

Algorithm for constructing plans without loops

1. Compute distances to goals for all states by strong preimages.
2. Synthesize an acyclic plan in forward direction starting from the initial states.

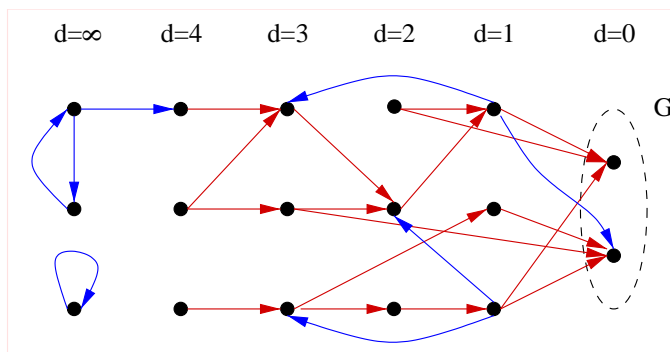
Operator selection: for each state, the operator has to decrease the remaining distance to goals by at least 1.

The algorithm: strong distances

```

PROCEDURE FOplan(I,O,G)
   $D_0 := G$ ;
   $i := 0$ ;
  WHILE  $I \not\subseteq D_i$  AND ( $i = 0$  OR  $D_{i-1} \neq D_i$ ) DO
     $D_i := D_{i-1} \cup \bigcup_{o \in O} \text{strong}_o(D_{i-1})$ ;
     $i := i + 1$ ;
  END
  
```

The algorithm: example of distances



The algorithm: plan construction

Initialize the edge labels and node set $N = \emptyset$, and call the recursive procedure FOplanconstruct.

```

 $N := \{0\}$ ;
 $l(j) := \emptyset$  for all  $j$ ;
 $cnt := 1$ ;
FOplanconstruct(0,I);
  
```

The algorithm: plan construction

```
PROCEDURE FOplanconstruct( $n, S$ )
FOR EACH  $o \in O$ 
... see next slide!
IF  $S \neq \emptyset$  THEN
  BEGIN (The remaining states are included in  $G$ .)
    cnt := cnt+1;
     $l(n) := l(n) \cup \{\langle S, cnt-1 \rangle\}$ ;
  END
END
```

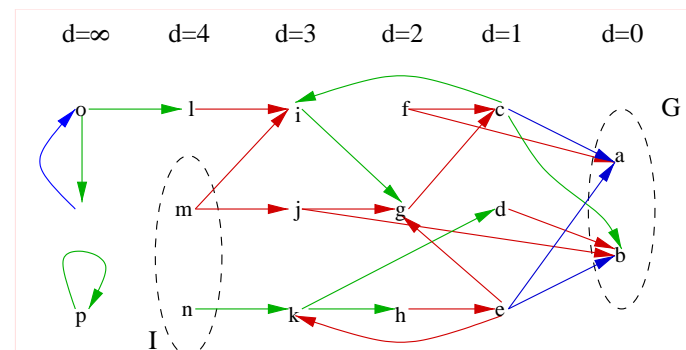
The algorithm: plan construction, cont'd

```
 $S' :=$  the maximal subset of  $S$  such that  $\text{progress}(o, S')$ ;
IF  $S' \neq \emptyset$  THEN
  BEGIN
     $S := S \setminus S'$ ;
    cnt := cnt+2;
     $N := N \cup \{\text{cnt-2}, \text{cnt-1}\}$ ;
     $l(n) := l(n) \cup \{\langle S', \text{cnt-2} \rangle\}$ ;
     $l(\text{cnt-2}) := \langle o, \text{cnt-1} \rangle$ ;
    FOplanconstruct(cnt-1,  $\text{img}_o(S')$ );
  END
```

The algorithm: a progress criterion

```
PROCEDURE progress( $o, S$ )
FOR  $j := 1$  TO  $i$  DO
  IF  $\text{img}_o(S \cap D_j) \not\subseteq D_{j-1}$  THEN RETURN false;
END
RETURN true;
```

The algorithm: plan construction, example

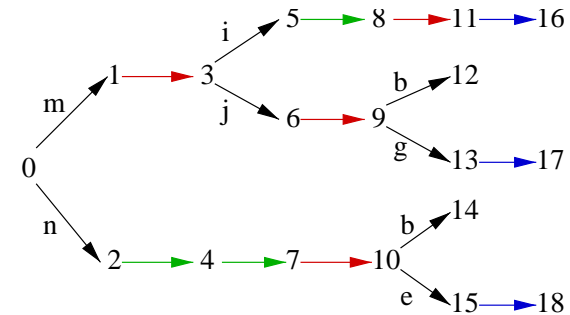


The algorithm: cont'd

Plan is $\langle \{0, \dots, 18\}, 0, l \rangle$ where

$l(0) = \{ \langle m, 1 \rangle, \langle n, 2 \rangle \}$	$l(1) = \langle R, 3 \rangle$
$l(2) = \langle G, 4 \rangle$	$l(3) = \{ \langle i, 5 \rangle, \langle j, 6 \rangle \}$
$l(4) = \langle G, 7 \rangle$	$l(5) = \langle G, 8 \rangle$
$l(6) = \langle R, 9 \rangle$	$l(7) = \langle R, 10 \rangle$
$l(8) = \langle R, 11 \rangle$	$l(9) = \{ \langle b, 12 \rangle, \langle g, 13 \rangle \}$
$l(10) = \{ \langle b, 14 \rangle, \langle e, 15 \rangle \}$	$l(11) = \langle B, 16 \rangle$
$l(12) = \emptyset$	$l(13) = \langle B, 17 \rangle$
$l(14) = \emptyset$	$l(15) = \langle B, 18 \rangle$
$l(16) = \emptyset, l(17) = \emptyset$	$l(18) = \emptyset$

The algorithm: plan construction, example



Memoryless plans

When we have **full observability**, a simpler definition of plans suffices.

- Plan is a mapping from states to operators.
- In program form: loop inside which one case statement chooses an operator for every state (or chooses termination.)

Memoryless plans: example

```

1: CASE
  a,b:      GOTO 5
  d,f,g,h,j,m: GOTO 2
  c,e:      GOTO 3
  i,k,n:    GOTO 4
2: RED
  GOTO 1
3: BLUE
  GOTO 1
4: GREEN
  GOTO 1
5:
  
```