Principles of AI Planning

17. Strong cyclic planning

Bernhard Nebel and Robert Mattmüller

February 5th, 2018
Strong cyclic plans
The simplest objective for nondeterministic planning is the one we have considered in the previous lecture: reach a goal state with certainty.

With this objective the nondeterminism can also be understood as an opponent like in 2-player games. The plan guarantees reaching a goal state no matter what the opponent does: plans are winning strategies.
In strong plans, goal states can be reached without visiting any state twice.

This property guarantees that the length of executions is bounded by some constant (which is smaller than the number of states.)

Some solvable problems are not solvable this way.

1. Action may fail to have any effect.
   - Hit a coconut to break it.

2. Action may fail and take us away from the goals.
   - Build a house of cards.

Consequences:

1. It is impossible to avoid visiting some states several times.
2. There is no finite upper bound on execution length.
Fairness assumption

For any nondeterministic operator $\langle \chi, \{e_1, \ldots, e_n\} \rangle$, the “probability” of every effect $e_i$, $i = 1, \ldots, n$, is greater than 0.

Alternatively: For each $s' \in \text{img}_o(s)$ the “probability” of reaching $s'$ from $s$ by $o$ is greater than 0.

This assumption guarantees that a strong cyclic plan reaches the goal almost certainly (with probability 1).

This is not compatible with viewing nondeterminism as an opponent in a 2-player game: the opponent’s strategy might rule out some of the choices $e_1, \ldots, e_n$. 
Need for strong cyclic plans

Example (Breaking a coconut)

- Initial state: coconut is intact.
- Goal state: coconut is broken.
- On every hit the coconut may or may not break.
- There is no finite upper bound on the number of hits.

This is equivalent to coin tossing.
Need for strong cyclic plans

Example

Example (Build a house of cards)
- Initial state: all cards lie on the table.
- Goal state: house of cards is complete.
- At every construction step the house may collapse.

Distance to $S_*$

$\infty$   $0$
We present two algorithms for strong cyclic planning:

- The **nested fixpoint algorithm** is conceptually simpler, but typically very costly, especially if not implemented symbolically.
  - Historically older
  - Uninformed
  - Considers entire state space

- The **determinization-based incremental planning algorithm** is a bit more complicated, but typically more efficient.
  - Historically newer, *state of the art*
  - Can use *informed* classical planner as sub-procedure
  - Often only considers small portion of state space
Nested Fixpoint Algorithm

Idea

■ Finds plans that may loop (strong cyclic plans).

■ The algorithm is rather tricky in comparison to the algorithm for strong plans.

■ Every state covered by a plan satisfies two properties:
  1. The state is **good**: there is at least one execution (= path in the graph defined by the plan) leading to a goal state.
  2. Every successor state is either a goal state or good.

■ The algorithm repeatedly eliminates states that are not good.
Nested Fixpoint Algorithm

Example

\[ S^* \]
All states are candidates for being good.
States from which goals are reachable in $\leq 1$ steps so that all immediate successors are possibly good.
States from which goals are reachable in \( \leq 2 \) steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 3$ steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 4$ steps so that all immediate successors are possibly good.
Eliminate states that turned out not to be good.
The set of possibly good states is now smaller.
States from which goals are reachable in \( \leq 1 \) steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 2$ steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 3$ steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 4$ steps so that all immediate successors are possibly good.
Eliminate states that turned out not to be good.
The set of possibly good states is now smaller.
States from which goals are reachable in \( \leq 1 \) steps so that all immediate successors are possibly good.
Nested Fixpoint Algorithm

Example

States from which goals are reachable in \( \leq 2 \) steps so that all immediate successors are possibly good.
States from which goals are reachable in $\leq 3$ steps so that all immediate successors are possibly good.
States from which goals are reachable in \( \leq 4 \) steps so that all immediate successors are possibly good.
Remaining states are all good.
A further iteration would not eliminate more states.
Assign each state an operator so that the successor states are goal states or good, and some of them are closer to goal states. Use \textit{weak distances} computed with \textit{weak preimages}. For this example this is trivial.
Strong cyclic plans

Recall the definition of cyclic strong plans:

**Definition (strong cyclic plan)**

Let $S$ be the set of states of a planning task $\Pi$. Then a strong cyclic plan for $\Pi$ is a function $\pi : S_\pi \rightarrow O$ for some subset $S_\pi \subseteq S$ such that

- $\pi(s)$ is applicable in $s$ for all $s \in S_\pi$,
- $S_\pi(s_0) \subseteq S_\pi \cup S_\ast$ ($\pi$ is closed), and
- $S_\pi(s') \cap S_\ast \neq \emptyset$ for all $s' \in S_\pi(s_0)$ ($\pi$ is proper).
The procedure `prune` finds a maximal set of states for which reaching goals with looping is possible.

It consists of two nested loops:

1. The outer loop iterates through \( i = 0, 1, 2, \ldots \) and produces a shrinking sequence of candidate good state sets \( C_0, C_1, \ldots, C_n \) until \( C_i = C_{i+1} \).

2. The inner loop identifies growing sets \( W_j \) of states from which a goal state can be reached with \( j \) steps without leaving the current set of candidate good states \( C_i \). The union of all \( W_0, W_1, \ldots \) will be \( C_{i+1} \).
Procedure \textit{prune}

\textbf{Definition}

\begin{verbatim}
def prune(S, O, S_*):
    C_0 := S
    for each \(i \in \mathbb{N}_1:\)
        W_0 := S_*
        for each \(j \in \mathbb{N}_1:\)
            W_j := W_{j-1} \cup \bigcup_{o \in O}(\text{wpreimg}_o(W_{j-1}) \cap \text{spreimg}_o(C_{i-1}))
        if \(W_j = W_{j-1}:\)
            break
    C_i := W_j
    if \(C_i = C_{i-1}:\)
        return \langle C_i, \langle W_0, \ldots, W_{j-1} \rangle \rangle
\end{verbatim}
Procedure `prune`

Correctness

Lemma (Procedure `prune`)

Let $S$ and $S_\star \subseteq S$ be sets of states and $O$ a set of operators. Then $\text{prune}(S, O, S_\star)$ terminates after a finite number of steps and returns $C \subseteq S$ such that there is a strategy $\pi : C \setminus S_\star \rightarrow O$ that is a strong cyclic plan (for the states for which it is defined) and maximal in the sense that there is no set $C' \supseteq C$ and a strong cyclic plan $\pi' : C' \setminus S_\star \rightarrow O$.

- The sets $W_j$ also returned by `prune` encode weak distances and can be used to define the strong cyclic plan $\pi$. 
The planning algorithm

```
def strong-cyclic-plan(⟨V, I, O, γ⟩):
    S := set of states over V
    S_γ := {s ∈ S | s = γ}
    ⟨C, (W_j)_{j=0,1,2,...}⟩ = prune(S, O, S_γ)
    if I ∉ C:
        return no solution
    for each s ∈ C:
        δ(s) := min{j ∈ N_0 | s ∈ W_j}
    for each s ∈ C \ S_γ:
        π(s) := some operator o ∈ O with img_o(s) ⊆ C
        and min{δ(s') | s' ∈ img_o(s)} < δ(s)
    return π
```
The procedure *prune* runs in polynomial time in the number of states because the number of iterations of each loop is at most $n$ – hence there are $O(n^2)$ iterations – and computation on each iteration takes polynomial time in the number of states.

Finding strong cyclic plans for full observability is in the complexity class EXPTIME.

The problem is also EXPTIME-hard.

Similar to strong planning, we can speed up the algorithm in many practical cases by using a symbolic implementation (e.g. with BDDs).
Determinization-based Incremental Alg.

Idea [Kuter/Nau/Reisner/Goldman, 2008; Fu/Ng/Bastiani/Yen, 2011]:

1. Pretend the planning task was deterministic: Turn each action $o = \langle \chi, E \rangle$ with $E = \{e_1, \ldots, e_n\}$ into $n$ actions $o_i = \langle \chi, e_i \rangle$ for $i = 1, \ldots, n$. Obtain classical problem $\Pi'$.  

2. Find classical plan $P$ in $\Pi'$. Add state-action mapping corresponding to $P$ to $\pi$.  

3. For each operator $o_i$ used in $P$ (in state $s$), identify original nondeterministic operator $o$ and states $S' = img_o(s)$.  

4. For each “open” state $s' \in S'$, go to 2.

Remark: May require backtracking, if some state used in a classical plan turns out not to admit a strong cyclic plan.
Definition (all-outcomes determinization)

Let $\Pi = \langle V, I, O, \gamma \rangle$ be a nondeterministic planning task. The all-outcomes determinization of $\Pi$ is the deterministic planning task $\Pi_{\text{det}} = \langle V, I, O_{\text{det}}, \gamma \rangle$, where $O_{\text{det}} = \bigcup_{o \in O} o_{\text{det}}$, and $\langle \chi, E \rangle_{\text{det}} = \{ \langle \chi, e \rangle \mid e \in E \}$. 
List of states to solve: \( \{s_0\} \)
Plan for $s_0$ in determinization: $\text{blue}_2, \text{red}_2, \text{red}$
“Undesired” outcomes of blue\textsubscript{1} and red\textsubscript{1} lead to new list of states to solve: \{s\textsubscript{1}, s\textsubscript{2}\}
Determinization-based Incremental Alg.

Example

Plan for $s_1$ in determinization: $\text{red}, \text{blue}_2, \text{red}_2, \text{red}$
No new “undesired” outcomes.
List of states to solve: \( \{s_2\} \)
Plan for $s_2$ in determinization: $\textit{blue}_1$
“Undesired” outcome of $blue_2$ in $s_2$ leads to goal state, too. List of states to solve: $\emptyset$. Strong cyclic plan found.
Determinization-based Incremental Alg.

Pseudocode

Procedure incremental-strong-cyclic-plan

```python
def incremental-strong-cyclic-plan(⟨V, I, O, γ⟩):
    π ← ∅; fail ← {I}
    while fail ≠ ∅:
        s ← SELECTANDREMOVEFROM(fail)
        π' ← DETSEARCH(⟨V, s, O_{det}, γ⟩)
        if π' = FAILURE:
            if s = I: return FAILURE
            else: BACKTRACK(s, π, ⟨V, I, O, γ⟩)
        else:
            π ← π ∪ π'
        fail ← {s ∈ S | s nongoal state reachable from I following π, but π(s) undefined}
    return π
```
If a deterministic search fails, the state $s$ from which it started cannot be part of a strong cyclic plan.

- If $s = I$, the whole given planning problem is unsolvable and the algorithm returns FAILURE.
- Otherwise, state $s$, which has already been added to the constructed policy $\pi$, has to be removed from $\pi$, and the algorithm has to ensure that $s$ will never be reconsidered again. This is accomplished by the procedure BACKTRACK.
Determinization-based Incremental Alg.

**Procedure backtrack**

```python
def backtrack(s, π, ⟨V, I, O, γ⟩):
    update π by deleting all entries that would immediately lead to s, i.e. π ← π \ \{ (s', π(s')) | s ∈ img_{π(s')} (s') \}
    add all states s' removed from π to the set of fail-states fail
    permanently mark all formerly assigned actions π(s') removed from π at s' as inapplicable in s' to avoid running into the same dead end again.
```
Determinization-based Incremental Alg.

- Iteratively solves all-outcomes determinizations of $\Pi$ with “fail-states” as initial states.
- Planner can choose desired outcome of each action.
- Deterministic plans are added to policy under construction.
- Corresponding undesired outcomes have to be added to the set of “fail-states” $\text{fail}$.
- Deterministic plans for “fail-states” are constructed until no more “fail-states” remain.
- Eventually, the algorithm either returns a strong cyclic plan or FAILURE if no such plan exists.
Theorem

Procedure incremental-strong-cyclic-plan, called with task Π, returns a strong cyclic plan for Π iff such a plan exists, and FAILURE, otherwise.
- Can use any classical planner for deterministic searches.
- Can benefit from heuristics etc. used there.
- Classical planner can be configured to prefer short solutions or solutions using deterministic actions induced by nondeterministic actions with few different outcomes (likely fewer new “fail-states”).
When to terminate a deterministic sub-search?
- At goal states?
- At states currently part of the partial solution?
- At parent of currently solved “fail-state”?

This can make a huge difference.

Similarly: Where should the heuristic guide the classical planner? Goals, partial solution, parent node?

Additional marking of nodes as definitely solved if this can be detected.

State reuse between subsequent classical planner calls.

Generalization of solved states by regression search from goal along weak (deterministic) plan (cf. [Muise/McIlraith/Beck, 2012]).
Maintenance goals
In this lecture, we usually limit ourselves to the problem of finding plans that reach a goal state.

In practice, planning is often about more general goals, where execution cannot be terminated.

1. An animal: find food, eat, sleep, find food, eat, sleep, ... 
2. Cleaning robot: keep the building clean.

These problems cannot be directly formalized in terms of reachability because infinite (unbounded) plan execution is needed.

We do not discuss this topic in full detail. However, to give at least a little impression of planning for temporally extended goals, we will discuss the simplest objective with infinite plan executions: maintenance.
Plan objectives
Maintenance

**Definition**

Let $\mathcal{T} = \langle V, I, O, \gamma \rangle$ be a planning task with state set $S$ and set of goal states $S_\star = \{ s \in S | s \models \gamma \}$. A strategy $\pi$ for $\mathcal{T}$ is called a plan for maintenance for $\mathcal{T}$ iff

- $\pi(s)$ is applicable in $s$ for all $s \in S_\pi$,
- $S_\pi(s_0) \subseteq S_\pi$, and
- $S_\pi(s_0) \subseteq S_\star$. 
The state of an animal is determined by three state values: hunger (0, 1, 2), thirst (0, 1, 2) and location (river, pasture, desert). There is also a special state called death.

- Thirst grows when not at river; at river it is 0.
- Hunger grows when not on pasture; on pasture it is 0.
- If hunger or thirst exceeds 2, the animal dies.
- The goal of the animal is to avoid death.
We can infer rules backwards starting from the death condition.

1. If in desert and $\text{thirst} = 2$, must go to river.
2. If in desert and $\text{hunger} = 2$, must go to pasture.
3. If on pasture and $\text{thirst} = 1$, must go to desert.
4. If at river and $\text{hunger} = 1$, must go to desert.

If the above rules conflict, the animal will die.
Algorithm for maintenance goals

Idea

Summary of the algorithm idea

Repeatedly eliminate from consideration those states that in one or more steps unavoidably lead to a non-goal state.

- A state is \(i\)-safe iff there is a plan that guarantees “survival” for the next \(i\) actions.
- A state is safe (or \(\infty\)-safe) iff it is \(i\)-safe for all \(i \in \mathbb{N}_0\).
- The 0-safe states are exactly the goal states: maintenance objective is satisfied for the current state.
- Given all \(i\)-safe states, compute all \(i + 1\)-safe states by using strong preimages.
- For some \(i \in \mathbb{N}_0\), \(i\)-safe states equal \(i + 1\)-safe states because there are only finitely many states and at each step and \(i + 1\)-safe states are a subset of \(i\)-safe states. Then \(i\)-safe states are also \(\infty\)-safe.
def maintenance-plan(⟨V, I, O, γ⟩):
    S := set of states over V
    Safe0 := {s ∈ S | s |= γ}
    for each i ∈ N:
        Safei := Safei−1 ∩ ∪o∈O spreimgo(Safei−1)
        if Safei = Safei−1:
            break
        if I ∉ Safei:
            return no solution
    for each s ∈ Safei:
        π(s) := some operator o ∈ O with imgo(s) ⊆ Safei
    return π
Maintenance goals

Transition system for the example
Maintenance goals

0-safe states
Maintenance goals

1-safe states

pasture

river
Maintenance goals

\( i \)-safe states for all \( i \geq 2 \)
Summary
Different planning objectives

Strong planning

Strong cyclic planning

Maintenance
We have considered different classes of solutions for planning tasks by defining different planning problems.
- strong planning problem: find a strong plan
- strong cyclic planning problem: find a strong cyclic plan
- ...

Alternatively, we could allow specifying goals in a modal logic like computational tree logic to directly express the type of plan we are interested in using modalities such as A (all), E (exists), G (globally), and F (finally).
- Weak planning: $\text{EF}\varphi$
- Strong planning: $\text{AF}\varphi$
- Strong cyclic planning: $\text{AGEF}\varphi$
- Maintenance: $\text{AG}\varphi$
We have extended our earlier planning algorithm from strong plans to strong cyclic plans.

The story does not end there: When considering infinitely executing plans, many more types of goals are feasible.

We considered maintenance as a simple example of a temporally extended goal.

In general, temporally extended goals be expressed in modal logics such as computational tree logic (CTL).

We presented dynamic programming (backward search) algorithms for strong cyclic and maintenance planning.

In practice, one might implement both algorithms by using binary decision diagrams (BDDs) as a data structure for state sets.