mathcal{I}}$
$\mathbf{1}$	self	$\{(d,d) \mid d \in \mathcal{D}\}$