

## Foundations of Artificial Intelligence

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### Exercise Sheet 5

**Due: Tuesday, June 9, 2009**

#### Exercise 5.1 (CNF Transformation, Resolution Method)

The following transformation rules hold, whereby propositional formulae can be transformed into equivalent formulae. Here,  $\varphi$ ,  $\psi$  and  $\chi$  are arbitrary propositional formulae:

$$\neg\neg\varphi \equiv \varphi \quad (1)$$

$$\neg(\varphi \vee \psi) \equiv \neg\varphi \wedge \neg\psi \quad (2)$$

$$\varphi \vee (\psi \wedge \chi) \equiv (\varphi \vee \psi) \wedge (\varphi \vee \chi) \quad (3)$$

$$\neg(\varphi \wedge \psi) \equiv \neg\varphi \vee \neg\psi \quad (4)$$

$$\varphi \wedge (\psi \vee \chi) \equiv (\varphi \wedge \psi) \vee (\varphi \wedge \chi) \quad (5)$$

Additionally, the operators  $\vee$  and  $\wedge$  are associative and commutative.  
Consider the formula  $((C \wedge \neg B) \leftrightarrow A) \wedge (\neg C \rightarrow A)$ .

- Transform the formula into a clause set  $K$  using the CNF transformation rules. Write down the steps.
- Afterwards, using the resolution method, show that  $K \models (\neg B \rightarrow (A \wedge C))$  holds.

#### Exercise 5.2 (Davis-Putnam Procedure)

Use the *Davis-Putnam* procedure to find a model for each of the following formulae, or prove that the particular formula has no model. At each step, indicate which rule you have applied.

(a)  $(P \vee \neg Q) \wedge (\neg P \vee Q) \wedge (Q \vee \neg R) \wedge (\neg Q \vee \neg R)$

(b)  $(P \vee Q \vee \neg R) \wedge (P \vee \neg Q) \wedge \neg P \wedge \neg R \wedge \neg U$

#### Exercise 5.3 (Modelling in First-Order Logic)

Represent the following sentences in first-order logic, using a consistent vocabulary.

- Everyone who has a good sense of hearing is able to sing.
- Nobody is a real musician, if he is not able to inspire his audience.
- Nobody, except for a real musician can write a symphony.

Which properties does someone necessarily possess when he wrote a symphony?

**Exercise 5.4** (Skolem Normal Form)

Transform the following formulae to Skolem normal form:

(a)  $F_1 = \forall x(\exists yR(x, y) \wedge \exists yR(y, x))$

(b)  $F_2 = \forall x\forall z(R(x, z) \Rightarrow \exists y(R(x, y) \wedge R(y, z)))$

(c)  $F_3 = \forall x\exists z(R(x, z) \wedge \neg\exists y(R(x, y) \wedge R(y, z)))$

**Exercise 5.5** (Herbrand Expansion)

Let  $F = \forall x\forall y(P(x, f(x, g(y))) \wedge P(h(y), f(y, y)))$ .

- (a) State ten minimally large terms from the Herbrand universe of  $F$ .
- (b) State five minimally large formulae from the Herbrand expansion of  $F$ .

The exercise sheets may and should be handed in and be worked on in groups of three (3) students. Please fill the cover sheet<sup>1</sup> and attach it to your solution.

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<sup>1</sup><http://www.informatik.uni-freiburg.de/~ki/teaching/ss09/gki/coverSheet-english.pdf>