Dynamic Epistemic Logic 7. Multi-Agent PathFInding with Destination Uncertainty

IBURG

ΞW

Albert-Ludwigs-Universität Freiburg

Bernhard Nebel and Robert Mattmüller July 10th, 2019





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Credits:

This presentation is based on our JAIR paper: "Implicitly Coordinated Multi-Agent Path Finding under Destination Uncertainty: Success Guarantees and Computational Complexity"



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Motivation

- All the examples of epistemic planning we have seen are rather simple.
- $\rightarrow\,$ Are there more complex settings, which could be dealt with in the framework?
 - There are no success guarantees for the execution of joint plans in the general case.
- → Are there possibly specialized epistemic planning situations, for which postive results are possible?
 - Propositional epistemic planning in general is undecidable (something, we have not seen yet).
- \rightarrow Are there decidable special cases?
- \Rightarrow MAPF/DU originally only used as a motivating example.

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(Classical) Multi-Agent Path Finding

Definition (Multi-agent path finding (MAPF) problem)

Given a set of *agents A*, an undirected, simple *graph* G = (V, E), an *initial state* modelled by an injective function $\alpha_0 : A \to V$, and a *goal state* modelled by another injective function α_* , can α_0 be *transformed* into α_* by *movements of single agents* without collisions?

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Existence problem: Does there exist a successful sequence of movements (= plan)?

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- Existence problem: Does there exist a successful sequence of movements (= plan)?
- Bounded existence problem: Does there exist a plan of a given length k or less?

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- Plan generation problem: Generate a plan.

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- Existence problem: Does there exist a successful sequence of movements (= plan)?
- Bounded existence problem: Does there exist a plan of a given length k or less?
- Plan generation problem: Generate a plan.
- Optimal plan generation problem: Generate a shortest plan.

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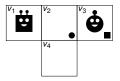
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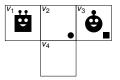
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$$G = (V, E) \text{ with } V = \{v_1, v_2, v_3, v_4\} \text{ and } E = \left\{\{v_1, v_2\}, \{v_2, v_3\}, \{v_2, v_4\}\right\}$$
$$A = \{S, C\} \text{ and } \alpha_0(S) = v_1, \alpha_0(C) = v_3, \alpha_*(S) = v_3, \alpha_*(C) = v_2$$

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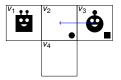
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Plan:

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Example

Can we find a (central) plan to move the square robot S to v_3

and the circle robot C to v_2 ?

 $G = (V, E) \text{ with } V = \{v_1, v_2, v_3, v_4\} \text{ and } E = \{\{v_1, v_2\}, \{v_2, v_3\}, \{v_2, v_4\}\}$ $A = \{S, C\}$ and $\alpha_0(S) = v_1, \alpha_0(C) = v_3, \alpha_*(S) = v_3, \alpha_*(C) = v_2$

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Plan: (C, v_3, v_2) ,

Can we find a (central) plan to move the square robot *S* to v_3 and the circle robot *C* to v_2 ?

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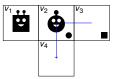
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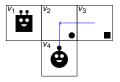






Example

Can we find a (central) plan to move the square robot *S* to v_3 and the circle robot *C* to v_2 ?



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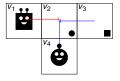
Literature

Plan: (C, v_3, v_2) , (C, v_2, v_4) ,



Example

Can we find a (central) plan to move the square robot *S* to v_3 and the circle robot *C* to v_2 ?



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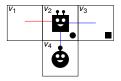
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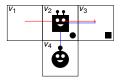
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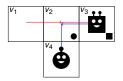
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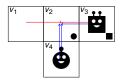
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Plan: (C, v_3, v_2) , (C, v_2, v_4) , (S, v_1, v_2) , (S, v_2, v_3) ,

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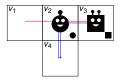
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 $\mathsf{Plan:}\;(C,v_3,v_2),\,(C,v_2,v_4),\,(S,v_1,v_2),\,(S,v_2,v_3),\!(C,v_4,v_2).$

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A special case: 15-puzzle



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1	2	3	4
12	13	14	5
11		15	6
10	9	8	7

Pictures from Wikipedia article on 15-Puzzle

A special case: 15-puzzle





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Pictures from Wikipedia article on 15-Puzzle

2

6

10

3

7

11

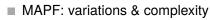
15

Δ

8

12





- Distributed MAPF (each agent plans on it own): DMAPF
- Distributed MAPF with destination uncertainty: MAPF/DU

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Sequential MAPF

- Sequential MAPF (or pebble motion on a graph) allows only one agent to move per time step.
- An agent $a \in A$ can move in one step from $s \in V$ to $t \in V$ transforming α to α' , if

$$\square \alpha(a) = s,$$

$$\langle \boldsymbol{s},t
angle \in \boldsymbol{E},$$

- there is no agent *b* such that $\alpha(b) = t$.
- In this case, α' is determined as follows:
 - $\blacksquare \alpha'(a) = t,$
 - for all agents $b \neq a : \alpha(b) = \alpha'(b)$,
- One usually wants to minimize the number of single movements (= sum-of-cost over all agents)

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- Parallel MAPF allows many agents to move in parallel, provided they do not collide.
- Two models:

Parallel MAPF

- Parallel: A chain of agents can move provided the first agent can move on an unoccupied vertex.
- Parallel with rotations: A closed cycle can move synchronously.
- In both cases, one is usually interested in the number of parallel steps (= make-span).
- However, also the sum-of-cost is sometimes considered.



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- There is a set of agents and a set of targets (of the same cardinality as the agent set).
- Each target must be reached by one agent.
- This means one first has to assign a target and then to solve the original MAPF problem.
- Interestingly, the problem as a whole is easier to solve (using flow-based techniques).

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Existence: For arbitrary graphs with at least one empty place, the problem is polynomial (O(|V|³) using Kornhauser's algorithm). For BIBOX on bi-connected with at least two empty places also cubic, but smaller constant.

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- Generation: O(|V|³), generating the same number of steps, again using Kornhauser's algorithm or BIBOX

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- Bounded existence: Is definitely in NP
 - If there exists a solution, then it is polynomially bounded.
 - A solution candidate can be checked in polynomial time for satisfying the conditions of being a movement plan with k of steps or less.
- Question: Is the problem also NP-hard?

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Definition (Exact Cover By 3-Sets (X3C) Problem)

Given a set of elements *U* and a collection of subsets $C = \{s_j\}$ with $s_j \subseteq U$ and $|s_j| = 3$. Is there a sub-collection of subsets $C' \subseteq C$ such that $\bigcup_{s \in C'} s = U$ and all subsets in C' are pairwise disjoint, i.e., $s_a \cap s_b = \emptyset$ for each $s_a, s_b \in C'$ with $s_a \neq s_b$?

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X3C is NP-complete.



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X3C is NP-complete.

Example

 $U = \{1, 2, 3, 4, 5, 6\}$

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Given a set of elements *U* and a collection of subsets $C = \{s_j\}$ with $s_j \subseteq U$ and $|s_j| = 3$. Is there a sub-collection of subsets $C' \subseteq C$ such that $\bigcup_{s \in C'} s = U$ and all subsets in C' are pairwise disjoint, i.e., $s_a \cap s_b = \emptyset$ for each $s_a, s_b \in C'$ with $s_a \neq s_b$?

X3C is NP-complete.

Example

 $U = \{1,2,3,4,5,6\}$ $C = \{\{1,2,3\},\{2,3,4\},\{2,5,6\},\{1,5,6\}\}$ UNI FREIBURG

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X3C is NP-complete.

Example

$$U = \{1,2,3,4,5,6\}$$

$$C = \{\{1,2,3\},\{2,3,4\},\{2,5,6\},\{1,5,6\}\}$$

$$C'_1 = \{\{1,2,3\},\{2,3,4\}\}$$
 is not a cover.

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X3C is NP-complete.

Example

$$U = \{1,2,3,4,5,6\}$$

$$C = \{\{1,2,3\},\{2,3,4\},\{2,5,6\},\{1,5,6\}\}$$

$$C'_1 = \{\{1,2,3\},\{2,3,4\}\} \text{ is not a cover.}$$

$$C'_2 = \{\{1,2,3\},\{2,3,4\},\{1,5,6\}\} \text{ is not an exact cover.}$$

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X3C is NP-complete.

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$$U = \{1,2,3,4,5,6\}$$

$$C = \{\{1,2,3\},\{2,3,4\},\{2,5,6\},\{1,5,6\}\}$$

$$C'_1 = \{\{1,2,3\},\{2,3,4\}\} \text{ is not a cover.}$$

$$C'_2 = \{\{1,2,3\},\{2,3,4\},\{1,5,6\}\} \text{ is not an exact cover.}$$

$$C'_3 = \{\{2,3,4\},\{1,5,6\}\} \text{ is an exact cover.}$$

Motivation

BUR

Multi-Agent Pathfinding

Definition and example

Computational Complextiy of MAPE

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

NP-hardness of MAPF: Reduction from X3C

$$C = \{\{1,2,3\},\{2,3,4\},\{2,5,6\},\{1,5,6\}\}$$



Motivation

Multi-Agent Pathfinding

Definition and example MAPE Variations

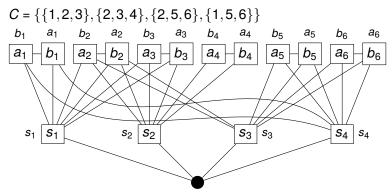
Computational Complextiy of MAPF

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MAPF/DU

Complexity

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Motivation

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Multi-Agent Pathfinding

Definition and example

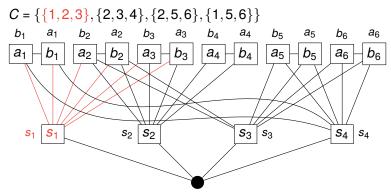
Computational Complextiy of MAPF

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Multi-Agent Pathfinding

Definition and example

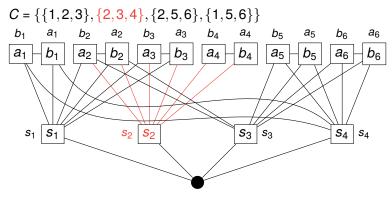
Computational Complextiy of MAPF

Going Beyond

MAPF/DU

Complexity

Summary & Outlook



Motivation

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Multi-Agent Pathfinding

Definition and example

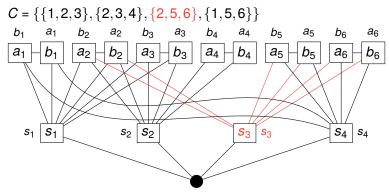
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Summary & Outlook



Motivation

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Multi-Agent Pathfinding

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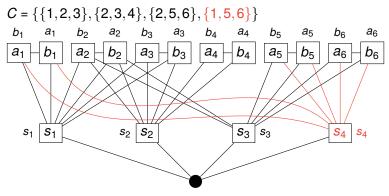
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Complexity

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Multi-Agent Pathfinding

Definition and example

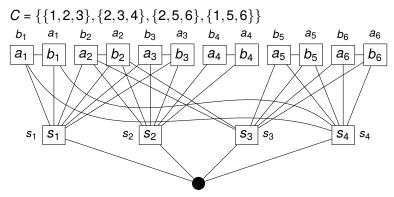
Computational Complextiy of MAPE

Going Beyond

MAPF/DU

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Summary & Outlook



Claim: There is an exact cover by 3-sets iff the constructed MAPF instance can be solved in at most k = 11/3|U| moves.

Motivation

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Definition and example

Computational Complextiy of MAPF

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MAPF/DU

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Summary & Outlook

Literature

Going Beyond: DMAPF and MAPF/DU

The common goal of all agents is still that everybody reaches its destination.

Motivation

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Multi-Agent Pathfinding

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MAPF/DU

Complexity

Summary & Outlook

- The common goal of all agents is still that everybody reaches its destination.
- Distributed MAPF (DMAPF):

Motivation

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Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

- The common goal of all agents is still that everybody reaches its destination.
- Distributed MAPF (DMAPF):
 - All agents plan and re-plan without communicating with their peers.

Motivation

2

Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

- The common goal of all agents is still that everybody reaches its destination.
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 - All agents plan and re-plan without communicating with their peers.
 - ightarrow Models multi-robot interactions without communication

Motivation

JRG

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Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

- The common goal of all agents is still that everybody reaches its destination.
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- MAPF under destination uncertainty (MAPF/DU):

Motivation

2

Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

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Motivation

> Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

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Motivation

> Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

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 - For each agent, there exists a set of possible destinations, which are common knowledge.
 - $\rightarrow\,$ Models multi-robot interactions without communication and with goal uncertainty

Motivation

2

Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook



- Note: Full observability = uniform knowledge
- General results apply: Joint execution successful if agents are optimally eager

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Motivation

Multi-Agent Pathfinding

Going Beyond

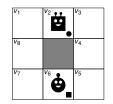
MAPF/DU

Complexity

Summary & Outlook

- Note: Full observability = uniform knowledge
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 - are optimally eager

S initially: $\langle (S, v_2, v_3), \ldots \rangle$ C initially: $\langle (C, v_6, v_5), \ldots \rangle$





Motivation

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S initially: $\langle (S, v_2, v_3), \ldots \rangle$ *C* initially: $\langle (C, v_6, v_5), \ldots \rangle$

are optimally eager

Note: Full observability = uniform knowledge

V1

 V_7

 V_2

 V_6

General results apply: Joint execution successful if agents

••

 V_5

- Summary & Outlook
- Beyond MAPE/DU

Motivation

Multi-Agent

Pathfinding Going

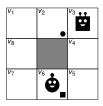
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- *C* initially: $\langle (C, v_6, v_5), \ldots \rangle$
 - C after (S, v_2, v_3) : $\langle (C, v_6, v_7), \ldots \rangle$



18/53

- Note: Full observability = uniform knowledge
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Motivation

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Summary & Outlook

18/53

Note: Full observability = uniform knowledge

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- *S* initially: $\langle (S, v_2, v_3), \ldots \rangle$
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- C after (S, v_2, v_3) : $\langle (C, v_6, v_7), \ldots \rangle$

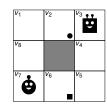
Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook





Note that solving MAPF optimally implies NP-hardness, while in general (sub-optimal) MAPF is polynomial. Motivation

JRG

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Multi-Agent Pathfinding

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MAPF/DU

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Summary & Outlook

- Note that solving MAPF optimally implies NP-hardness, while in general (sub-optimal) MAPF is polynomial.
- An alternative to avoid infinite executions is to plan conservatively: When replanning, consider the already executed part as a prefix of the plan from the original configuration (history dependent policy!).

Motivation

Multi-Agent Pathfinding

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MAPF/DU

Complexity

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Motivation

Multi-Agent Pathfinding

Going Beyond

MAPF/DU

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Summary & Outlook

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- \rightarrow Because plans should be cycle-free, no agent will ever consider to revisit a previously visited configuration!
 - We do not have to solve NP-hard problems every time we (re-)plan!

Motivation

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Multi-Agent Pathfinding

Going Beyond

MAPF/DU

Complexity

Summary & Outlook

In the worst-case (if scheduled unfavorably), the agents might consider all of the exponentially many configurations



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ZW

Multi-Agent Pathfinding

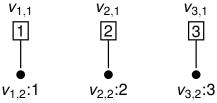
Going Beyond

MAPF/DU

Complexity

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Motivation

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Multi-Agent Pathfinding

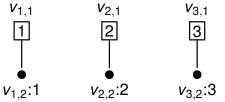
Going Beyond

MAPF/DU

Complexity

Summary & Outlook

In the worst-case (if scheduled unfavorably), the agents might consider all of the exponentially many configurations



Assuming here a schedule similar to a Gray counter (a counter that changes only one bit at a time), we visit all possible configurations.

Motivation

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Multi-Agent Pathfinding

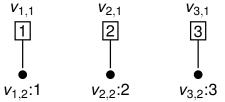
Going Beyond

MAPF/DU

Complexity

Summary & Outlook

In the worst-case (if scheduled unfavorably), the agents might consider all of the exponentially many configurations



- Assuming here a schedule similar to a Gray counter (a counter that changes only one bit at a time), we visit all possible configurations.
- In addition, the rule-based polynomial planing methods may not create the right plan prefixes, except when all agents use the identical algorithm.

Motivation

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Multi-Agent Pathfinding

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Multi-Agent Pathfinding

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MAPF/DU

Implicitly Coordinated Branching Plans Strong plans Stepping Stones Execution cost Execution

Complexity

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Literature

MAPF/DU

MAPF/DU representation & state space

■ In addition to the sets of agents *A*, the graph G = (V, E), and the assignment of agents to nodes α , we need a function to represent the *possible destinations* $\beta : A \rightarrow 2^V$.

Motivation

DRD

2

Multi-Agent Pathfinding

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Complexity

Summary & Outlook

- In addition to the sets of agents *A*, the graph G = (V, E), and the assignment of agents to nodes α , we need a function to represent the *possible destinations* $\beta : A \rightarrow 2^V$.
- We assume that the set of possible destinations are pairwise disjoint (this can be relaxed, though).

Motivation

D^RC

2

Multi-Agent Pathfinding

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- An *objective state* is given by the pair $s = \langle \alpha, \beta \rangle$ representing the common knowledge of all agents.

Motivation

D^RC

2

Multi-Agent Pathfinding

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- A *subjective state* of agent *i* is given by $s^i = \langle \alpha, \beta, i, v \rangle$ with $v \in \beta(i)$ or $v = \bot$ if $\beta(i) = \emptyset$, representing the private knowledge of agent *i*.

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Summarv &

Outlook

- In addition to the sets of agents *A*, the graph G = (V, E), and the assignment of agents to nodes α , we need a function to represent the *possible destinations* $\beta : A \rightarrow 2^V$.
- We assume that the set of possible destinations are pairwise disjoint (this can be relaxed, though).
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- A *MAPF/DU instance* is given by $\langle A, G, s_0, \alpha_* \rangle$, where $s_0 = \langle \alpha_0, \beta_0 \rangle$.

Motivation

2

Multi-Agent Pathfinding

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Complexity

Summary & Outlook

Goal specification is complicated since all agents should reach their true destination. Motivatio

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Complexity

Summary & Outlook

- Goal specification is complicated since all agents should reach their true destination.
- \rightarrow Can be specified using a conjunction of implications: If *x* is *i*'s destination, then try to achieve a configuration where *i* is on *x*.

Motivation

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Multi-Agent Pathfinding

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Complexity

Summary & Outlook

- Goal specification is complicated since all agents should reach their true destination.
- \rightarrow Can be specified using a conjunction of implications: If *x* is *i*'s destination, then try to achieve a configuration where *i* is on *x*.
 - In order to form an implicitly coordinated plan, the last agent moving needs to know that all other agents have reached their destinations.

Motivation

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Complexity

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 - In order to form an implicitly coordinated plan, the last agent moving needs to know that all other agents have reached their destinations.
- → Use a public announcement when destination has been reached so that true destination becomes common knolwedge. We require that the agent does not move afterwards, but this can be varied.

Motivation

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- → Use a public announcement when destination has been reached so that true destination becomes common knolwedge. We require that the agent does not move afterwards, but this can be varied.
 - ightarrow Common goal can be stated as common knowledge goal.

Motivation

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Summary & Outlook

Simplifying properties of MAPF/DU

- Since there are no private announcements and no non-deterministic effects, the update operation is simple.
- $\rightarrow\,$ Simple ontic or public announcement update, never adding any worlds.
 - Perspective-shifts are also simple.
- → If shifting perspective to agent *i*, simply assume all combinations of destinations from the agents as possible.
 Perhaps branch over *i*'s possible destinations.

Motivation

Multi-Agent Pathfinding

Going Beyond

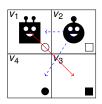
MAPF/DU

Implicitly Coordinated Branching Plans Strong plans Stepping Stones Execution cost Execution guarantees

Complexity

Summary & Outlook

Square agent *S* wants to go to v_3 and knows that circle agent *C* wants to go to v_1 or v_4 .



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Strong plans

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Execution cost

Execution

Complexity

Summary & Outlook

- Square agent *S* wants to go to v_3 and knows that circle agent *C* wants to go to v_1 or v_4 .
- C wants to go to v₄ and knows that S wants to go to v₂ or v₃.



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Multi-Agent Pathfinding

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Strong plans

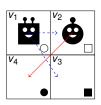
Stepping Stone

Execution cost

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Summary & Outlook



- Square agent *S* wants to go to v_3 and knows that circle agent *C* wants to go to v_1 or v_4 .
- C wants to go to v₄ and knows that S wants to go to v₂ or v₃.
- Let us assume *S* forms a plan in which it moves in order to empower *C* to reach their common goal.

Motivation

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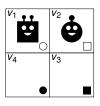
Stepping Stone

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- C wants to go to v₄ and knows that S wants to go to v₂ or v₃.
- Let us assume S forms a plan in which it moves in order to empower C to reach their common goal.
 - S needs *shifting its perspective* in order to plan for all possible destinations of *C* (*branching on destinations*).

Motivation

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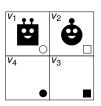
Stepping Stones

Execution cost

guarantees

Complexity

Summary & Outlook





- Square agent *S* wants to go to v_3 and knows that circle agent *C* wants to go to v_1 or v_4 .
- *C* wants to go to *v*₄ and knows that *S* wants to go to *v*₂ or *v*₃.
- Let us assume S forms a plan in which it moves in order to empower C to reach their common goal.
- S needs shifting its perspective in order to plan for all possible destinations of C (branching on destinations).
- Planning for C, S must forget about its own destination.

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Strong plans

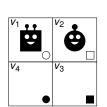
Stepping Stones

Execution cost

guarantees

Complexity

Summary & Outlook



Branching plans: Building blocks

Branching plans consist of:

- Movement actions: ((agent), (sourcenode), (targetnode)), i.e., a movement of an agent
- Success announcement: ((agent), S), after that all agents know that the agent has reached its destination and it cannot move anymore
- Perspective shift: [(agent):...], i.e., from here on we assume to plan with the knowledge of agent (agent). This can be unconditional or conditional on (agent)'s destinations.
- Branch on all destinations:

 $(?\langle dest_1 \rangle \{...\},...,?\langle dest_n \rangle \{...\})$, where all destinations of the current agent have to be listed. For each case we try to find a successful plan to reach the goal state.

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Movement actions modify α in the obvious way.

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Summary & Outlook

- Movement actions modify α in the obvious way.
- A success announcement of agent *i* transforms β to β' such that $\beta'(i) = \emptyset$ in order to signal that *i* cannot move anymore.



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Summary & Outlook

- **Movement actions** modify α in the obvious way.
- A success announcement of agent *i* transforms β to β' such that $\beta'(i) = \emptyset$ in order to signal that *i* cannot move anymore.
- A perspective shift from *i* to *j* with subsequent branching on destinations transforms the subjective state sⁱ = ⟨α, β, i, v_i⟩ to a set of subjective states s^{j_k} = ⟨α, β, j, v_{j_k}⟩ with all v_{j_k} ∈ β(j).

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Summary & Outlook

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- A perspective shift from *i* to *j* with subsequent branching on destinations transforms the subjective state sⁱ = ⟨α, β, i, v_i⟩ to a set of subjective states s^{j_k} = ⟨α, β, j, v_{j_k}⟩ with all v_{j_k} ∈ β(j).
- A perspective shift from *i* to *j* without subsequent branching on destinations induces the same transformation, but enforces that the subsequent plans are the same for all states subjective states s^{jk}.

Motivation

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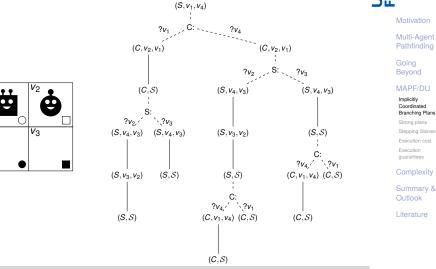
Execution cost

guarantees

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Branching plan: Example

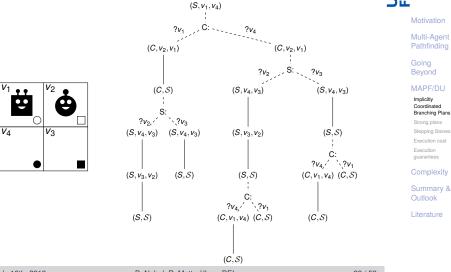


 V_1

 V_4

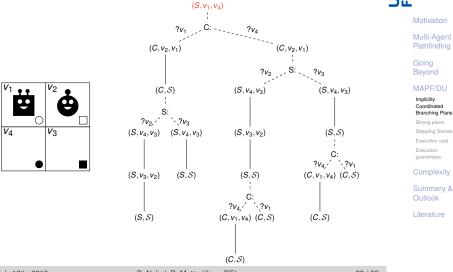
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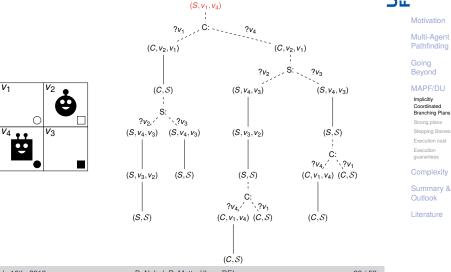
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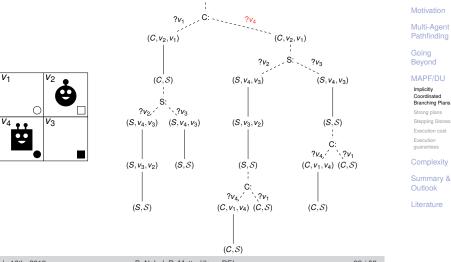
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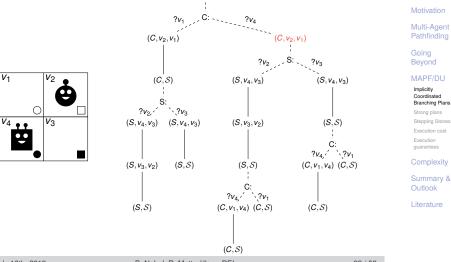


 (S, v_1, v_4)

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 (S, v_1, v_4)

 V_1

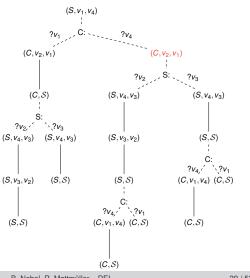
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 V_1 V_2 V4 . V₃

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Beyond

MAPF/DU

Strong plans

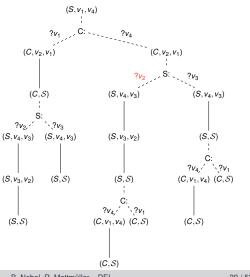
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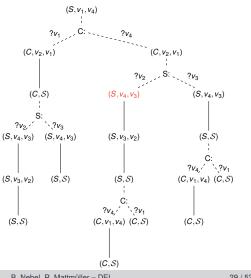
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 V_1 V_2 V4 . V₃



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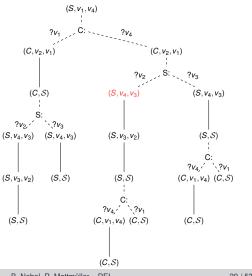
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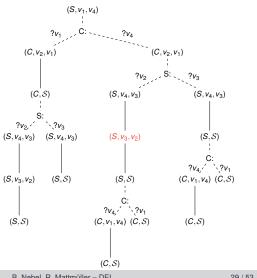
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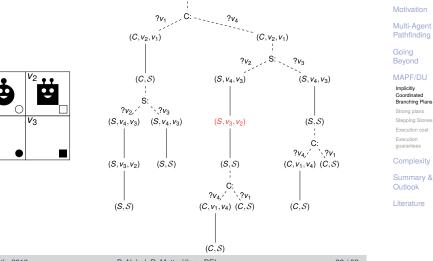
Execution

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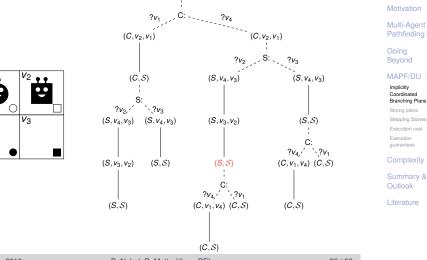
 (S, v_1, v_4)

 V_1

 V_4

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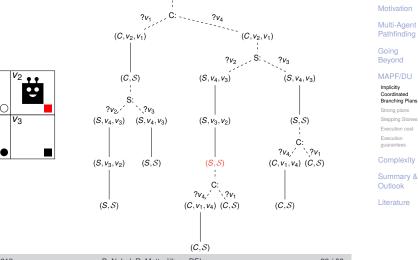
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 V_1

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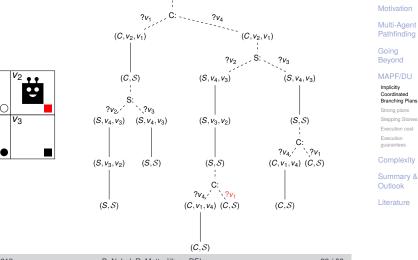
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 V_1

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 (S, v_1, v_4)

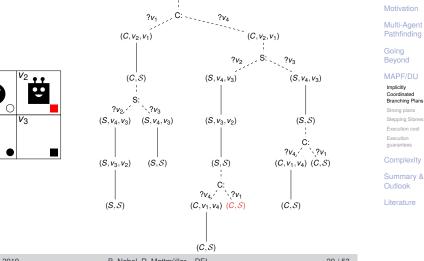
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 (S, v_1, v_4)

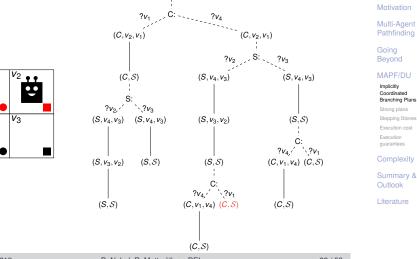
 V_1

 V_4

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Branching plan: Subjective execution example



 (S, v_1, v_4)

 V_1

 V_4

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Similar to the notion of strong plans in non-deterministic single-agent planning, we define *i-strong plans* for an agent *i* to be:

- *cycle-free*, i.e., not visiting the same objective state twice;
- always successful, i.e. always ending up in a state such that all agents have announced success;
- covering, i.e., for all combinations of possible destinations of agents different from *i*, success can be reached.

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- A plan is called *subjectively strong* if it is *i*-strong for some agent *i*.
- A plan is called *objectively strong* if it is *i*-strong for each agent *i*.
- An instance is *objectively* or *subjectively solvable* if there exists an objectively or subjectively strong plan, respectively.

Motivation

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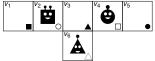
Execution cost

Execution guarantees

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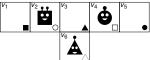
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→ There does not exist a *T*-strong plan, but an *S*- and a *C*-strong plan.



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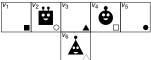
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- \rightarrow There does not exist a *T*-strong plan, but an *S* and a *C*-strong plan.
 - Difference between subjective and objective solvability concerns only the first acting agent!



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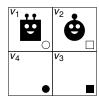
Execution cost

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Summary & Outlook

A stepping stone for agent *i* is a state in which *i* can move to each of its possible destinations, announcing success, and afterwards, for each possible destination, there exists an *i*-strong plan to solve the resulting states.



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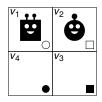
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- A stepping stone for agent *i* is a state in which *i* can move to each of its possible destinations, announcing success, and afterwards, for each possible destination, there exists an *i*-strong plan to solve the resulting states.
- S can create a stepping stone for C by moving from v_1 via v_4 to v_3 .



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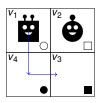
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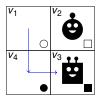
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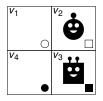
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July 10th, 2019

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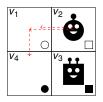
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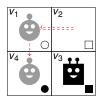
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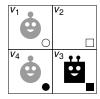
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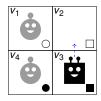
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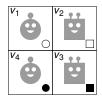
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Theorem

Given an i-solvable MAPF/DU instance, there exists an i-strong branching plan such that the only branching points are those utilizing stepping stones. UNI FREIBURG

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Theorem

Given an i-solvable MAPF/DU instance, there exists an i-strong branching plan such that the only branching points are those utilizing stepping stones.

Proof sketch.

Remove non-stepping stone branching points by picking one branch without success announcement.

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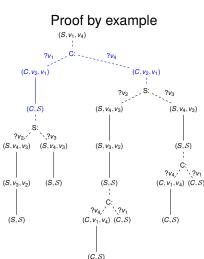
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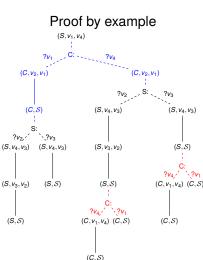
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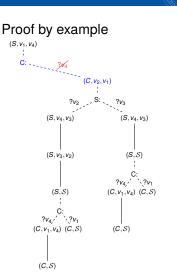
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The *execution cost* of a branching plan is the number of atomic actions of the longest execution trace.



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The *execution cost* of a branching plan is the number of atomic actions of the longest execution trace.

Theorem

Given an *i*-solvable MAPF/DU instance over a graph G = (V, E), then there exists an *i*-strong branching plan with execution cost bounded by $O(|V|^4)$.

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Summary & Outlook

The *execution cost* of a branching plan is the number of atomic actions of the longest execution trace.

Theorem

Given an *i*-solvable MAPF/DU instance over a graph G = (V, E), then there exists an *i*-strong branching plan with execution cost bounded by $O(|V|^4)$.

Proof sketch.

Direct consequence of the stepping stone theorem and the maximal number of movements in the MAPF problem.

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Joint execution is defined similarly to the fully observable case: One agent is chosen; afterwards the plan is tracked or the agent has to replan. Motivation

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Summary & Outlook

- Joint execution is defined similarly to the fully observable case: One agent is chosen; afterwards the plan is tracked or the agent has to replan.
- In the MAPF/DU framework not all agents might have a plan initially!

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IRG

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- One might hope that optimally eager agents are always successful.
- In epistemic planning this was proven to be true only in the *uniform knowledge* case.

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- Joint execution is defined similarly to the fully observable case: One agent is chosen; afterwards the plan is tracked or the agent has to replan.
- In the MAPF/DU framework not all agents might have a plan initially!
- One might hope that optimally eager agents are always successful.
- In epistemic planning this was proven to be true only in the *uniform knowledge* case.
- We do not have uniform knowledge ... and indeed, execution cycles cannot be excluded.

Motivation

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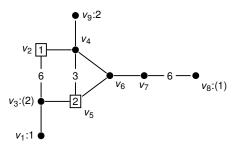
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A number on an edge means that there are as many nodes on a line.

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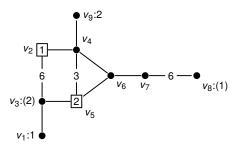
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A number on an edge means that there are as many nodes on a line.

Agent 2 has a shortest eager plan moving first to v_6 .



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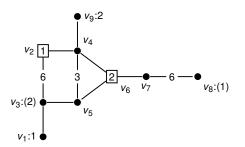
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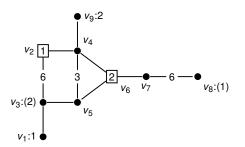
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Execution cost

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A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
- Agent 1 has then a shortest eager plan moving first to v_4 .



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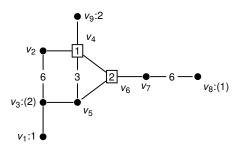
Strong plans

= ...

Execution guarantees

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Summary & Outlook



A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
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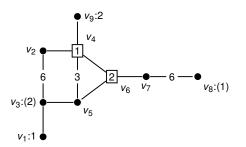
Strong plans

- -

Execution guarantees

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Summary & Outlook



A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
- Agent 1 has then a shortest eager plan moving first to v_4 .
- Agent 2 has then a shortest eager plan moving first to v_5 .

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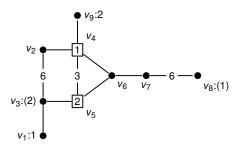
Stepping Stope

Execution cost

Execution guarantees

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Summary & Outlook



A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
- Agent 1 has then a shortest eager plan moving first to v_4 .
- Agent 2 has then a shortest eager plan moving first to v_5 .

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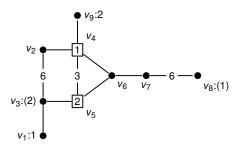
Execution cost

Execution guarantees

Complexity

Summary & Outlook

A counter example



A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
- Agent 1 has then a shortest eager plan moving first to v_4 .
- Agent 2 has then a shortest eager plan moving first to v₅.
- Agent 1 has then a shortest eager plan moving first to v₂.

Multi-Agent

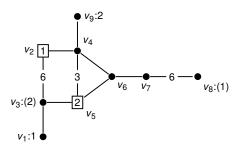
Pathfinding

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A number on an edge means that there are as many nodes on a line.

- Agent 2 has a shortest eager plan moving first to v_6 .
- Agent 1 has then a shortest eager plan moving first to v_4 .
- Agent 2 has then a shortest eager plan moving first to v₅.
- Agent 1 has then a shortest eager plan moving first to v_2 .

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MAPE/DU

Execution guarantees

Summary & Outlook

- Perhaps conservatism can help!
- Similarly to DMAPF, conservative replanning means that the already executed actions are used as a prefix in the plan to be generated.
- Differently from DMAPF, we assume that after a success announcement, the initial state is modified so that the *real destination* of the agent is known in the initial state.
- Otherwise we could not solve instances that are only subjectively solvable.

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Summary & Outlook

Conservative, eager agents are always successful, but might visit the entire state space before terminating.



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- Conservative, eager agents are always successful, but might visit the entire state space before terminating.
- Adding optimal eagerness can help to reduce the execution length.

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- Conservative, eager agents are always successful, but might visit the entire state space before terminating.
- Adding optimal eagerness can help to reduce the execution length.

Theorem

For solvable MAPF/DU instances, joint execution and replanning by conservative, optimally eager agents is always successful and the execution length is polynomial.

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Summary & Outlook

- Conservative, eager agents are always successful, but might visit the entire state space before terminating.
- Adding optimal eagerness can help to reduce the execution length.

Theorem

For solvable MAPF/DU instances, joint execution and replanning by conservative, optimally eager agents is always successful and the execution length is polynomial.

Proof idea.

After the second agent starts to act, all agents have an identical perspective and for this reason produce objectively strong plans with the same execution costs, which can be shown to be bounded polynomially using the stepping stone theorem.

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Assume S moves first to v_4 .



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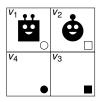
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Assume S moves first to v_4 .



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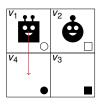
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Assume S moves first to v_4 .

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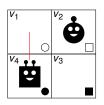
Stepping Stones

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Summary & Outlook



- Assume *S* moves first to *v*₄.
- Assume *C* re-plans. From now on, in replanning from the beginning, it has to do a perspective shift to *S*, because it now has to extend the partial plan starting with (S, v_4, v_1) , i.e., it has to create an objectively strong plan.

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- Assume that C moves now to v_1 .

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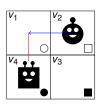
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Execution

Execution guarantees

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Summary & Outlook



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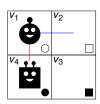
Stepping Stone:

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Summary & Outlook



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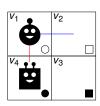
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Summary & Outlook



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- Assume that C moves now to v_1 .
- From now on, also *S* has to make a perspective shift to *C*, effectively "forgetting" its own destination, i.e., it also has to create a objectively strong plan.

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Computational Complexity of MAPF/DU planning

Computational Complexity of MAPF/DU planning

- MAPF planning is a polynomial problem.
- DMAPF is NP-hard, because we need optimal MAPF solutions.
- What about MAPF/DU?
- \rightarrow We look at the problem of deciding whether exists a plan with execution costs *k* (decision problem corresponding to the optimization problem)

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Summary & Outlook

Problems are categorized into *complexity classes* according to the requirements of computational resources:

- The class of problems decidable on deterministic Turing machines in polynomial time: P
 - Problems in P are assumed to be *efficiently solvable* (although this might not be true if the exponent is very large)
 - In practice, a reasonable definition
- The class of problems decidable on *non-deterministic Turing machines* in *polynomial time*, i.e., having a poly. length accepting computation for all positive instances: NP
- More classes are definable using other resource bounds on time and memory

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Computational Complexity: PSPACE

There are problems even more difficult than NP ...



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Computational Complexity: PSPACE

There are problems even more difficult than NP ...

Definition ((N)PSPACE)

PSPACE (NPSPACE) is the class of decision problems that can be decided on deterministic (non-deterministic) Turing machines using only *polynomially many tape cells*.

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There are problems even more difficult than NP ...

Definition ((N)PSPACE)

PSPACE (NPSPACE) is the class of decision problems that can be decided on deterministic (non-deterministic) Turing machines using only *polynomially many tape cells*.

Some facts about PSPACE:

- PSPACE is closed under complements (... as all other deterministic classes)
- PSPACE is *identical* to NPSPACE (because non-deterministic Turing machines can be simulated on deterministic TMs using only quadratic space: Savitch's Theorem)
- NPCPSPACE (because in polynomial time one can "visit" only polynomial space, i.e., NPCNPSPACE)



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Definition (PSPACE-completeness)

A decision problem (or language) is PSPACE-complete if it is in PSPACE and all other problems in PSPACE can be polynomially reduced to it.



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Computational Complexity: PSPACE-completeness

Definition (PSPACE-completeness)

A decision problem (or language) is PSPACE-complete if it is in PSPACE and all other problems in PSPACE can be polynomially reduced to it.

PSPACE-complete problems appear to be "harder" than *NP-complete* problems from a *practical point of view*.

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Computational Complexity: PSPACE-completeness

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PSPACE-complete problems appear to be "harder" than *NP-complete* problems from a *practical point of view*.

An example for a PSPACE-complete problem is **QBF**:

Instance: A Boolean formula φ with quantification over the variables in φ :

 $\forall x_1 \exists x_2 \dots \forall x_n : \varphi$

Question: Is the QBF formula true under all assignments described by the quantification?

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Theorem

Deciding whether there exists an eager MAPF/DU i-strong or objectively strong plan with execution cost k or less is PSPACE-complete.

Proof sketch.

Since plans have polynomial depth, all execution traces can be generated non-deterministically and tested using only polynomial space, i.e., PSPACE-membership.

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Example construction for

 $\forall x_1 \exists x_2 \forall x_3 :$

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_1 \vee x_2 \vee x_3)$$

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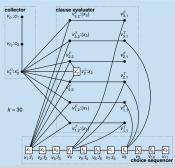
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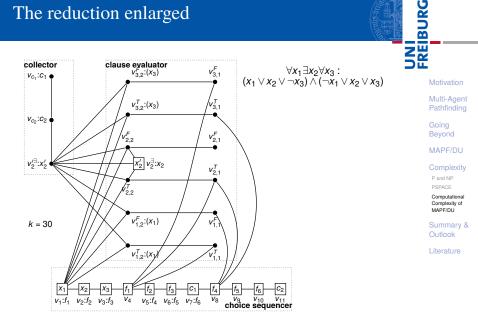
Literature

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The reduction enlarged



Complexity with a fixed number of agents

These results probably imply that the technique could not be used online.

For a fixed number of agents, however, the bounded planning problem is polynomial.

Theorem

For a fixed number c of agents, deciding whether there exists a MAPF/DU i-strong or objectively strong plan with execution cost of k or less can be done in time $O(n^{c^2+c})$.

That means, for two agents, it takes "only" $O(n^6)$ time – but in practice it should be faster.

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Summary & Outlook

- UNI FREIBURG
- 1 Determine in the state space of all node assignments the distance to the initial state using Dijkstra: $O(|V|^4)$ time.



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- 1 Determine in the *state space of all node assignments* the distance to the initial state using Dijkstra: $O(|V|^4)$ time.
- 2 For each of the $O(|V|^2)$ configurations check, whether it is a *potential stepping stone* for one agent, i.e., whether all potential destinations of this agent are reachable using Dijkstra on the modified graph, where the other agent blocks the way: $O(|V|^4)$ time.

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Summary & Outlook

- **1** Determine in the *state space of all node assignments* the distance to the initial state using Dijkstra: $O(|V|^4)$ time.
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- Solution For all $O(|V|^2)$ potential stepping stones, check whether for each of the O(|V|) possible destination of the first agent, the second agent can reach its possible destinations and use Dijkstra to compute the shortest path: altogether $O(|V|^5)$ time.

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Summary & Outlook

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- Solution For all $O(|V|^2)$ potential stepping stones, check whether for each of the O(|V|) possible destination of the first agent, the second agent can reach its possible destinations and use Dijkstra to compute the shortest path: altogether $O(|V|^5)$ time.
- Consider all stepping stones and minimize over the maximum plan depth. Among the minimal plans select those that are eager for the planning agent.

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- DMAPF generalizes the MAPF problem by dropping the assumption that plans are generated centrally and then communicated.
- MAPF/DU generalizes the MAPF problem further by dropping the assumptions that destinations are common knowledge.
- A solution concept for this setting are *i*-strong branching plans corresponding to implicitly coordinated policies in the area of epistemic planning.
- The backbone of such plans are stepping stones.
- Joint execution can be guaranteed to be successful and polynomially bounded if all agents are conservative and optimally eager.
- While plan existence in general is PSPACE-complete, it is polynomial for a fixed number of agents.

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Outlook

- $\rightarrow\,$ Do the results still hold for planar graphs?
 - Is MAPF/DU plan existence also PSPACE-complete?
 - How would more general forms of describing the common knowledge about destinations affect the results?
- $\rightarrow~$ Overlap of destinations or general Boolean combinations
- Can we get similar results for other execution semantics?
- → Concurrent executions of actions
- Can we be more aggressive in expectations about possible destinations?
- $\rightarrow\,$ Use forward induction, i.e., assume that actions in the past were rational.
- Are other forms of implicit coordination possible?
- → More communication? Coordination in competitive scenarios?

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