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# Logik für Informatiker: PROLOG Part 1: Introduction

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(original slides by Peter Flach)

### Prolog

- PROgramming in LOGic
- <u>PRO</u>gramming <u>Language</u> <u>Of</u> <u>G</u>od ©
- Why Prolog ?
  - completely different programming paradigm
  - a declarative programming language (as Haskell)
    - focus on WHAT instead of HOW
  - based on first order logic

used especially in artificial intelligence, natural language processing, search problems, expert systems, databases, ...

very elegant / powerful / compact

## What we shall do

- short introduction to logic
- Prolog as a programming language
- various programming techniques
- examples from artificial intelligence
- a bit of theory / a lot of practice
- some larger programs

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## Materials

 Simply Logical" by Peter Flach, Addison-Wiley, 1994. (printed on demand) – 1<sup>st</sup> few chapters (free download: http://www.cs.bris.ac.uk/~flach/SimplyLogical.html)

PROLOG. Programming for Artificial Intelligence, by Ivan Bratko, Addison-Wesley, Third Edition 2001, next chapters ... (Old german version might be available)

Prolog:

- YAP Prolog (http://sourceforge.net/projects/yap/ or http://www.ncc.up.pt/~vsc/Yap/ )
- SWI Prolog (http://www.swi-prolog.org/)
- Many more: GNU Prolog, Visual Prolog, …



London Underground example

connected(bond\_street,oxford\_circus,central). connected(oxford\_circus,tottenham\_court\_road,central). connected(bond\_street,green\_park,jubilee). connected(green\_park,charing\_cross,jubilee). connected(green\_park,piccadilly\_circus,piccadilly). connected(piccadilly\_circus,leicester\_square,piccadilly). connected(green\_park,oxford\_circus,victoria). connected(oxford\_circus,piccadilly\_circus,bakerloo). connected(piccadilly\_circus,charing\_cross,bakerloo). connected(tottenham\_court\_road,leicester\_square,northern). connected(leicester\_square,charing\_cross,northern).

# London Underground in Prolog (1)

Two stations are nearby if they are on the same line with at most one other station in between:

```
nearby(bond_street,oxford_circus).
nearby(oxford_circus,tottenham_court_road).
nearby(bond_street,tottenham_court_road).
nearby(bond_street,green_park).
nearby(green_park,charing_cross).
nearby(bond_street,charing_cross).
nearby(green_park,piccadilly_circus).
```

#### or better

```
nearby(X,Y):-connected(X,Y,L).
nearby(X,Y):-connected(X,Z,L),connected(Z,Y,L).
```

Facts: unconditional truths Rules/Clauses: conditional truths

Both definitions are equivalent.

London Underground in Prolog (2)

- Query: which station is nearby Tottenham Court Road?
  - ?- nearby(tottenham\_court\_road, W).
- Prefix ?- means it's a query and not a fact.
- Answer to query is:
   {W -> leicester\_square}
   a so-called substitution.
- When nearby defined by facts, substitution found by simple *matching*.

Answering queries (1)

## Exercise 1.1:

Define a predicate not\_too\_far, which is true, if two stations are one the same or a different line, with at most one station in between.

Exercise 1.1

#### Remember nearby:

```
nearby(X,Y):-connected(X,Y,L).
nearby(X,Y):-connected(X,Z,L),connected(Z,Y,L).
```

Then not\_too\_far is:

```
not_too_far(X,Y):-connected(X,Y,L).
not_too_far(X,Y):-connected(X,Z,L1),connected(Z,Y,L2).
```

This can be rewritten with don't cares:

not\_too\_far(X,Y):-connected(X,Y,\_).
not\_too\_far(X,Y):-connected(X,Z,\_),connected(Z,Y,\_).

If clauses are involved, then answering a query can take several steps.

```
    ?- nearby(tottenham_court_road, W).
    matches conclusion (head) of clause
        nearby(X,Y) :- connected(X,Y,L).
    with the substitution
        {X -> tottenham court road, Y -> W}
```

Subsequently,

?- connected(tottenham\_court\_road, W, L).
is to prove.

• Here, looking up the facts is sufficient for answering the query: {W -> leicester\_square, L-> northern}

```
Result:
```

```
{W -> leicester_square}
```

## Answering queries (2)



Resolution

#### To answer a query

?- Q1, Q2, ..., Qn. find a clause A :- B1,..., Bm such that A matches Q1, and then answer the query ?- B1,..., Bm, Q2,..., Qn.

- Adds a procedural interpretation to the declarative interpretation of a logical formula
- Resolution proof: reductio ad absurdum; proof by refutation
- Start: clause with empty head (conclusion), e.g.:
  - :- nearby(tottenham\_court\_road, W).

(= negation of nearby(...))

 Contradiction is found, if empty clause is derived. Empty clause: premise (body) is always true, because non-existing.

Recursion (1)

- Up to now: rules (clauses) and facts
- Particular kind of rules; rules that are defined by recurring to themselves: recursion

```
IF N = 0 THEN FAC:= 1
ELSE FAC:= N*FAC(N-1)
```

- Recursion is (except for *failure-driven loops*) the only construct for loops in Prolog.
- Example relation reachable Could be defined by enumeration of facts or by nonrecursive rules for routes of length 1, 2, etc.

A station is reachable from another if they are on the same line, or with one, two, ... changes:

#### or better

Recursion (2)

...

```
reachable(X,Y):-connected(X,Y,L).
reachable(X,Y):-connected(X,Z,L),reachable(Z,Y).
```

Recursive definition:

- Examples so far have shown:
   Prolog performs *search* in order to answer queries
- Backtracking: returning to previous choice points, if proof fails at some point.

Simply Logical – Chapterig. 1.3, p.9



Recursion (4)

```
reachable0(X,Y):-
    connected(X,Y,L).
reachable1(X,Y,Z):-
    connected(X,Z,L1),
    connected(Z,Y,L2).
reachable2(X,Y,Z1,Z2):-
    connected(X,Z1,L1),
    connected(Z1,Z2,L2),
    connected(Z2,Y,L3).
```

One clause for each route of length n.

Solution: functors

Are used to construct complex objects out of simpler ones.

e.g., route(oxford\_circus, tottenham\_court\_road)

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Structured terms (1)





Structured terms (2)

- Built-in data type in Prolog
- Functor: "."
- Tree notation with functor as well as linear notation possible
- Empty list: []
- •.(a, .(b, .(c, [])))
- Linear: [a, b, c]
- •.(First, Rest)
- [First|Rest]

Lists (1)

[First,Second,Third|Rest]

Lists (2)



This list can be written in many ways:

- . (a,.(b,.(c,[])))
- [a|[b|[c|[]]]
- [a|[b|[c]]]
- [a|[b,c]]
- [a,b,c]

•

**[a,b|[**C]]





```
Definition of predicate append:
```

```
append (X, Y, Z) is true, if appending lists x and y gives list z:
append([], X, X).
append([X|R], Y, [X|Z]) := append(R, Y, Z).
?- append([1,2], [3,4], [1,2,3,4]).
yes
?- append([1,2], B, [1,2,3]).
B = [3]
?- append(A, [2,3], [1,2,3]).
\mathbf{A} = [1]
?- append(A, B, [1,2]).
A = [], B = [1,2];
A = [1], B = [2];
A = [1, 2], B = []
```

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# Simple list processing

- Prolog has very simple syntax
  - constants, variables, and structured terms refer to objects
    - variables start with uppercase character
    - functors are never evaluated, but are used for naming
  - predicates express relations between objects
  - clauses express true statements
    - each clause independent of other clauses
- Queries are answered by matching with head of clause
  - there may be more than one matching clause
    - query answering is search process
  - query may have 0, 1, or several answers
  - no pre-determined input/output pattern (usually)

# Summary