Principles of Al Planning Nondeterministic planning with partial observability

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Nondeterministic planning with partial observability

Planning with partial observability is harder than both the fully observable and unobservable cases:

- Memoryless plans (where the next action to take only depends on the current situation) as in the fully observable case are not sufficient.
 - Of course, we cannot define a memoryless plan based on individual states because limited observability makes some states indistinguishable.
 - It is also not sufficient to consider memoryless plans where the action to take is based on the current observation class.
- Conformant (i.e., non-branching) plans as in the unobservable case are also clearly not powerful enough.

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Strong planning

- We will (mostly) consider the strong planning problem.
- Generalizations to the strong cyclic planning are similar to the fully observable case.

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Algorithms

Similar to other variants of the planning problem, there are three major approaches to nondeterministic planning with partial observability:

- Reduction to another problem
- Forward search
- Backward search

We will consider one example for each of these.

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Algorithms Three approaches

Reduction to another problem:

Reduce to planning with full observability.

Forward search (progression):

- Define the search space as an AND/OR tree.
- Define a heuristic function for such trees.
- Use a tree search algorithm such as AO* or Proof Number Search.

Backward search (regression):

- Start from the set of goal states.
- Find state sets from which already generated state sets can be reached by applying operators and making observations.

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- Memoryless plans are not sufficient for the partially observable case because a plan must take into account the knowledge collected in previous observations etc.
- During plan execution, this knowledge is represented in the current belief state.
- One idea for solving a partially observable task \mathcal{T} is to map it to a fully observable task \mathcal{T}' where each belief state of \mathcal{T} corresponds to a state of \mathcal{T}' .

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Reduction to fully observable case State variables

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

State variables

- For each state $s \in S$, there is one state variable $v_s \in A'$.
- Intuition: v_s is true in a state of \mathcal{T}' iff it is possible that we are currently in s.
- Formally: $A' := \{ v_s \mid s \in S \}$

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Initial state formula

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

Initial state formula

- The initial state of \mathcal{T}' is fully deterministic (in terms of A'), as there is only one possible initial belief state in \mathcal{T} .
- For all states s in the initial belief state of \mathcal{T} , variable v_s is initially true. Other variables are initially false.
- Formally: $I' := \bigwedge_{s \in S, s \models I} v_s \wedge \bigwedge_{s \in S, s \not\models I} \neg v_s$.

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Initial state formula

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

Goal formula

- ullet A goal belief state of ${\mathcal T}$ is one where all possible states satisfy G.
- This is equivalent to saying that no state in the current belief state violates G.
- We can express that by saying that none of the variables v_s for states s violating G are true.
- Formally: $G' := \bigwedge_{s \in S, s \not\models G} \neg v_s$.

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Initial state formula

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

Operators (preconditions)

- Each operator $o = \langle c, e \rangle \in O$ is translated to an operator $o' = \langle c', e' \rangle \in O'$.
- ullet To test whether operator o is applicable, we must verify that all states in the current belief state satisfy c.
- Again, this is equivalent to saying that no state in the current belief state violates c.
- Formally: $c' := \bigwedge_{s \in S, s \not\models c} \neg v_s$.

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Initial state formula

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

Operators (effects)

- Each operator $o=\langle c,e\rangle\in O$ is translated to an operator $o'=\langle c',e'\rangle\in O'.$
- After applying operator o, we can possibly be in state $s \in S$ iff we were previously in some state in which o is applicable and from which applying o can lead to s.
- This is modeled by an effect $((\bigvee_{t \in \mathit{preimg}_o(s)} v_t) \rhd v_s) \land (\neg(\bigvee_{t \in \mathit{preimg}_o(s)} v_t) \rhd \neg v_s).$
- $\begin{array}{ll} \bullet \text{ Formally:} & e' := \bigwedge_{s \in S} (((\bigvee_{t \in \mathit{preimg}_o(s)} v_t) \vartriangleright v_s) \land \\ & (\neg (\bigvee_{t \in \mathit{preimg}_o(s)} v_t) \vartriangleright \neg v_s)). \end{array}$

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 We have translated state variables, initial state formula, goal formula and operators.

Is that it?

- So far, our translation is independent of the set of observable variables V!
- Moreover, the resulting planning task is deterministic!

Is there an error in our modeling?

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Is there an error in our modeling?

- No, but it is not complete yet: There are solvable partially observable tasks \mathcal{T} for which \mathcal{T}' (as defined so far) is unsolvable.
- The reason for this is that he have not yet modeled the possibility of observing state variables.

Modeling observations requires introducing nondeterminism in \mathcal{T}' .

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Reduction to fully observable case Observations

Let $\mathcal{T} = \langle A, I, O, G, V \rangle$ be the input task with state set S. We define the fully observable task $\mathcal{T}' = \langle A', I', O', G', A' \rangle$.

Observations

- In general, our formalism allows observations to be general formulas over V. However, it is sufficient to only consider atomic observations $u \in V$.
- If we observe u in a belief state b, we can end up in two different belief states: one containing exactly the states of b where u is true, and one containing exactly the states of b where u is false.
- In other words, either the belief states where u is false or the belief states where u is true are ruled out.
- Formally: Translate observation of $u \in V$ into an operator $\langle \top, e'_u \rangle \in O'$ with $e'_u := (\bigwedge_{s \in S, s \not\models u} \neg v_s) | (\bigwedge_{s \in S, s \models u} \neg v_s).$

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Reduction to fully observable case Discussion

- Note that the reduction works both for strong and for strong cyclic planning.
- The reduction has a significant drawback: Since it introduces as many state variables as there are states in the original task, the resulting problem is exponentially larger than the original one.
- This will usually not be practical.
- On the other hand, there does not really exist any truly "practical" algorithm for nondeterministic planning with partial observability.

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Reduction to fully observable case Complexity result

- Using an exponential-time planning algorithm for fully observable planning, \mathcal{T}' can be solved in time $O(c^{\|\mathcal{T}'\|})$, and $\|\mathcal{T}'\| = O(c^{\|\mathcal{T}\|})$.
- Thus, we have a double-exponential $(O(c^{c^{||\mathcal{T}||}}))$ algorithm for nondeterministic planning for partial observability.
- We will later prove that this is worst-case optimal.

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Search in AND/OR trees

In forward search, plans are represented as trees whose nodes represent the situations arising during plan execution.

- The root node represents the initial situation.
- OR nodes correspond to choosing and applying operators.
 - Note how these relate to operators in T' in the earlier reduction.
- AND nodes correspond to making observations.
 - Note how these relate to nondeterminism in \mathcal{T}' in the earlier reduction.

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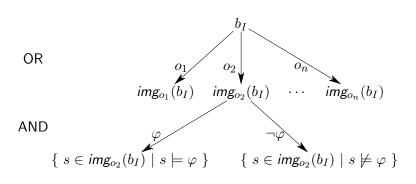
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Summary

Search in AND/OR trees $_{\text{Example}}$



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AND/OR trees Formal definition

Definition

An AND/OR tree is a labeled rooted tree where

- internal nodes are labeled with (∧) or (∨)
 (AND nodes/OR nodes), and
- leaves are labeled with (⊤) or (⊥) (true leaves/false leaves).

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AND/OR trees Truth value

Definition

An AND/OR tree evaluates to true iff

- it is a true leaf.
- it is an OR node with a child that evaluates to true, or
- it is an AND node whose children all evaluate to true.

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Partial plan trees

Definition

A partial plan tree for a nondeterministic planning task $\langle A,I,O,G,V\rangle$ with state set S is an AND/OR tree with the following properties:

- Each node n has an associated belief state b(n).
- If n is the root node, then $b(n) = \{ s \in S \mid s \models I \}.$
- A leaf node n is labeled with (\top) iff $b(n) \models G$. In this case it is called a goal node, otherwise an open node.
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Partial plan trees

Definition (ctd.)

A partial plan tree for a nondeterministic planning task $\langle A,I,O,G,V\rangle$ with state set S is an AND/OR tree with the following properties:

- ...
- An OR node n (also called an operator node) has one child n_o for each operator $o \in O$ applicable in b(n), with associated belief state $b(n_o) = app_o(b(n))$.
- An AND node n (also called an observation node) has an associated formula $\varphi(n)$ over V. It has two children:
 - $\bullet \ n^\top \text{ with } b(n^\top) = \{ \ s \in b(n) \mid s \models \varphi \}$
 - $\bullet \ n^\perp \text{ with } b(n^\perp) = \{ \ s \in b(n) \mid s \not\models \varphi \}.$

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Forward planning as search in partial plan trees

- Clearly, a partial plan tree represents a strategy.
- This strategy is a strong plan iff the tree evaluates to true.

We thus obtain a (nondeterministic) forward search algorithm:

Forward search in partial plan trees

def expand-tree(\mathcal{T}):

Set T to the partial plan tree for $\mathcal T$ that consists of a single leaf, labeled with the initial belief state.

while T evaluates to false:

Choose some open leaf n in T.

Replace n by an operator or observation node, adding the necessary children to T.

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Search in AND/OR trees

• There is a conflict between plan size and observing:

- With many observations, plans become very big.
- With few observations, it may be impossible to find a plan.

Trying out all possible ways to branch is not feasible. No good general solutions to this problem exist.

- AND-OR search algorithms use heuristics for making branching decisions.
 - But they do not really work well...

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Backward search algorithms

- Backward search algorithms are similar in flavour to the ones for fully observable problems.
- Backward steps with operator application:
 - Compute strong preimages.
- Backward steps with observations:
 - Compute union of belief states from disjoint observational classes.
 - Note: Can always take subsets of solved belief states to make them disjoint.

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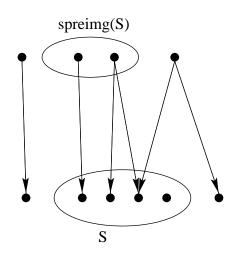
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Backward search algorithms

Regression: strong preimages



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- Let C_1, \ldots, C_n be different observational classes.
- Let B_1, \ldots, B_n be belief states with $B_i \subseteq C_i$ for all $i = 1, \ldots, n$ for which we have a solution plan.
- Then we can find a plan for $B=B_1\cup\cdots\cup B_n$ by first observing in which class C_i we are and then applying the corresponding plan for B_i .

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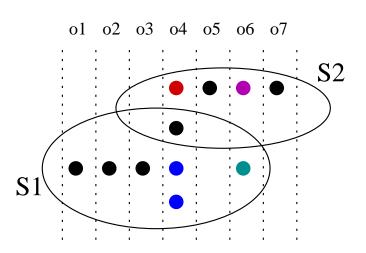
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Example: Combining two belief states



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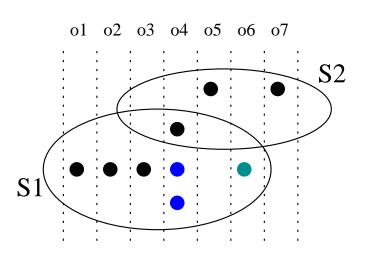
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Example: Combining two belief states, option 1



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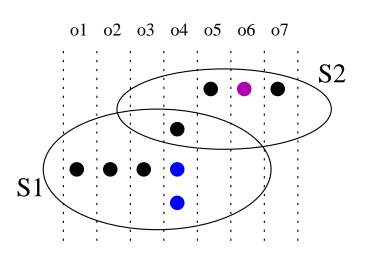
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Example: Combining two belief states, option 2



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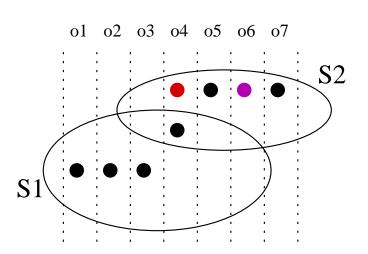
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Example: Combining two belief states, option 3



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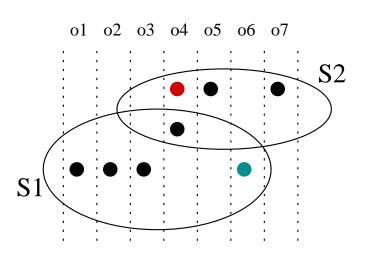
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Observations in backward search

Example: Combining two belief states, option 4



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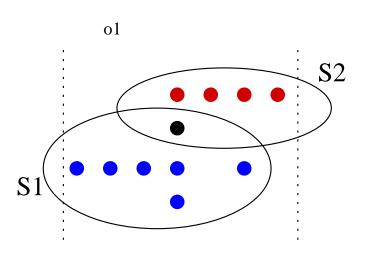
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Observations Algorithm Example

No observability \Rightarrow no branching

Only one observational class: no choice between subplans



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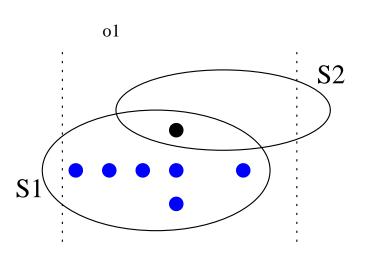
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> Observations Algorithm Example

No observability \Rightarrow no branching

No choice between subplans during execution: option 1



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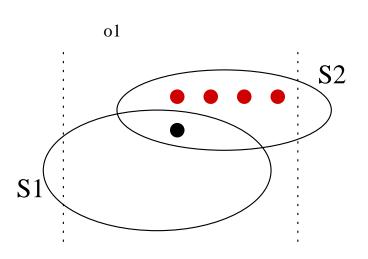
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No observability \Rightarrow no branching

No choice between subplans during execution: option 2



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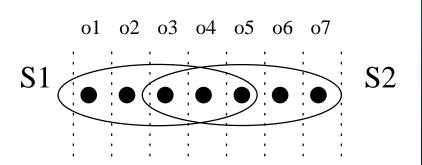
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Full observability ⇒ arbitrary branching A different plan can be used for every state



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A systematic backward algorithm

Idea: always split belief states into all observational classes.

Initially, the set of solved belief states includes the set $b_G \cap C_i$ for each observational class C_i , where b_G is the belief state containing all states satisfying the goal.

Then iterate the following steps:

- Pick one belief state b_i for each observational class and compute their union b.
- ② If b includes all initial states \rightsquigarrow solution.
- **3** Otherwise, compute the strong preimage of b with respect to some operator o.
- Split the resulting set of states to belief states for different observational classes and add them to the set of solved belief states.

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Blocks world with three blocks

- Goal: all blocks are on the table
- Only the variables clear(X) are observable.
- A block can be moved onto the table if the block is clear.
- 8 observational classes corresponding to the 8 valuations of {clear(A), clear(B), clear(C)} (one of the valuations does not correspond to a blocks world state).

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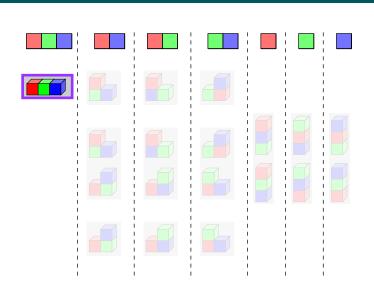
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Example: goal belief state



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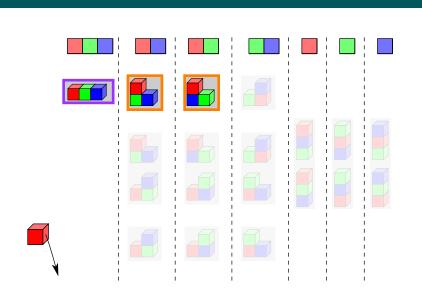
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Example: backward step with red-block-onto-table



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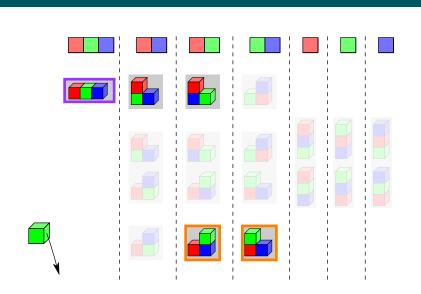
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Example: backward step with green-block-onto-table



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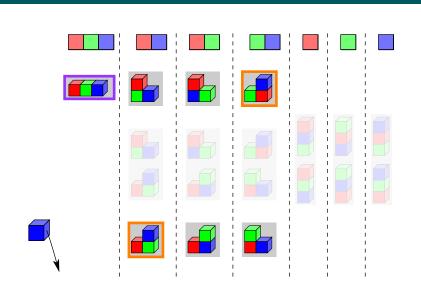
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Example: backward step with blue-block-onto-table



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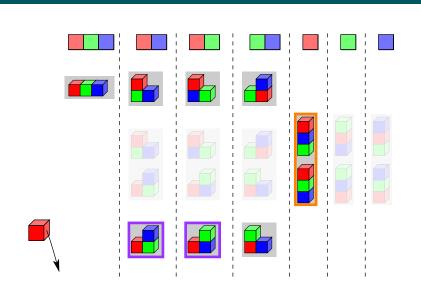
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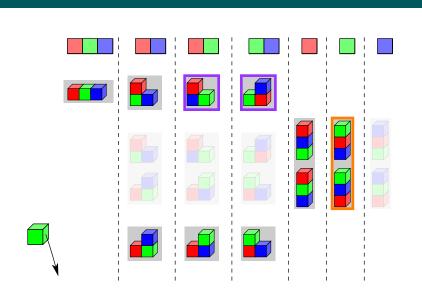
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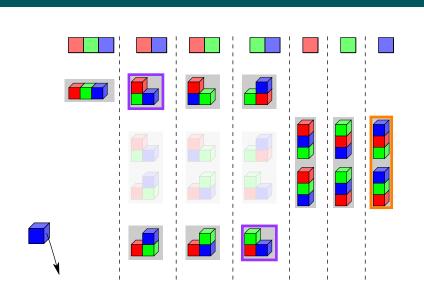
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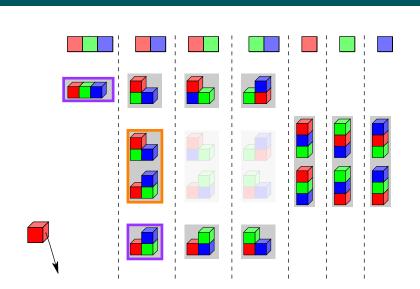
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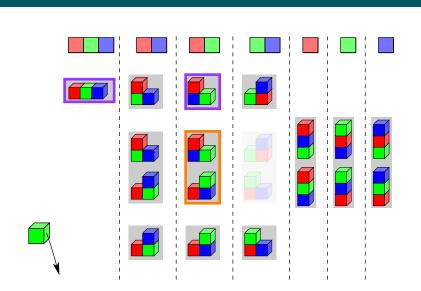
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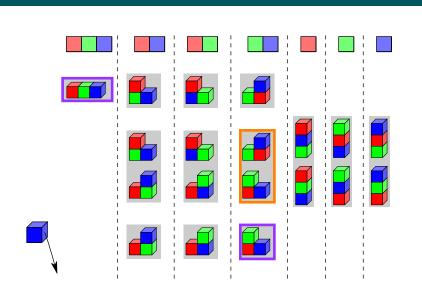
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Example: backward step with blue-block-onto-table



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Summary

- Planning with partial observability in general requires more general classes of plans than the fully observable and unobservable special cases.
- It appears to be significantly harder.
- Algorithmic ideas are similar to the simpler cases:
 - Reduction to full observability by viewing belief states as states.
 - Forward search in AND/OR trees.
 - Dynamic-programming style backward construction of solvable belief states, starting from goal belief states.

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