

Principles of Knowledge Representation and Reasoning

Description Logics – Algorithms

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Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Description Logics – Algorithms

- 1 Motivation
- 2 Structural Subsumption Algorithms
- 3 Tableau Subsumption Method

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

- *Satisfiability* or *subsumption* of concept descriptions
- *Satisfiability* or *instance relation* in ABoxes
- **Structural subsumption algorithms**
 - *Normalization* of concept descriptions and *structural comparison*
 - very fast, but can only be used for small DLs
- **Tableau algorithms**
 - Similar to modal tableau methods
 - Meanwhile the method of choice

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

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Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Reasoning Problems & Algorithms

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature

Structural Subsumption Algorithms

- **Small Logic** \mathcal{FL}^-

- $C \sqcap D$
- $\forall r.C$
- $\exists r$ (simple existential quantification)

- **Idea**

- 1 In the conjunction, collect all *universally quantified expressions* (also called *value restrictions*) with the same role and build *complex value restriction*:

$$\forall r.C \sqcap \forall r.D \rightarrow \forall r.(C \sqcap D).$$

- 2 Compare all conjuncts with each other. For each conjunct in the subsuming concept there should be a *corresponding one* in the subsumed one.

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Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Structural Subsumption Algorithms

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Structural Subsumption Algorithms

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural Subsumption Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau Subsumption Method

Literature

Structural Subsumption Algorithms

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KRR

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Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Structural Subsumption Algorithms

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Example

$$\begin{aligned}D &= \text{Human} \sqcap \exists \text{has-child} \sqcap \forall \text{has-child}.\text{Human} \sqcap \\ &\quad \forall \text{has-child}.\exists \text{has-child} \\ C &= \text{Human} \sqcap \text{Female} \sqcap \exists \text{has-child} \sqcap \\ &\quad \forall \text{has-child}.\text{(Human} \sqcap \text{Female} \sqcap \exists \text{has-child)}\end{aligned}$$

Check: $C \sqsubseteq D$

- 1 Collect value restrictions in D :
... $\forall \text{has-child}.\text{(Human} \sqcap \exists \text{has-child)}$
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$\rightsquigarrow C$ is subsumed by D !

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Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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$$C = \text{Human} \sqcap \text{Female} \sqcap \exists \text{has-child} \sqcap \\ \forall \text{has-child}.\text{(Human} \sqcap \text{Female} \sqcap \exists \text{has-child)}$$

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

$$\begin{aligned} D &= \text{Human} \sqcap \exists \text{has-child} \sqcap \forall \text{has-child}.\text{Human} \sqcap \\ &\quad \forall \text{has-child}.\exists \text{has-child} \\ C &= \text{Human} \sqcap \text{Female} \sqcap \exists \text{has-child} \sqcap \\ &\quad \forall \text{has-child}.\text{(Human} \sqcap \text{Female} \sqcap \exists \text{has-child)} \end{aligned}$$

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Example

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Subsumption Algorithm

SUB(C, D) algorithm:

- 1 Reorder terms (*commutativity*, *associativity* and *value restriction law*):

$$C = \bigwedge A_i \cap \bigwedge \exists r_j \cap \bigwedge \forall r_k : C_k$$

$$D = \bigwedge B_l \cap \bigwedge \exists s_m \cap \bigwedge \forall s_n : D_n$$

- 2 For each B_l in D , is there an A_i in C with $A_i = B_l$?
- 3 For each $\exists s_m$ in D , is there an $\exists r_j$ in C with $s_m = r_j$?
- 4 For each $\forall s_n : D_n$ in D , is there a $\forall r_k : C_k$ in C such that $C_k \sqsubseteq D_n$ and $s_n = r_k$?

$\rightsquigarrow C \sqsubseteq D$ iff all questions are answered positively

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Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

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- 2 For each B_l in D , is there an A_i in C with $A_i = B_l$?
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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Subsumption Algorithm

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Subsumption Algorithm

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

Theorem (Soundness)

$$SUB(C, D) \Rightarrow C \sqsubseteq D$$

Proof sketch.

Reordering of terms (1):

a) Commutativity and associativity are trivial

b) Value restriction law. We show: $(\forall r.(C \sqcap D))^{\mathcal{I}} = (\forall r.C \sqcap \forall r.D)^{\mathcal{I}}$

Assumption: $d \in (\forall r.(C \sqcap D))^{\mathcal{I}}$

Case 1: $\nexists e : (d, e) \in r^{\mathcal{I}} \quad \checkmark$

Case 2: $\exists e : (d, e) \in r^{\mathcal{I}} \Rightarrow e \in (C \sqcap D)^{\mathcal{I}} \Rightarrow e \in C^{\mathcal{I}}, e \in D^{\mathcal{I}}$

Since e is arbitrary: $d \in (\forall r.C)^{\mathcal{I}}, d \in (\forall r.D)^{\mathcal{I}}$ then

d must also be conjunction, i.e.,

$$(\forall r.(C \sqcap D))^{\mathcal{I}} \subseteq (\forall r.C \sqcap \forall r.D)^{\mathcal{I}}$$

Other direction is similar

(2+3+4): Induction on the nesting depth of \forall -expressions □

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Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Soundness

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Reordering of terms (1):

a) Commutativity and associativity are trivial

b) Value restriction law. We show: $(\forall r.(C \sqcap D))^{\mathcal{I}} = (\forall r.C \sqcap \forall r.D)^{\mathcal{I}}$

Assumption: $d \in (\forall r.(C \sqcap D))^{\mathcal{I}}$

Case 1: $\nexists e : (d, e) \in r^{\mathcal{I}} \quad \checkmark$

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d must also be conjunction, i.e.,

$$(\forall r.(C \sqcap D))^{\mathcal{I}} \subseteq (\forall r.C \sqcap \forall r.D)^{\mathcal{I}}$$

Other direction is similar

(2+3+4): Induction on the nesting depth of \forall -expressions □

KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption

Method

Literature

Soundness

Theorem (Soundness)

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

Theorem (Soundness)

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(2+3+4): Induction on the nesting depth of \forall -expressions □

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

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(2+3+4): Induction on the nesting depth of \forall -expressions □

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

Theorem (Soundness)

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(2+3+4): Induction on the nesting depth of \forall -expressions □

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

Theorem (Soundness)

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Other direction is similar

(2+3+4): Induction on the nesting depth of \forall -expressions □

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Soundness

Theorem (Soundness)

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(2+3+4): Induction on the nesting depth of \forall -expressions



KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Completeness

Theorem (Completeness)

$$C \sqsubseteq D \Rightarrow \text{SUB}(C, D)$$

Proof idea.

One shows the contrapositive:

$$\neg \text{SUB}(C, D) \Rightarrow C \not\sqsubseteq D$$

Idea: If one of the rules leads to a negative answer, we use this to construct an interpretation with a special element d such that

$$d \in C^{\mathcal{I}}, \text{ but } d \notin D^{\mathcal{I}}$$



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Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Completeness

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Completeness

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Completeness

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Generalizing the Algorithm

Extensions of \mathcal{FL}^- by

- $\neg A$ (*atomic negation*),
- $(\leq nr)$, $(\geq nr)$ (*cardinality restrictions*),
- $r \circ s$ (*role composition*)

does not lead to any problems.

However: If we use full existential restrictions, then it is very unlikely that we can come up with a *simple* structural subsumption algorithm – having the same flavor as the one above.

More precisely: There is (most probably) no algorithm that uses polynomially many reorderings and simplifications and allows for a simple structural comparison

Reason: Subsumption for $\mathcal{FL}^- + \exists r.C$ is NP-hard (Nutt).

KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Generalizing the Algorithm

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Generalizing the Algorithm

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Generalizing the Algorithm

Extensions of \mathcal{FL}^- by

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

Generalizing the Algorithm

Extensions of \mathcal{FL}^- by

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
ABox Reasoning

Tableau
Subsumption
Method

Literature

Generalizing the Algorithm

Extensions of \mathcal{FL}^- by

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Idea

Example

Algorithm

Soundness

Completeness

Generalizations

ABox Reasoning

Tableau

Subsumption
Method

Literature

ABox Reasoning

Idea: *abstraction* + *classification*

- *Complete* AB_{OX} by propagating value restrictions to role fillers
- Compute for each object its *most specialized concepts*
- These can then be handled using the ordinary subsumption algorithm

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Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Idea
Example
Algorithm
Soundness
Completeness
Generalizations
AB_{OX} Reasoning

Tableau
Subsumption
Method

Literature

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- *Complete* ABBox by propagating value restrictions to role fillers
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- **Logic** \mathcal{ALC}
 - $C \sqcap D$
 - $C \sqcup D$
 - $\neg C$
 - $\forall r.C$
 - $\exists r.C$
- **Idea**: Decide (un-)satisfiability of a concept description C by trying to *systematically construct* a model for C . If that is successful, C is satisfiable. Otherwise C is unsatisfiable.

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Example: Subsumption in a TBox

TBox

Hermaphrodite \doteq Male \sqcap Female

Parents-of-sons-and-daughters \doteq

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female}$

Parents-of-hermaphrodite $\doteq \exists \text{has-child.Hermaphrodite}$

Query

Parents-of-sons-and-daughters $\sqsubseteq_{\mathcal{T}}$

Parents-of-hermaphrodites

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Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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Parents-of-hermaphrodites

KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Reductions

1 *Unfolding*

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqsubseteq$
 $\exists \text{has-child.}(\text{Male} \sqcap \text{Female})$

2 *Reduction to unsatisfiability*

Is

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqcap$
 $\neg(\exists \text{has-child.}(\text{Male} \sqcap \text{Female}))$
unsatisfiable?

3 *Negation normal form* (move negations inside):

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqcap$
 $\forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female})$

4 *Try to construct a model*

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Reductions

1 *Unfolding*

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqsubseteq$
 $\exists \text{has-child.}(\text{Male} \sqcap \text{Female})$

2 *Reduction to unsatisfiability*

Is
 $\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqcap$
 $\neg(\exists \text{has-child.}(\text{Male} \sqcap \text{Female}))$
unsatisfiable?

3 *Negation normal form* (move negations inside):

$\exists \text{has-child.Male} \sqcap \exists \text{has-child.Female} \sqcap$
 $\forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female})$

4 *Try to construct a model*

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Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Model Construction (1)

- ① **Assumption:** There exists an object x in the interpretation of our concept:

$$x \in (\exists \dots)^{\mathcal{I}}$$

- ② This implies that x is in the interpretation of all conjuncts:

$$x \in (\exists \text{has-child.Male})^{\mathcal{I}}$$

$$x \in (\exists \text{has-child.Female})^{\mathcal{I}}$$

$$x \in (\forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female}))^{\mathcal{I}}$$

- ③ This implies that there should be objects y and z such that $(x, y) \in \text{has-child}^{\mathcal{I}}$, $(x, z) \in \text{has-child}^{\mathcal{I}}$, $y \in \text{Male}^{\mathcal{I}}$ and $z \in \text{Female}^{\mathcal{I}}$ and ...

KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

**Model
Construction**

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

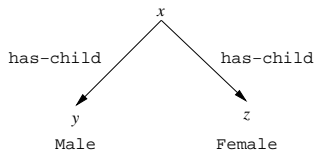
Space
Complexity

ABox Reasoning

Model Construction (2)

$x : \exists \text{has-child.Male}$

$x : \exists \text{has-child.Female}$



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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

**Model
Construction**

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

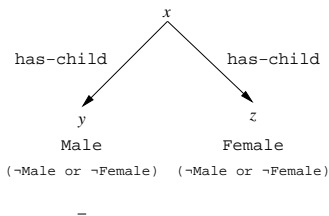
ABox Reasoning

Model Construction (3)

$x : \exists \text{has-child.Male}$

$x : \exists \text{has-child.Female}$

$x : \forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female})$



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Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

**Model
Construction**

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

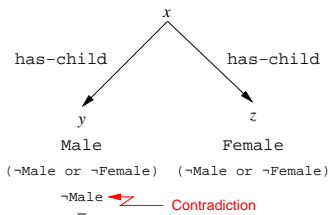
Model Construction (4)

$x : \exists \text{has-child.Male}$

$x : \exists \text{has-child.Female}$

$x : \forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female})$

$y : \neg \text{Male}$



KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Model Construction (5)

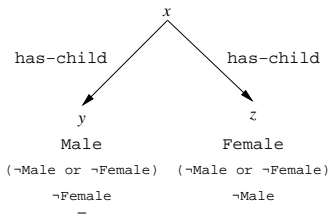
$x : \exists \text{has-child.Male}$

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$y : \neg \text{Female}$

$z : \neg \text{Male}$



\rightsquigarrow Model **constructed!**

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

**Model
Construction**

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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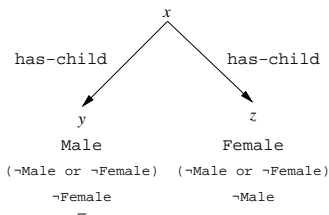
$x : \exists \text{has-child.Male}$

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$x : \forall \text{has-child.}(\neg \text{Male} \sqcup \neg \text{Female})$

$y : \neg \text{Female}$

$z : \neg \text{Male}$



\rightsquigarrow Model **constructed!**

KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (1): NNF

$C \equiv D$ iff $C \sqsubseteq D$ and $D \sqsubseteq C$.

Now we have the following equivalences:

$$\neg(C \sqcap D) \equiv \neg C \sqcup \neg D$$

$$\neg(C \sqcup D) \equiv \neg C \sqcap \neg D$$

$$\neg\neg C \equiv C$$

$$\neg(\forall r.C) \equiv \exists r.\neg C$$

$$\neg(\exists r.C) \equiv \forall r.\neg C$$

These equivalences can be used to move all negations signs to the inside, resulting in concept description where only concept names are negated: **negation normal form (NNF)**

Theorem (NNF)

The negation normal form of an \mathcal{ALC} concept can be computed in polynomial time.

KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural Subsumption Algorithms

Tableau Subsumption Method

Example

Reductions: Unfolding & Unsatisfiability

Model Construction

Equivalences & NNF

Constraint Systems

Transforming Constraint Systems

Invariances

Soundness and Completeness

Space Complexity

ABox Reasoning

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (2): Constraint Systems

A **constraint** is a syntactical object of the form: $x: C$ or xry , where C is a concept description in NNF, r is a role name and x and y are *variable names*.

Let \mathcal{I} be an interpretation. An \mathcal{I} -assignment α is a function that maps each variable symbol to an object of the universe \mathcal{D} .

A *constraint* $x: C$ (xry) is **satisfied** by an \mathcal{I} -assignment α , if $\alpha(x) \in C^{\mathcal{I}}$ ($(\alpha(x), \alpha(y)) \in r^{\mathcal{I}}$).

A **constraint system** S is a finite, non-empty set of constraints. An \mathcal{I} -assignment α satisfies S if α satisfies each constraint in S . S is **satisfiable** if there exists \mathcal{I} and α such that α satisfies S .

Theorem

An ACC concept C in NNF is satisfiable iff the system $\{x: C\}$ is satisfiable.

KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (2): Constraint Systems

A **constraint** is a syntactical object of the form: $x: C$ or xry , where C is a concept description in NNF, r is a role name and x and y are *variable names*.

Let \mathcal{I} be an interpretation. An **\mathcal{I} -assignment** α is a function that maps each variable symbol to an object of the universe \mathcal{D} .

A **constraint** $x: C$ (xry) is **satisfied** by an \mathcal{I} -assignment α , if $\alpha(x) \in C^{\mathcal{I}}$ ($(\alpha(x), \alpha(y)) \in r^{\mathcal{I}}$).

A **constraint system** S is a finite, non-empty set of constraints. An \mathcal{I} -assignment α satisfies S if α satisfies each constraint in S . S is **satisfiable** if there exists \mathcal{I} and α such that α satisfies S .

Theorem

An ACC concept C in NNF is satisfiable iff the system $\{x: C\}$ is satisfiable.

KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl, Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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An \mathcal{ALC} concept C in NNF is satisfiable iff the system $\{x: C\}$ is satisfiable.

KRR

Nebel, Wöflf,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (3): Transforming Constraint Systems

Transformation rules:

- 1 $S \rightarrow_{\sqcap} \{x: C_1, x: C_2\} \cup S$
if $(x: C_1 \sqcap C_2) \in S$ and either $(x: C_1)$ or $(x: C_2)$ or both are not in S .
- 2 $S \rightarrow_{\sqcup} \{x: D\} \cup S$
if $(x: C_1 \sqcup C_2) \in S$ and neither $(x: C_1) \in S$ nor $(x: C_2) \in S$ and $D = C_1$ *or* $D = C_2$.
- 3 $S \rightarrow_{\exists} \{xry, y: C\} \cup S$
if $(x: \exists r.C) \in S$, y is a *fresh variable*, and there is no z s.t. $(xrz) \in S$ and $(z: C) \in S$.
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Deterministic rules (1,3,4) vs. non-deterministic (2).

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems
**Transforming
Constraint
Systems**

Invariances
Soundness and
Completeness

Space
Complexity
ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems
**Transforming
Constraint
Systems**

Invariances
Soundness and
Completeness

Space
Complexity
ABox Reasoning

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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (4): Invariances

Theorem (Invariance)

Let S and T be constraint systems:

- 1 If T has been generated by applying a deterministic rule to S , then S is satisfiable iff T is satisfiable.*
- 2 If T has been generated by applying a non-deterministic rule to S , then S is satisfiable if T is satisfiable. Furthermore, if a non-deterministic rule can be applied to S , then it can be applied such that S is satisfiable iff the resulting system T is satisfiable.*

Theorem (Termination)

Let C be an \mathcal{ALC} concept description in NNF. Then there exists no infinite chain of transformations starting from the constraint system $\{x: C\}$.

KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Tableau Method (5): Soundness and Completeness

A constraint system is called **closed** if no transformation rule can be applied.

A **clash** is a pair of constraints of the form $x : A$ and $x : \neg A$, where A is a concept name.

Theorem (Soundness and Completeness)

A closed constraint system is satisfiable iff it does not contain a clash.

Proof idea.

\Rightarrow : obvious. \Leftarrow : Construct a model by using the concept labels.

KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wöflfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example

Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

Space Requirements

Because the tableau method is *non-deterministic* ($\rightarrow \sqcup$ rule) ... there could be exponentially many closed constraint systems in the end.

Interestingly, even one constraint system can have *exponential size*.

Example:

$$\begin{array}{c} \exists r. A \sqcap \exists r. B \sqcap \\ \forall r. \left(\begin{array}{c} \exists r. A \sqcap \exists r. B \sqcap \\ \forall r. \left(\begin{array}{c} \exists r. A \sqcap \exists r. B \sqcap \\ \forall r. (\dots) \end{array} \right) \end{array} \right) \end{array}$$

However: One can modify the algorithm so that it needs only poly. space.

Idea: Generating a y only for one $\exists r. C$ and then proceeding into the depth.

KRR

Nebel, Wölfel,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems
Transforming
Constraint
Systems

Invariances
Soundness and
Completeness

Space
Complexity
ABox Reasoning

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KRR

Nebel, Wölfel,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction
Equivalences &
NNF

Constraint
Systems
Transforming
Constraint
Systems

Invariances
Soundness and
Completeness

Space
Complexity
ABox Reasoning

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KRR

Nebel, Wölfel,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

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KRR

Nebel, Wölfel,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Example
Reductions:
Unfolding &
Unsatisfiability

Model
Construction

Equivalences &
NNF

Constraint
Systems

Transforming
Constraint
Systems

Invariances

Soundness and
Completeness

Space
Complexity

ABox Reasoning

ABox satisfiability can also be decided using the tableau method if we can add constraints of the form $x \neq y$ (for *UNA*):

- *Normalize* and *unfold* and add inequalities for all pairs of objects mentioned in the ABox.
- Strictly speaking, in *ALC* we do not need this because we are never *forced* to identify two objects.






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KRR

Nebel, Wöfl,
Ragni

Motivation

Structural
Subsumption
Algorithms

Tableau
Subsumption
Method

Literature