Observability

- Robot can see and hear only the immediate surroundings.
- Poker player cannot see the opponents' cards.
- Formalization: split state variables to those that can be, and to those that cannot be observed.
- More general formalization: set of atomic and compound formulas; truth-values of these formulas can be observed, but not necessarily the truth-values of their subformulas.

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Problem definition

A 5-tuple $\langle P, I, O, G, B \rangle$ consisting of

- a set *P* of state variables,
- a propositional formula *I* over *P*,
- a set O of operators,
- a propositional formula G over P, and
- a set $B \subseteq P$ of observable state variables.

is a problem instance in nondeterministic planning with partial observability.

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Conditional plans: definition

B = the observable state variables, O = the operators. A plan is $\langle N, b, l \rangle$ where

- *N* is a finite set of nodes,
- $b \in N$ is the initial node,
- $l: N \to (O \times N) \cup 2^{\mathcal{L} \times N}$ assigns each node
- an operator and a successor node $\langle o, n \rangle \in O \times N$ or
- a set of conditions and successor nodes $\langle \phi, n \rangle$, where $n \in N$ and ϕ is a formula over B.

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Restrictions on observability

- Let $\langle P, I, O, G, B \rangle$ be a problem instance in conditional planning.
- If P = B, the problem instance is *fully observable*.

If $B = \emptyset$, the problem instance is *unobservable*.

If $B \neq \emptyset$ and $P \neq B$, the problem instance is *partially observable*. We also use this term for the general case that includes full observability and unobservability.

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Observational classes

- When only variables in B = {p₁,..., p_m} are observable, then it is not possible to distinguish between states s and s' such that s(p) = s'(p) for all p ∈ B.
- Observability partitions states to classes S₁ ∪ S₂ ∪ · · · ∪ S_n of indistinguishable states. One class for every valuation of the observable state variables.

$$S = S_1 \cup S_2 \cup \dots \cup S_n,$$

$$S_i \cap S_j = \emptyset \text{ for any } \{i, j\} \subset \{1, \dots, n\} \text{ such that } i \neq j.$$

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Observational classes: cardinality

Full observability: state space $S = \{s_1, \ldots, s_n\}$ is partitioned to singleton classes $S_1 = \{s_1\}, S_2 = \{s_2\}, \ldots, S_n = \{s_n\}.$

Unobservability: The partition has only one class $S_1 = S$ consisting of all the states.

Partial observability: the number and cardinality of S_i may be anywhere between 1 and n.

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Belief states and the belief space

- Current state is not in general known during plan execution. Instead, the a set of possible current states is known.
- A set of possible current states is a belief state.
- The set of all belief states is the belief space.
- If there are n observationally indistinguishable states, then there are 2^n belief states.

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1. Let F be a belief state (a set of states). 2. Operator o is executable in F if it is executable in every s ∈ F. 3. When o is executed, possible next states are E = img_o(F). 4. Observations select some observational class S_j. 5. New belief state is F' = img_o(F) ∩ S_j. ⇒ There is an arc from F to F' in the transition graph for the belief space.









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Belief space: after WWWWWNN



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Belief space: after WWWWWWNNNNN



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Belief space: after WWWWWWNNNNNNE







Regression for nondetermin	nistic operators	
1. Nondeterministic operator $o = \langle c, e \rangle$ we normal formal II so that every e_i is defined as the every e_i is defined as the every e_i is defined as the every e_i and e_i as the every e_i and e_i as the every e_i as the e	with $e = p_1 e_1 \cdots p_n e_n$ in terministic.	
2. $\operatorname{regr}_{\langle c,e angle}(\phi)$ is now defined as $\operatorname{regr}_{\langle c,e_1}$	$_{\langle}(\phi)\wedge\cdots\wedgeregr_{\langle c,e_n angle}(\phi).$	
THEOREM. If $S' = \{s' s' \models \phi\}$, then regr _o (S')}.	$spreimg_o(S') = \{s s \models$	
PROOF IDEA: To guarantee reaching S be reached from s with every effect e_1 ,.	" from a state s , S' must, e_n .	
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Algorithms for UO planning: heuristic search

Regression + heuristic search (A*, best first, ...)

- new problem: Testing $I \models \operatorname{regr}_{o}(\phi)$ is co-NP-hard.
- heuristic 1: strong distances as in FO algorithms

If only one initial state, distance estimation as in deterministic planning can be used, but even accurate estimates for deterministic planning may be very misleading.

• heuristic 2: size of belief state

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Algorithms for UO planning: heuristic search

Use strong distances of states as a heuristic:

 $\begin{array}{rcl} D_0 &=& G\\ D_{i+1} &=& D_i \cup \bigcup_{o \in O} \textit{spreimg}_o(D_i) \textit{ for all } i \geq 1 \end{array}$

A lower bound on plan length for belief state Z is j if $Z \subseteq D_j$ and $Z \not\subseteq D_{j-1}$.

This is an **admissible heuristic** (does not overestimate the cost).

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Algorithms for UO planning: heuristic search

- Use the cardinality of belief states as a heuristic.
- Backward search: Prefer operators that *increase the size of the belief state*, i.e. find a plan suffix that reaches a goal state from more starting states.
- Forward search: Prefer operators that *decrease the size of the belief state*, i.e. reduce the uncertainty about the current state and make reaching goals easier.
- This heuristic is not admissible.

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