Principles of AI Planning

6. Planning as search: search algorithms

Albert-Ludwigs-Universität Freiburg

UNI

Bernhard Nebel and Robert Mattmüller

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Introduction to search algorithms for planning

Introduction

Nodes and states Search for plannin Common procedures

Uninforme search

Heuristic search

Our plan for the next lectures



Choices to make:

- search direction: progression/regression/both → previous chapter

- search control: heuristics, pruning techniques
 → next chapters

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- Search algorithms are used to find solutions (plans) for transition systems in general, not just for planning tasks.
- Planning is one application of search among many.
- In this chapter, we describe some popular and/or representative search algorithms, and (the basics of) how they apply to planning.
- Most of this is review of material that should be known (details: Russell and Norvig's textbook).

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Search states vs. search nodes





In search, one distinguishes:

- search states s → states (vertices) of the transition system
- \blacksquare search nodes $\sigma \leadsto$ search states plus information on where/when/how they are encountered during search

What is in a search node?

Different search algorithms store different information in a search node σ , but typical information includes:

- \blacksquare state(σ): associated search state
- \blacksquare parent(σ): pointer to search node from which σ is reached
- **action**(σ): action leading from $state(parent(\sigma))$ to $state(\sigma)$
- $g(\sigma)$: cost of σ (length of path from the root node)

For the root node, $parent(\sigma)$ and $action(\sigma)$ are undefined.

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Search states vs. planning states



Search states \neq (planning) states:

- Search states don't have to correspond to states in the planning sense.
 - progression: search states ≈ (planning) states
 - lacktriangledown regression: search states pprox sets of states (formulae)
- Search algorithms for planning where search states are planning states are called state-space search algorithms.
- Strictly speaking, regression is not an example of state-space search, although the term is often used loosely.
- However, we will put the emphasis on progression, which is almost always state-space search.

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Required ingredients for search





A general search algorithm can be applied to any transition system for which we can define the following three operations:

- init(): generate the initial state
- \blacksquare is-goal(s): test if a given state is a goal state
- succ(s): generate the set of successor states of state s, along with the operators through which they are reached (represented as pairs $\langle o, s' \rangle$ of operators and states)

Together, these three functions form a search space (a very similar notion to a transition system).

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Search for planning: progression





Let $\Pi = \langle A, I, O, \gamma \rangle$ be a planning task.

Search space for progression search

states: all states of Π (assignments to A)

$$lacksquare$$
 succ $(s) = \{\langle o, s' \rangle \mid \text{applicable } o \in O, s' = app_o(s)\}$

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Search for planning: regression



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Let $\Pi = \langle A, I, O, \gamma \rangle$ be a planning task.

Search space for regression search

states: all formulae over A (how many?)

■ init() =
$$\gamma$$

■ is-goal
$$(\varphi) = \begin{cases} \text{true} & \text{if } I \models \varphi \\ \text{false} & \text{otherwise} \end{cases}$$

■
$$\operatorname{succ}(\varphi) = \{\langle o, \varphi' \rangle \mid o \in O, \\ \varphi' = \operatorname{regr}_o(\varphi), \varphi' \text{ is satisfiable}\}$$
 (modified if splitting is used)

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Classification of search algorithms



uninformed search vs. heuristic search:

- uninformed search algorithms only use the basic ingredients for general search algorithms
- heuristic search algorithms additionally use heuristic functions which estimate how close a node is to the goal

systematic search vs. local search:

- systematic algorithms consider a large number of search nodes simultaneously
- local search algorithms work with one (or a few) candidate solutions (search nodes) at a time
- not a black-and-white distinction; there are crossbreeds
 (e.g., enforced hill-climbing)

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uninformed vs. heuristic search:

- For satisficing planning, heuristic search vastly outperforms uninformed algorithms on most domains.
- For optimal planning, the difference is less pronounced.

systematic search vs. local search:

- For satisficing planning, the most successful algorithms are somewhere between the two extremes.
- For optimal planning, systematic algorithms are required.

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Common procedures for search algorithms



Before we describe the different search algorithms, we introduce three procedures used by all of them:

- make-root-node: Create a search node without parent.
- make-node: Create a search node for a state generated as the successor of another state.
- extract-solution: Extract a solution from a search node representing a goal state.

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Procedure make-root-node





make-root-node: Create a search node without parent.

Procedure make-root-node

```
def make-root-node(s):
\sigma := \mathbf{new} \text{ node}
state(\sigma) := s
parent(\sigma) := \text{ undefined}
action(\sigma) := \text{ undefined}
g(\sigma) := 0
\mathbf{return} \ \sigma
```

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Procedure make-node





make-node: Create a search node for a state generated as the successor of another state.

Procedure make-node

```
def make-node(\sigma, o, s):
\sigma' := \mathbf{new} \text{ node}
state(\sigma') := s
parent(\sigma') := \sigma
action(\sigma') := o
g(\sigma') := g(\sigma) + 1
\mathbf{return} \ \sigma'
```

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Procedure extract-solution



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extract-solution: Extract a solution from a search node representing a goal state.

Procedure extract-solution

```
\begin{aligned} \textbf{def} & \text{ extract-solution}(\sigma): \\ & \textit{ solution} := \textbf{new} \text{ list} \\ & \textbf{while } \textit{ parent}(\sigma) \text{ is defined:} \\ & \textit{ solution}. \text{push-front}(\textit{action}(\sigma)) \\ & \sigma := \textit{ parent}(\sigma) \\ & \textbf{return } \textit{ solution} \end{aligned}
```

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Uninformed search

duplicate detect Breadth-first wit duplicate detect Random walk

Heuristic search

Summary

Uninformed search algorithms

Uninformed search algorithms

- UNI
- Uninformed algorithms are less relevant for planning than heuristic ones, so we keep their discussion brief.
- Uninformed algorithms are mostly interesting to us because we can compare and contrast them to related heuristic search algorithms.

Popular uninformed systematic search algorithms:

- breadth-first search
- depth-first search
- iterated depth-first search

Popular uninformed local search algorithms:

random walk

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Breadth-first w/o duplicate detection Breadth-first with duplicate detection

Heuristic

Breadth-first search

```
queue := \mathbf{new} \text{ fifo-queue} queue. \mathsf{push-back}(\mathsf{make-root-node}(\mathsf{init}())) \mathsf{while} \ \mathsf{not} \ queue. \mathsf{empty}(): \sigma = queue. \mathsf{pop-front}() \mathsf{if} \ \mathsf{is-goal}(\mathsf{state}(\sigma)): \mathsf{return} \ \mathsf{extract-solution}(\sigma) \mathsf{for} \ \mathsf{each} \ \langle o, s \rangle \in \mathsf{succ}(\mathit{state}(\sigma)): \sigma' := \mathsf{make-node}(\sigma, o, s) queue. \mathsf{push-back}(\sigma') \mathsf{return} \ \mathsf{unsolvable}
```

- Possible improvement: duplicate detection (see next slide).
- Another possible improvement: test if σ' is a goal node; if so, terminate immediately. (We don't do this because it obscures the similarity to some of the later algorithms.)

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Uninformed search Breadth-first w/o

duplicate detection

Breadth-first with

Random walk

search





Breadth-first search with duplicate detection

```
queue := new fifo-queue
queue.push-back(make-root-node(init()))
closed := \emptyset
while not queue.empty():
      \sigma = queue.pop-front()
     if state(\sigma) \notin closed:
           closed := closed \cup \{state(\sigma)\}\
           if is-goal(state(\sigma)):
                 return extract-solution(\sigma)
           for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                 \sigma' := \mathsf{make-node}(\sigma, o, s)
                 queue.push-back(\sigma')
return unsolvable
```

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Breadth-first w/c duplicate detect

> Breadth-first with duplicate detection Random walk

Heuristic





Breadth-first search with duplicate detection

```
queue := new fifo-queue
queue.push-back(make-root-node(init()))
closed := \emptyset
while not queue.empty():
      \sigma = queue.pop-front()
     if state(\sigma) \notin closed:
           closed := closed \cup \{state(\sigma)\}\
           if is-goal(state(\sigma)):
                 return extract-solution(\sigma)
           for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                 \sigma' := \mathsf{make-node}(\sigma, o, s)
                 queue.push-back(\sigma')
return unsolvable
```

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Random walk



Random walk

```
\begin{split} \sigma := & \mathsf{make\text{-}root\text{-}node}(\mathsf{init}()) \\ & \mathsf{forever} \colon \\ & \mathsf{if} \ \mathsf{is\text{-}goal}(\mathsf{state}(\sigma)) \colon \\ & \mathsf{return} \ \mathsf{extract\text{-}solution}(\sigma) \\ & \mathsf{Choose} \ \mathsf{a} \ \mathsf{random} \ \mathsf{element} \ \langle o, s \rangle \ \mathsf{from} \ \mathsf{succ}(\mathsf{state}(\sigma)) . \\ & \sigma := & \mathsf{make\text{-}node}(\sigma, o, s) \end{split}
```

- The algorithm usually does not find any solutions, unless almost every sequence of actions is a plan.
- Often, it runs indefinitely without making progress.
- It can also fail by reaching a dead end, a state with no successors. This is a weakness of many local search approaches.

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duplicate detection Breadth-first with duplicate detection Bandom walk

leuristic

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Heuristic search algorithms

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Heuristic search algorithms: systematic



 Heuristic search algorithms are the most common and overall most successful algorithms for classical planning.

Popular systematic heuristic search algorithms:

- greedy best-first search
- A*
- weighted A*
- IDA*
- depth-first branch-and-bound search
- ...

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Heuristic search algorithms: local



 Heuristic search algorithms are the most common and overall most successful algorithms for classical planning.

Popular heuristic local search algorithms:

- hill-climbing
- enforced hill-climbing
- beam search
- tabu search
- genetic algorithms
- simulated annealing
- **.** . . .

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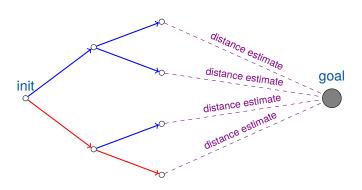
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Heuristic search: idea







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Required ingredients for heuristic search



A heuristic search algorithm requires one more operation in addition to the definition of a search space.

Definition (heuristic function)

Let Σ be the set of nodes of a given search space.

A heuristic function or heuristic (for that search space) is a function $h: \Sigma \to \mathbb{N}_0 \cup \{\infty\}$.

The value $h(\sigma)$ is called the heuristic estimate or heuristic value of heuristic h for node σ . It is supposed to estimate the distance from σ to the nearest goal node.

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What exactly is a heuristic estimate?



What does it mean that *h* "estimates the goal distance"?

- For most heuristic search algorithms, *h* does not need to have any strong properties for the algorithm to work (= be correct and complete).
- However, the efficiency of the algorithm closely relates to how accurately *h* reflects the actual goal distance.
- For some algorithms, like A*, we can prove strong formal relationships between properties of *h* and properties of the algorithm (optimality, dominance, run-time for bounded error, ...)
- For other search algorithms, "it works well in practice" is often as good an analysis as one gets.

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Heuristics Systematic or

Local search

Heuristics applied to nodes or states?



- Most texts apply heuristic functions to states, not nodes.
- This is slightly less general than our definition:
 - Given a state heuristic h, we can define an equivalent node heuristic as $h'(\sigma) := h(state(\sigma))$.
 - The opposite is not possible. (Why not?)
- There is good justification for only allowing state-defined heuristics: why should the estimated distance to the goal depend on how we ended up in a given state s?
- We call heuristics which don't just depend on $state(\sigma)$ pseudo-heuristics.
- In practice there are sometimes good reasons to have the heuristic value depend on the generating path of σ (e.g., landmark pseudo-heuristic, Richter et al. 2008).

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Perfect heuristic



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minoduction

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Summary

Let Σ be the set of nodes of a given search space.

Definition (optimal/perfect heuristic)

The optimal or perfect heuristic of a search space is the heuristic h^* which maps each search node σ to the length of a shortest path from $state(\sigma)$ to any goal state.

Note: $h^*(\sigma) = \infty$ iff no goal state is reachable from σ .

Properties of heuristics



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Summary

A heuristic h is called

- safe if $h^*(\sigma) = \infty$ for all $\sigma \in \Sigma$ with $h(\sigma) = \infty$
- **goal-aware** if $h(\sigma) = 0$ for all goal nodes $\sigma \in \Sigma$
- admissible if $h(\sigma) \le h^*(\sigma)$ for all nodes $\sigma \in \Sigma$
- consistent if $h(\sigma) \le h(\sigma') + 1$ for all nodes $\sigma, \sigma' \in \Sigma$ such that σ' is a successor of σ

Relationships?

```
open := new min-heap ordered by (\sigma \mapsto h(\sigma))
open.insert(make-root-node(init()))
closed := \emptyset
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed:
            closed := closed \cup \{state(\sigma)\}\
            if is-goal(state(\sigma)):
                  return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                  \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < \infty:
                        open.insert(\sigma')
return unsolvable
```

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Summarv

Properties of greedy best-first search



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- Summary

- one of the three most commonly used algorithms for satisficing planning
- complete for safe heuristics (due to duplicate detection)
- suboptimal unless h satisfies some very strong assumptions (similar to being perfect)
- invariant under all strictly monotonic transformations of h
 (e.g., scaling with a positive constant or adding a
 constant)



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A* (with duplicate detection and reopening)

```
open := new min-heap ordered by (\sigma \mapsto g(\sigma) + h(\sigma))
open.insert(make-root-node(init()))
closed := \emptyset
distance := \emptyset
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed or g(\sigma) < distance(state(\sigma)):
            closed := closed \cup \{state(\sigma)\}\
            distance(state(\sigma)) := g(\sigma)
            if is-goal(state(\sigma)):
                  return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                  \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < \infty: open.insert(\sigma')
return unsolvable
```

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Example





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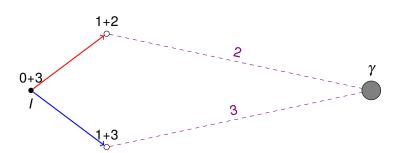
Systematic search



Example







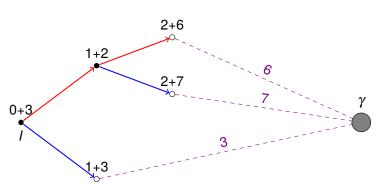
search

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Example





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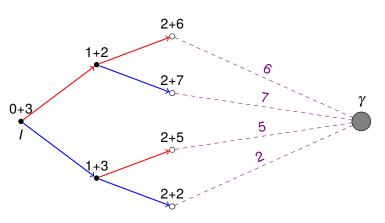
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Example





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Example



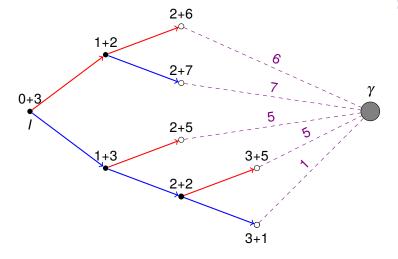




search

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Terminology for A*



) II

- f value of a node: defined by $f(\sigma) := g(\sigma) + h(\sigma)$
- generated nodes: nodes inserted into open at some point
- **expanded nodes:** nodes σ popped from *open* for which the test against *closed* and *distance* succeeds
- reexpanded nodes: expanded nodes for which $state(\sigma) \in closed$ upon expansion (also called reopened nodes)

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Properties of A*



- the most commonly used algorithm for optimal planning
- rarely used for satisficing planning
- complete for safe heuristics (even without duplicate detection)
- optimal if h is admissible (even without duplicate detection)
- never reopens nodes if h is consistent

Implementation notes:

- in the heap-ordering procedure, it is considered a good idea to break ties in favour of lower h values
- can simplify algorithm if we know that we only have to deal with consistent heuristics
- common, hard to spot bug: test membership in closed at the wrong time

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```
open := new min-heap ordered by (\sigma \mapsto g(\sigma) + W \cdot h(\sigma))
open.insert(make-root-node(init()))
closed := \emptyset
distance := \emptyset
while not open.empty():
      \sigma = open.pop-min()
      if state(\sigma) \notin closed or g(\sigma) < distance(state(\sigma)):
            closed := closed \cup \{state(\sigma)\}\
            distance(\sigma) := g(\sigma)
            if is-goal(state(\sigma)):
                  return extract-solution(\sigma)
            for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                  \sigma' := \mathsf{make-node}(\sigma, o, s)
                  if h(\sigma') < \infty: open.insert(\sigma')
return unsolvable
```

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Properties of weighted A*



The weight $W \in \mathbb{R}_0^+$ is a parameter of the algorithm.

- \blacksquare for W = 0, behaves like breadth-first search
- \blacksquare for W = 1, behaves like A^*
- lacksquare for $W o \infty$, behaves like greedy best-first search

Properties:

- one of the most commonly used algorithms for satisficing planning
- for W > 1, can prove similar properties to A*, replacing optimal with bounded suboptimal: generated solutions are at most a factor W as long as optimal ones

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Hill-climbing

- can easily get stuck in local minima where immediate improvements of $h(\sigma)$ are not possible
- many variations: tie-breaking strategies, restarts

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Enforced hill-climbing: procedure improve

```
def improve(\sigma_0):
      queue := new fifo-queue
      queue.push-back(\sigma_0)
      closed := \emptyset
      while not queue.empty():
            \sigma = queue.pop-front()
            if state(\sigma) \notin closed:
                  closed := closed \cup \{state(\sigma)\}\
                   if h(\sigma) < h(\sigma_0):
                         return \sigma
                  for each \langle o, s \rangle \in \text{succ}(state(\sigma)):
                         \sigma' := \mathsf{make-node}(\sigma, o, s)
                         queue.push-back(\sigma')
      fail
```

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Summary

 \leadsto breadth-first search for more promising node than σ_0





Enforced hill-climbing

```
\sigma := \mathsf{make\text{-}root\text{-}node}(\mathsf{init}())

while not is-goal(state(\sigma)):

\sigma := \mathsf{improve}(\sigma)

return extract-solution(\sigma)
```

- one of the three most commonly used algorithms for satisficing planning
- \blacksquare can fail if procedure improve fails (when the goal is unreachable from σ_0)
- complete for undirected search spaces (where the successor relation is symmetric) if $h(\sigma) = 0$ for all goal nodes and only for goal nodes

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Local search





- distinguish: planning states, search states, search nodes
 - planning state: situation in the world modelled by the task
 - search state: subproblem remaining to be solved
 - In state-space search (usually progression search), planning states and search states are identical.
 - In regression search, search states usually describe sets of states ("subgoals").
 - search node: search state + info on "how we got there"
- search algorithms mainly differ in order of node expansion
 - uninformed vs. informed (heuristic) search
 - local vs. systematic search

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Summary (ctd.)



- heuristics: estimators for "distance to goal node"
 - usually: the more accurate, the better performance
 - desiderata: safe, goal-aware, admissible, consistent
 - the ideal: perfect heuristic h*
- most common algorithms for satisficing planning:
 - greedy best-first search
 - weighted A*
 - enforced hill-climbing
- most common algorithm for optimal planning:
 - A*

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