

Introduction to Game Theory

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2.6. Zero-sum Games and Nash Equilibria

Definition: A zero-sum game is a finite 2 player strategic game

$$G = \langle \{1, 2\}, (A)_{1, 2}, (u_i)_{1, 2} \rangle$$

such that for all profiles $a, b \in A$

$$u_1(a) + u_2(b) = 0.$$

Constant sum games: $\exists c \in \mathbb{R}$
 $u_1(a) + u_2(a) = c$
 for all $a \in A$.

Example:

		L	M	R	
		1	8, 8 3, -3 -6, 6 -6		
		1	2, -2 -1, 1 3, -3 -1		
		3	-6, 6 4, -4 8, -8 -6		
		8	4	6	-8

Idea: Choose action that is best for "you"
 make the assumption that the other player wants to minimize your utility.

Definition

Let $G \subseteq \mathbb{Z}^G$. $x^* \in A_1$ is called a maximizer for player 1, if

$$\min_{y \in A_2} u_1(x^*, y) \geq \min_{y \in A_2} u_1(x, y) \text{ for each } x \in A_1$$

$y^* \in A_2$ is maximizer of player 2 if

$$\min_{x \in A_1} u_2(x, y^*) \geq \min_{x \in A_1} u_2(x, y) \text{ for each } y \in A_2.$$

Consider as an example:

	L	R	
T	1, -1	2, -2	← 1
D	-2, 2	-4, 4	← -4
	↓	↓	
	-2	-2	

Profile (T, L) is a NE and a pair of maximizers.

We show: If G has a NE then any profile α that is NE is a pair of maximizers.

Lemma: Let G be a LSG. Then

$$\max_{y \in A_2} \min_{x \in A_1} u_2(x, y) =$$

$$= \min_{y \in A_2} \max_{x \in A_1} u_1(x, y)$$

Proof: Obviously, for each real-valued function f , it holds

$$(1) \quad \min_z (-f(z)) = -\max_z (f(z))$$

Thus, we obtain

$$(2) \quad - \min_{x \in A_1} u_2(x, y) \stackrel{(1)}{=} \max_{x \in A_1} -u_2(x, y)$$

$$\stackrel{-u_2 = u_1}{=} \max_{x \in A_1} u_1(x, y)$$

Thus,

$$\max_{y \in A_2} \min_{x \in A_1} u_2(x, y) \stackrel{(1)}{=} -\min_{y \in A_2} -(\min_{x \in A_1} u_2(x, y))$$

$$= -\min_{y \in A_2} \max_{x \in A_1} u_1(x, y) \quad \blacksquare$$

Proposition : Let G be a 2SGo

- (a) Whenever (x^*, y^*) is a NE of G , then
 x^* and y^* are maximinimizers of player 1
and 2, respectively.
- (b) If (x^*, y^*) is a NE of G , then

$$\max_{x \in A_1} \min_{y \in A_2} u_1(x, y)$$

$$= \min_{y \in A_2} \max_{x \in A_1} u_1(x, y)$$

$$= u_1(x^*, y^*) \quad \leftarrow -u_2(x^*, y^*)$$

Thus all NE of G have the same payoff.

- (c) If $\max_x \min_y u_1(x, y) = \min_y \max_x u_1(x, y)$ and
 x^* and y^* are maximinimizers of player 1 and
2, respectively, then (x^*, y^*) is a NE. In particular,
if (x_1^*, y_1^*) and (x_2^*, y_2^*) are NE of G , then so are (x_1^*, y_2^*)
and (x_2^*, y_1^*) .

Proof:

(a) + (b). Let (x^*, y^*) be a NE of \mathcal{G}_0 .

By defn of NE,

$$u_2(x^*, y^*) \geq u_2(x^*, y) \text{ for each } y \in A_2$$

From $u_2 = -u_1$ it follows

$$u_1(x^*, y^*) \leq u_1(x^*, y), \text{ for each } y \in A_2.$$

and hence

$$\begin{aligned} u_1(x^*, y^*) &= \min_{y \in A_2} u_1(x^*, y) \\ (1) \quad &\leq \max_{x \in A_1} \min_{y \in A_2} u_1(x, y). \end{aligned}$$

Moreover by def of NE,

$$u_1(x^*, y^*) \geq u_1(x, y^*) \text{ for each } x \in A_1.$$

and hence,

$$u_1(x^*, y^*) = \min_{y \in A_2} u_1(x, y) \text{ for each } x \in A_1$$

Let thus,

$$(2) \quad u_1(x^*, y^*) \geq \max_{x \in A_1} \min_{y \in A_2} u_1(x, y).$$

From (1) & (2) we obtain:

$$(3) \quad u_1(x^*, y^*) = \max_{x \in A_1} \min_{Y \in A_2} u_1(x, y)$$

Thus by (1),

$$\min_{Y \in A_2} u_1(x^*, y) \geq \min_y u_1(x, y) \quad \text{for each } x \in A_1.$$

i.e. x^* is a maximinizer of player 1.

In a similar way, prove that y^* is a maximinizer of player 2 and that

$$(4) \quad u_2(x^*, y^*) = \max_{Y \in A_2} \min_{x \in A_1} u_2(x, y)$$

This shows (a).

Moreover, we have:

$$\begin{aligned} u_1(x^*, y^*) &\stackrel{(3)}{=} \max_{x \in A_1} \min_{Y \in A_2} u_1(x, y) \\ &= -\max_{Y \in A_2} -\min_{x \in A_1} u_1(x, y) \\ &\stackrel{\text{Lemma}}{=} -\max_{Y \in A_2} \min_{x \in A_1} u_2(x, y) \end{aligned}$$

This shows (b).

References

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