## Planning with partial observability

- Combines the difficulties needed in the unobservable and fully observable cases.
- Sequential plans (like with unobservability) or state $\rightarrow$ action plans (like with full observability) do not suffice.
- In principle solvable by reduction to nondeterministic fully observable planning in the belief space. But this is impractical because of $2^{n}$ belief states for $n$ states.
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5. Create node $n^{\prime}$. Assign $l(n):=\left\langle o, n^{\prime}\right\rangle$ and $B S\left(n^{\prime}\right):=\operatorname{img}_{o}(B)$.
6. Make $n^{\prime}$ a branch node with $l\left(n^{\prime}\right)=\left\{\left\langle C_{1}, n_{1}\right\rangle, \ldots,\left\langle C_{m}, n_{m}\right\rangle\right\}$ where $n_{1}, \ldots, n_{m}$ are new nodes. Assign $B S\left(n_{i}\right):=i m g_{o}(B) \cap$ $C_{i}$ for every $i \in\{1, \ldots, m\}$.
7. Go to step 3

Nondeterministic choice in step 4 is implemented as search
Prevention of infinite plans: no node $n^{\prime}$ following $n$ may fulfill $B S(n) \subseteq B S\left(n^{\prime}\right)$.

## A simple forward search algorithm

1. Initialize the plan with node $b$ and assign $B S(b):=I$.
2. Make $b$ a branch node with $l(b)=\left\{\left\langle C_{1}, n_{1}\right\rangle, \ldots,\left\langle C_{m}, n_{m}\right\rangle\right\}$ where $n_{1}, \ldots, n_{m}$ are new nodes. Assign $B S\left(n_{i}\right):=I \cap C_{i}$ for all $i \in\{1, \ldots, m\}$.
3. Choose a node $n$ with $B S(n) \nsubseteq G$ and with $l(n)=\emptyset$.

If there is no such node, plan is complete.
4. Nondeterministically choose $o \in O$ that is applicable in $B$.

## Backward search algorithms

- Flavor similar to the backward algorithms for fully observable problems.
- Backward steps with operator applications: strong preimages.
- Backward steps with branching: we present a new construction for doing this.


## Branching in backward search

- Let the observational classes be $C_{1}, \ldots, C_{n}$.
- Let $S_{1}, S_{2}, \ldots, S_{n}$ be sets of states with plans so that for all $i, j$ such that $i \neq j$ there is no $C \in\left\{C_{1}, \ldots, C_{n}\right\}$, such that $S_{i} \cap C \neq \emptyset$ and $S_{j} \cap C \neq \emptyset$.

Now they can be combined to $S=S_{1} \cup \cdots \cup S_{n}$ that has a plan starting with a branch.

- Where do such sets $S_{i}$ come from?


## Regression/preimages




## Branching



## Combination 11



## Combination 12



Combination 22
$\begin{array}{lllllll}\text { o1 } & \text { o2 } & \text { o3 } & \text { o4 } & \text { o5 } & \text { o6 } & \text { o7 }\end{array}$


No observability $\Rightarrow$ No branching



Combination with full observability


## Algorithm outline

- Pick from each observational class one belief state.
- Compute the strong preimage of their union w.r.t. operator $o$.
- Split the resulting set of states to belief states for different observational classes.
- Objective: obtain new belief states, preferably closer to $I$.


## Example: backward search in the belief space

- 3 blocks A, B and C
- Goal: all blocks are on the table
- Only the variables clear $(X)$ are observable.
- A block can be moved onto the table if the block is clear.
- 8 observational classes corresponding to the 8 valuations of $\{$ clear $(A)$, clear $(B)$, clear $(C)\}$ (one of the valuations does not correspond to a blocks world state.)


## Algorithm idea: construction of plans

If plans for belief states $Z_{1}, \ldots, Z_{n}$, respectively corresponding to observational classes $C_{1}, \ldots, C_{n}$, were $\pi_{1}, \ldots, \pi_{n}$, the plan for a new belief state is

1. Apply $o$
2. If new current state is in $C_{i}$ for $i \in\{1, \ldots, n\}$, continue with $\pi_{i}$.
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## Example: goal belief state



Example: backup step with A-onto-table


Example: backup step with C-onto-table


Example: backup step with B-onto-table


Example: backup step with A-onto-table


Example: backup step with B-onto-table


Example: backup step with A-onto-table


Example: backup step with B-onto-table


Example: backup step with C-onto-table


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