

# An Introduction to Game Theory

## Part I: Strategic Games

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## Strategic Game

- A **strategic game**  $G$  consists of
  - a finite set  $N$  (the set of **players**)
  - for each player  $i \in N$  a non-empty set  $A_i$  (the set of **actions** or **strategies** available to player  $i$ ), whereby  $A = \prod_i A_i$
  - for each player  $i \in N$  a function  $u_i: A \rightarrow \mathbb{R}$  (the **utility** or **payoff** function)
  - $G = (N, (A_i), (u_i))$
- If  $A$  is finite, then we say that the game is *finite*

## Playing the Game

- Each player  $i$  makes a **decision** which action to play:  $a_i$
- All players make their moves simultaneously leading to the **action profile**  $a^* = (a_1, a_2, \dots, a_n)$
- Then each player gets the **payoff**  $u_i(a^*)$
- Of course, each player tries to maximize its own payoff, but what is the right decision?
- Note:** While we want to maximize our payoff, we are not interested in harming our opponent. It just does not matter to us what he will get!
  - If we want to model something like this, the payoff function must be changed

## Notation

- For 2-player games, we use a matrix, where the strategies of **player 1** are the **rows** and the strategies of **player 2** the **columns**
- The payoff for every action profile is specified as a pair  $x, y$ , whereby  $x$  is the value for player 1 and  $y$  is the value for player 2
- Example: For (T, R), **player 1** gets  $x_{12}$ , and **player 2** gets  $y_{12}$

	Player 2 L action	Player 2 R action
Player 1 T action	$x_{11}, y_{11}$	$x_{12}, y_{12}$
Player 1 B action	$x_{21}, y_{21}$	$x_{22}, y_{22}$

## Example Game: Bach and Stravinsky

- Two people want to out together to a concert of music by either Bach or Stravinsky. Their main concern is to go out together, but one prefers, the other Stravinsky. Will they meet?
- This game is also called the *Battle of the Sexes*

	Bach	Stravinsky
Bach	2, 1	0, 0
Stravinsky	0, 0	1, 2

## Example Game: Hawk-Dove

- Two animals fighting over some prey.
- Each can behave like a dove or a hawk
- The best outcome is if oneself behaves like a hawk and the opponent behaves like a dove
- This game is also called *chicken*.

	Dove	Hawk
Dove	3, 3	1, 4
Hawk	4, 1	0, 0

## Example Game: Prisoner's Dilemma

- Two suspects in a crime are put into separate cells.
- If they both confess, each will be sentenced to 3 years in prison.
- If only one confesses, he will be freed.
- If neither confesses, they will both be convicted of a minor offense and will spend one year in prison.

	Don't confess	Confess
Don't confess	3,3	0,4
Confess	4,0	1,1

## Solving a Game

- What is the right move?
- Different possible **solution concepts**
  - Elimination of strictly or weakly **dominated** strategies
  - Maximin** strategies (for minimizing the loss in zero-sum games)
  - Nash equilibrium**
- How difficult is it to compute a solution?
- Are there always solutions?
- Are the solutions unique?

## Strictly Dominated Strategies

- Notation:**
  - Let  $a = (a_i)$  be a strategy profile
  - $a_{-i} := (a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_n)$
  - $(a_{-i}, a'_i) := (a_1, \dots, a_{i-1}, a'_i, a_{i+1}, \dots, a_n)$
- Strictly dominated strategy:**
  - An strategy  $a_j^* \in A_j$  is **strictly dominated** if there exists a strategy  $a'_j$  such that for all strategy profiles  $a \in A$ :  
 $u_j(a_{-j}, a'_j) > u_j(a_{-j}, a_j^*)$
- Of course, it is **not rational** to play **strictly dominated strategies**

## Iterated Elimination of Strictly Dominated Strategies

- Since strictly dominated strategies will never be played, one can **eliminate** them from the game
- This can be done **iteratively**
- If this converges to a single strategy profile, the result is **unique**
- This can be regarded as the **result** of the game, because it is the **only rational outcome**

## Iterated Elimination: Example

- Eliminate:**
  - b4, dominated by b3
  - a4, dominated by a1
  - b3, dominated by b2
  - a1, dominated by a2
  - b1, dominated by b2
  - a3, dominated by a2
- **Result: (a2,b2)**

	b1	b2	b3	b4
a1	1,7	2,5	7,2	0,1
a2	5,2	3,3	5,2	0,1
a3	7,0	2,5	0,4	0,1
a4	0,0	0,2	0,0	0,1

## Iterated Elimination: Prisoner's Dilemma

- Player 1 reasons that "not confessing" is strictly dominated and eliminates this option
- Player 2 reasons that player 1 will not consider "not confessing". So he will eliminate this option for himself as well
- So, they both confess

	Don't confess	Confess
Don't confess	3,3	0,4
Confess	4,0	1,1

## Weakly Dominated Strategies

- Instead of strict domination, we can also go for weak domination:

- An strategy  $a_j^* \in A_j$  is **weakly dominated** if there exists a strategy  $a_j'$  such that for all strategy profiles  $a \in A$ :

$$u_j(a_{-j}, a_j') \geq u_j(a_{-j}, a_j^*)$$

and for at least one profile  $a \in A$ :

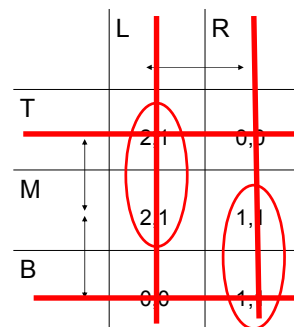
$$u_j(a_{-j}, a_j') > u_j(a_{-j}, a_j^*).$$

## Results of Iterative Elimination of Weakly Dominated Strategies

- The result is not necessarily unique

- Example:

- Eliminate
  - T ( $\leq$ M)
  - L ( $\leq$ R)
 ➤ Result: (1,1)
- Eliminate:
  - B ( $\leq$ M)
  - R ( $\leq$ L)
 ➤ Result (2,1)



## Existence of Dominated Strategies

- Dominating strategies are a convincing **solution concept**
- Unfortunately, often dominated strategies do not exist
- What do we do in this case?

➤ **Nash equilibrium**

	Dove	Hawk
Dove	3,3	1,4
Hawk	4,1	0,0

## Nash Equilibrium

- A **Nash equilibrium** is an action profile  $a^* \in A$  with the property that for all players  $i \in N$ :  
 $u_i(a^*) = u_i(a_{-i}^*, a_i^*) \geq u_i(a_{-i}^*, a_i) \forall a_i \in A_i$
- In words, it is an action profile such that there is **no incentive** for any agent **to deviate** from it
- While it is less convincing than an action profile resulting from iterative elimination of dominated strategies, it is still a **reasonable solution concept**
- If there exists a **unique solution** from **iterated elimination of strictly dominated strategies**, then it is also a **Nash equilibrium**

## Example Nash-Equilibrium: Prisoner's Dilemma

- Don't – Don't
  - not a NE
- Don't – Confess (and vice versa)
  - not a NE
- Confess – Confess
  - NE

	Don't confess	Confess
Don't confess	3,3	0,4
Confess	4,0	1,1

## Example Nash-Equilibrium: Hawk-Dove

- Dove-Dove:
  - not a NE
- Hawk-Hawk
  - not a NE
- Dove-Hawk
  - is a NE
- Hawk-Dove
  - is, of course, another NE
- So, NEs are not necessarily unique

	Dove	Hawk
Dove	3,3	1,4
Hawk	4,1	0,0

## Auctions

- An **object** is to be **assigned** to a player in the set  $\{1, \dots, n\}$  in exchange for a payment.
- Players  $i$  **evaluation** of the object is  $v_i$ , and  $v_1 > v_2 > \dots > v_n$ .
- The mechanism to assign the object is a **sealed-bid auction**: the players simultaneously submit bids (non-negative real numbers)
- The object is given to the player with the lowest index among those who submit the highest bid in exchange for the payment
- The payment for a **first price** auction is the highest bid.
- What are the Nash equilibria in this case?

## Formalization

- Game  $G = (\{1, \dots, n\}, (A_i), (u_i))$
- $A_i$ : bids  $b_i \in \mathbb{R}^+$
- $u_i(b_{-i}, b_i) = v_i - b_i$  if  $i$  has won the auction, 0 otherwise
- Nobody would bid more than his valuation, because this could lead to negative utility, and we could easily achieve 0 by bidding 0.

## Nash Equilibria for First-Price Sealed-Bid Auctions

- The Nash equilibria of this game are all profiles  $b$  with:
  - $b_i \leq b_1$  for all  $i \in \{2, \dots, n\}$ 
    - No  $i$  would bid more than  $v_2$  because it could lead to negative utility
    - If a  $b_i$  (with  $< v_2$ ) is higher than  $b_1$ , player 1 could increase its utility by bidding  $v_2 + \epsilon$
    - So  $i$  wins in all NEs
  - $v_1 \geq b_1 \geq v_2$ 
    - Otherwise, player 1 either loses the bid (and could increase its utility by bidding more) or would have itself negative utility
  - $b_j = b_1$  for at least one  $j \in \{2, \dots, n\}$ 
    - Otherwise player 1 could have gotten the object for a lower bid

## Another Game: Matching Pennies

- Each of two people chooses either **Head** or **Tail**. If the choices differ, player 1 pays player 2 a euro; if they are the same, player 2 pays player 1 a euro.
- This is also a **zero-sum** or **strictly competitive** game
- No NE at all! What shall we do here?

	Head	Tail
Head	1,-1	-1,1
Tail	-1,1	1,-1

## Conclusions

- **Strategic games** are one-shot games, where everybody plays its move simultaneously
- The game outcome is the **action profile** resulting from the individual choices.
- Each player gets a payoff based on its **payoff function** and the resulting action profile.
- **Iterated elimination of strictly dominated strategies** is a convincing solution concept, but unfortunately, most of the time it does not yield a unique solution
- **Nash equilibrium** is another solution concept: Action profiles, where **no player has an incentive to deviate**
- It also might **not be unique** and there can be even infinitely many NEs.
- Also, there is no guarantee for the **existence** of a NE