





# Advanced AI Techniques

## I. Bayesian Networks / 1. Probabilistic Independence and Separation in Graphs

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## **1. Basic Probability Calculus**

## **2. Separation in undirected graphs**

## **3. Separation in directed graphs**

## **4. Markov networks**

## **5. Bayesian networks**



Joint probability distributions

Pain Weightloss Vomiting Adeno	Y				N			
	Y		N		Y		N	
	Y	N	Y	N	Y	N	Y	N
Y	0.220	0.220	0.025	0.025	0.095	0.095	0.010	0.010
N	0.004	0.009	0.005	0.012	0.031	0.076	0.050	0.113

Figure 1: Joint probability distribution of four random variables  $P$  (pain),  $W$  (weightloss),  $V$  (vomiting) and  $A$  (adeno).



## Marginal probability distributions

**Definition 1.** Let  $p$  be a the joint probability of the random variables  $\mathcal{X} := \{X_1, \dots, X_n\}$  and  $\mathcal{Y} \subseteq \mathcal{X}$  a subset thereof. Then

$$p(\mathcal{Y} = y) := p^{\downarrow \mathcal{Y}}(y) := \sum_{x \in \text{dom } \mathcal{X} \setminus \mathcal{Y}} p(\mathcal{X} = x, \mathcal{Y} = y)$$

is a probability distribution of  $\mathcal{Y}$  called **marginal probability distribution**.

**Example 1.**

Vomiting	Y	N
Adeno Y	0.350	0.350
N	0.090	0.210

Pain Weightloss Vomiting	Y		N		N		N	
	Y		N		Y		N	
	Y	N	Y	N	Y	N	Y	N
Adeno Y	0.220	0.220	0.025	0.025	0.095	0.095	0.010	0.010
N	0.004	0.009	0.005	0.012	0.031	0.076	0.050	0.113

Figure 2: Joint probability distribution of four random variables  $P$  (pain),  $W$  (weightloss),  $V$  (vomiting) and  $A$  (adeno).



## Marginal probability distributions / example

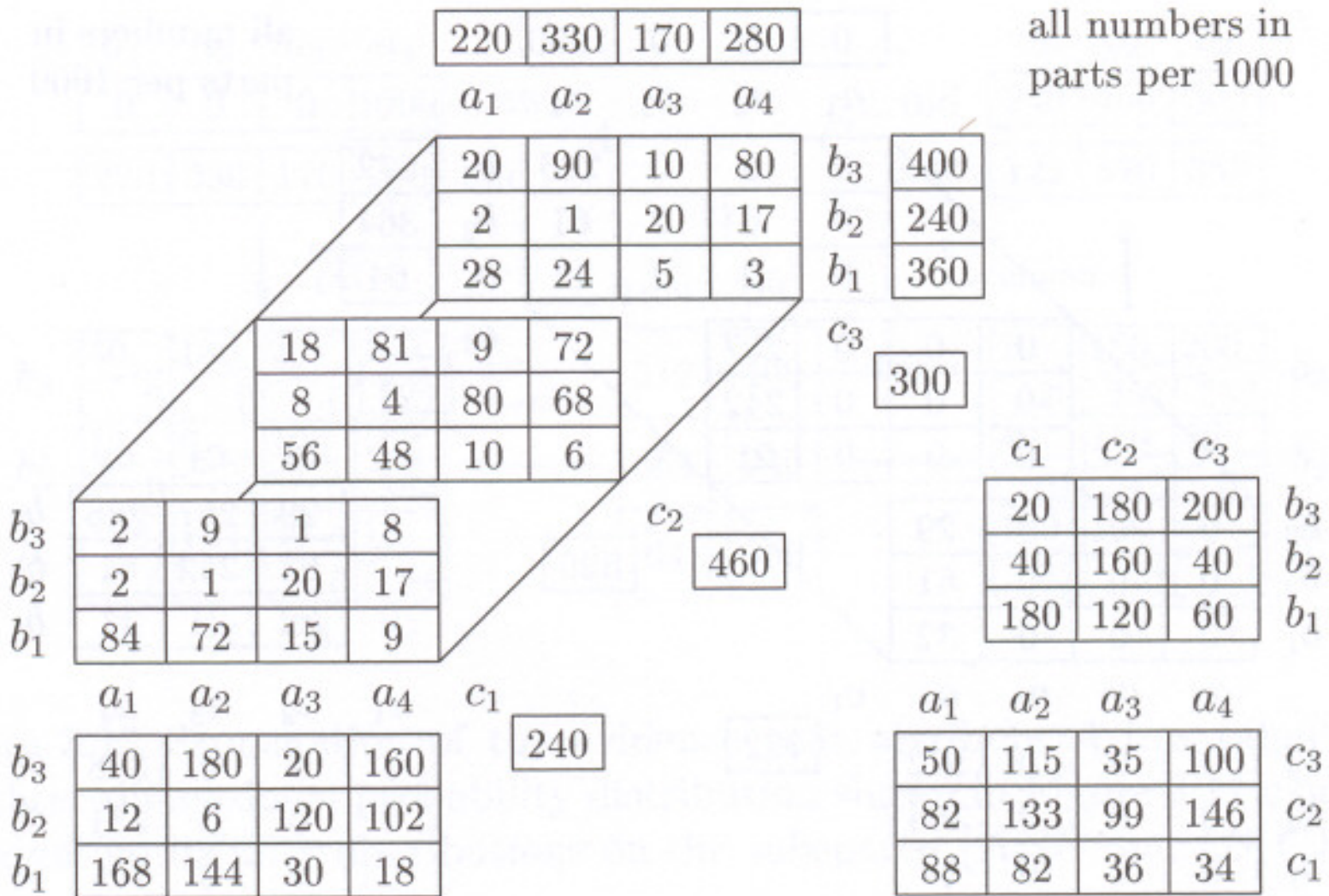


Figure 3: Joint probability distribution and all of its marginals [BK02, p. 75].



## Conditional probability distributions

**Definition 2.** By  $p > 0$  we mean

$$p(x) > 0, \quad \text{for all } x \in \prod \text{dom}(p)$$

Then  $p$  is called **non-extreme**.

For a JPD  $p$  and a subset  $\mathcal{Y} \subseteq \text{dom}(p)$  of its variables with  $p^{\downarrow \mathcal{Y}} > 0$  we define

$$p^{\downarrow \mathcal{Y}} := \frac{p}{p^{\downarrow \mathcal{Y}}}$$

as **conditional probability distribution of  $p$  w.r.t.  $\mathcal{Y}$** .

A conditional probability distribution w.r.t.  $\mathcal{Y}$  sums to 1 for all fixed values of  $\mathcal{Y}$ , i.e.,

$$(p^{\downarrow \mathcal{Y}})^{\downarrow \mathcal{Y}} \equiv 1$$

**Example 2.** Let  $p$  be the JPD

$$p := \begin{pmatrix} 0.4 & 0.1 \\ 0.2 & 0.3 \end{pmatrix}$$

on two variables  $R$  (rows) and  $C$  (columns) with the domains  $\text{dom}(R) = \text{dom}(C) = \{1, 2\}$ .

The conditional probability distribution w.r.t.  $C$  is

$$p^{\downarrow C} := \begin{pmatrix} 2/3 & 1/4 \\ 1/3 & 3/4 \end{pmatrix}$$



## Chain rule

**Lemma 1 (Chain rule).** *Let  $X_1, X_2, \dots, X_n$  be variables. Then*

$$p(X_1, X_2, \dots, X_n) = p(X_n | X_1, \dots, X_{n-1}) \cdots p(X_2 | X_1) \cdot p(X_1)$$



## Independent variables

**Definition 3.** Two sets  $\mathcal{X}, \mathcal{Y}$  of variables are called **independent**, when all pairs of events  $\mathcal{X} = x$  and  $\mathcal{Y} = y$  are independent, i.e.

$$p(\mathcal{X} = x, \mathcal{Y} = y) = p(\mathcal{X} = x) \cdot p(\mathcal{Y} = y)$$

for all  $x$  and  $y$  or equivalently

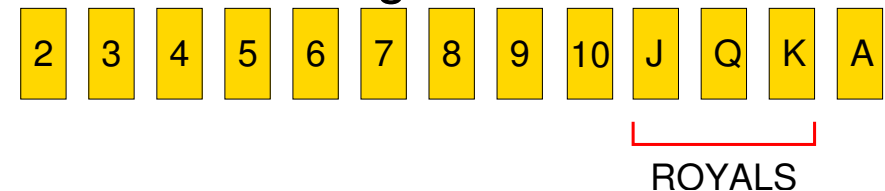
$$p(\mathcal{X} = x | \mathcal{Y} = y) = p(\mathcal{X} = x)$$

for  $y$  with  $p(\mathcal{Y} = y) > 0$ .

**Example 3.** Let  $\Omega$  be the cards in an ordinary deck and

- $R$  = true, if a card is royal,
- $T$  = true, if a card is a ten or a jack,
- $S$  = true, if a card is spade.

Cards for a single color:



$S$	$R$	$T$	$p(R, T   S)$
Y	Y	Y	1/13
		N	2/13
	N	Y	1/13
		N	9/13
N	Y	Y	3/39 = 1/13
		N	6/39 = 2/13
	N	Y	3/39 = 1/13
		N	27/39 = 9/13

$R$	$T$	$p(R, T)$
Y	Y	4/52 = 1/13
	N	8/52 = 2/13
N	Y	4/52 = 1/13
	N	36/52 = 9/13



## Conditionally independent variables

**Definition 4.** Let  $\mathcal{X}$ ,  $\mathcal{Y}$ , and  $\mathcal{Z}$  be sets of variables.

$\mathcal{X}, \mathcal{Y}$  are called **conditionally independent given  $\mathcal{Z}$** , when for all events  $\mathcal{Z} = z$  with  $p(\mathcal{Z} = z) > 0$  all pairs of events  $\mathcal{X} = x$  and  $\mathcal{Y} = y$  are conditionally independent given  $\mathcal{Z} = z$ , i.e.

$$p(\mathcal{X} = x, \mathcal{Y} = y, \mathcal{Z} = z) = \frac{p(\mathcal{X} = x, \mathcal{Z} = z) \cdot p(\mathcal{Y} = y, \mathcal{Z} = z)}{p(\mathcal{Z} = z)}$$

for all  $x, y$  and  $z$  (with  $p(\mathcal{Z} = z) > 0$ ), or equivalently

$$p(\mathcal{X} = x | \mathcal{Y} = y, \mathcal{Z} = z) = p(\mathcal{X} = x | \mathcal{Z} = z)$$

We write  $I_p(\mathcal{X}, \mathcal{Y} | \mathcal{Z})$  for the statement, that  $\mathcal{X}$  and  $\mathcal{Y}$  are conditionally independent given  $\mathcal{Z}$ .



## Conditionally independent variables

**Example 4.** Assume  $S$  (shape),  $C$  (color), and  $L$  (label) be three random variables that are distributed as shown in figure 4.

We show  $I_p(\{L\}, \{S\} | \{C\})$ , i.e., that label and shape are conditionally independent given the color.

$C$	$S$	$L$	$p(L C, S)$
black	square	1	$2/6 = 1/3$
		2	$4/6 = 2/3$
	round	1	$1/3$
		2	$2/3$
white	square	1	$1/2$
		2	$1/2$
	round	1	$1/2$
		2	$1/2$

$C$	$L$	$p(L C)$
black	1	$3/9 = 1/3$
	2	$6/9 = 2/3$
white	1	$2/4 = 1/2$
	2	$2/4 = 1/2$

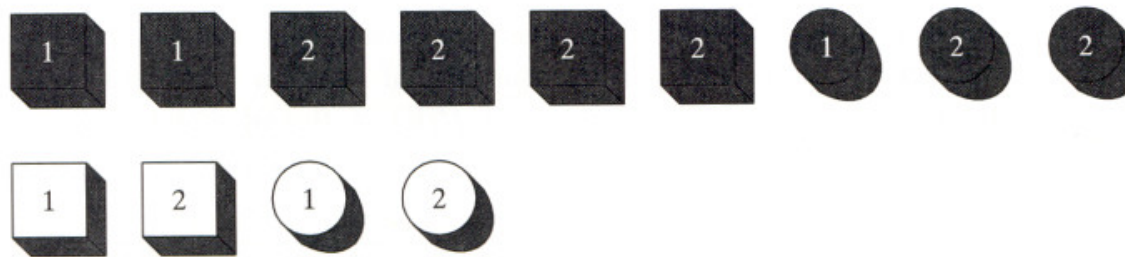


Figure 4: 13 objects with different shape, color, and label [Nea03, p. 8].



## **1. Basic Probability Calculus**

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## Separation in graphs (u-separation)

**Definition 5.** Let  $G := (V, E)$  be a graph. Let  $Z \subseteq V$  be a subset of vertices. We say, two vertices  $x, y \in V$  are **u-separated by  $Z$  in  $G$** , if every path from  $x$  to  $y$  contains some vertex of  $Z$  ( $\forall p \in G^* : p_1 = x, p_{|p|} = y \Rightarrow \exists i \in \{1, \dots, n\} : p_i \in Z$ ).

Let  $X, Y, Z \subseteq V$  be three disjoint subsets of vertices. We say, the vertices  $X$  and  $Y$  are **u-separated by  $Z$  in  $G$** , if every path from any vertex from  $X$  to any vertex from  $Y$  is separated by  $Z$ , i.e., contains some vertex of  $Z$ .

We write  $I_G(X, Y|Z)$  for the statement, that  $X$  and  $Y$  are u-separated by  $Z$  in  $G$ .

$I_G$  is called **u-separation relation in  $G$** .

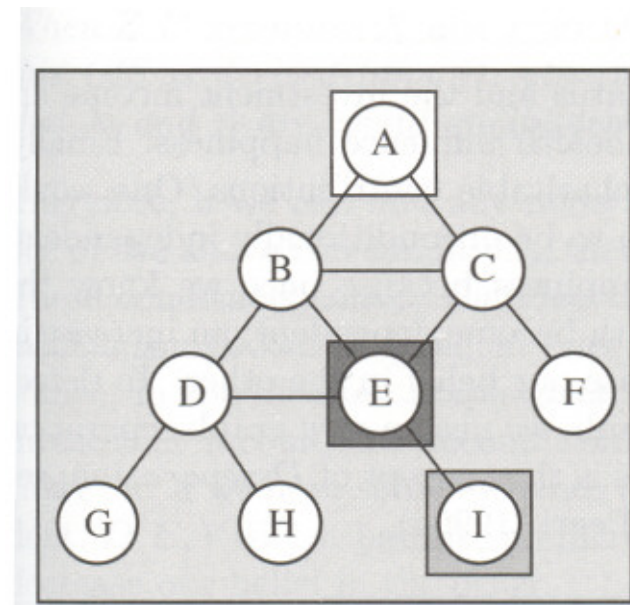


Figure 5: Example for u-separation [CGH97, p. 179].



## Separation in graphs (u-separation)

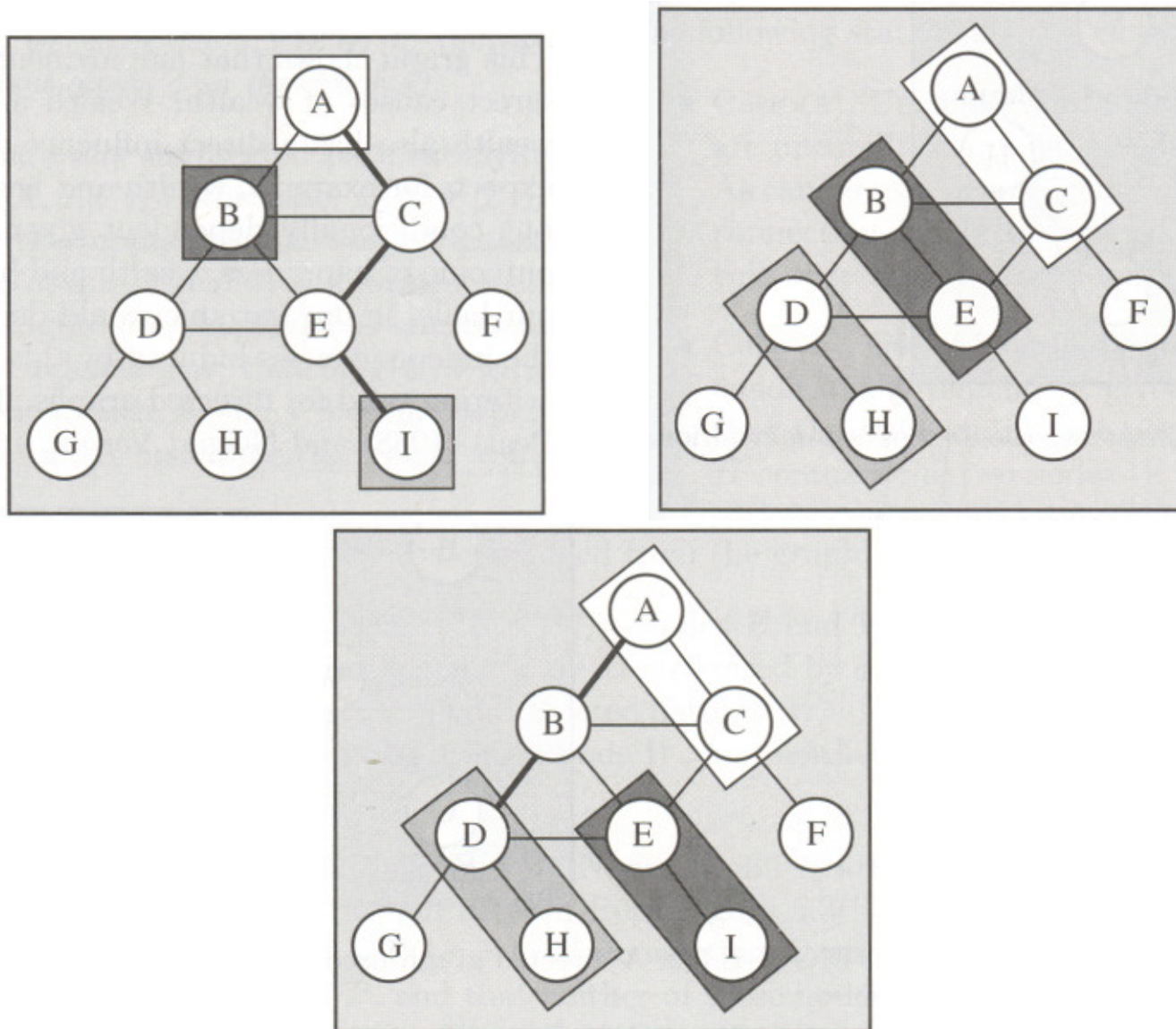


Figure 6: More examples for u-separation [CGH97, p. 179].



## Checking u-separation

To test, if for a given graph  $G = (V, E)$  two given sets  $X, Y \subseteq V$  of vertices are u-separated by a third given set  $Z \subseteq V$  of vertices, we may use standard breadth-first search to compute all vertices that can be reached from  $X$  (see, e.g., [OW02], [CLR90]).

```

1 breadth-first search( $G, X$ ) :
2    $border := X$ 
3    $reached := \emptyset$ 
4   while  $border \neq \emptyset$  do
5        $reached := reached \cup border$ 
6        $border := fan_G(border) \setminus reached$ 
7   od
8   return  $reached$ 

```

Figure 7: Breadth-first search algorithm for enumerating all vertices reachable from  $X$ .

For checking u-separation we have to tweak the algorithm

1. not to add vertices from  $Z$  to the border and
2. to stop if a vertex of  $Y$  has been reached.

```

1 check-u-separation( $G, X, Y, Z$ ) :
2    $border := X$ 
3    $reached := \emptyset$ 
4   while  $border \neq \emptyset$  do
5        $reached := reached \cup border$ 
6        $border := fan_G(border) \setminus reached \setminus Z$ 
7       if  $border \cap Y \neq \emptyset$ 
8           return  $false$ 
9       fi
10  od
11  return  $true$ 

```

Figure 8: Breadth-first search algorithm for checking u-separation of  $X$  and  $Y$  by  $Z$ .



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## Chains

**Definition 6.** Let  $G := (V, E)$  be a directed graph. We can construct an **undirected skeleton**  $u(G) := (V, u(E))$  of  $G$  by dropping the directions of the edges:

$$u(E) := \{\{x, y\} \mid (x, y) \in E \text{ or } (y, x) \in E\}$$

The paths on  $u(G)$  are called **chains of  $G$** :

$$G^\blacktriangle := u(G)^*$$

i.e., a chain is a sequence of vertices that are linked by a forward or a backward edge. If we want to stress the directions of the linking edges, we denote a chain  $p = (p_1, \dots, p_n) \in G^\blacktriangle$  by

$$p_1 \leftarrow p_2 \rightarrow p_3 \leftarrow \dots \leftarrow p_{n-1} \rightarrow p_n$$

The notions of **length**, **subchain**, **interior** and **proper** carry over from undirected paths to chains.

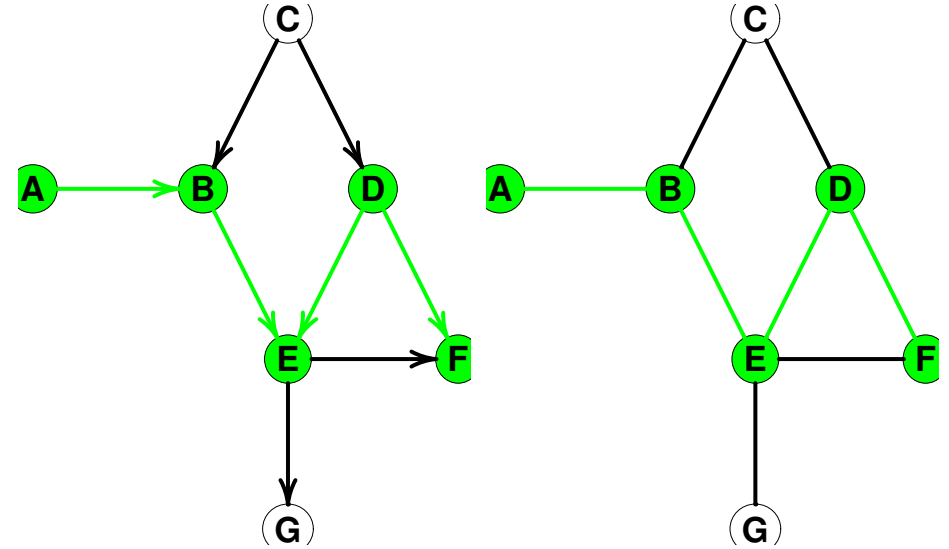


Figure 9: Chain  $(A, B, E, D, F)$  on directed graph and path on undirected skeleton.



## Blocked chains

**Definition 7.** Let  $G := (V, E)$  be a directed graph. We call a chain

$$p_1 \rightarrow p_2 \leftarrow p_3$$

a **head-to-head meeting**.

Let  $Z \subseteq V$  be a subset of vertices.

Then a chain  $p \in G^\Delta$  is called **blocked at position  $i$  by  $Z$** , if for its subchain  $(p_{i-1}, p_i, p_{i+1})$  there is

$$\begin{cases} p_i \in Z, & \text{if not } p_{i-1} \rightarrow p_i \leftarrow p_{i+1} \\ p_i \notin Z \cup \text{anc}(Z), & \text{else} \end{cases}$$

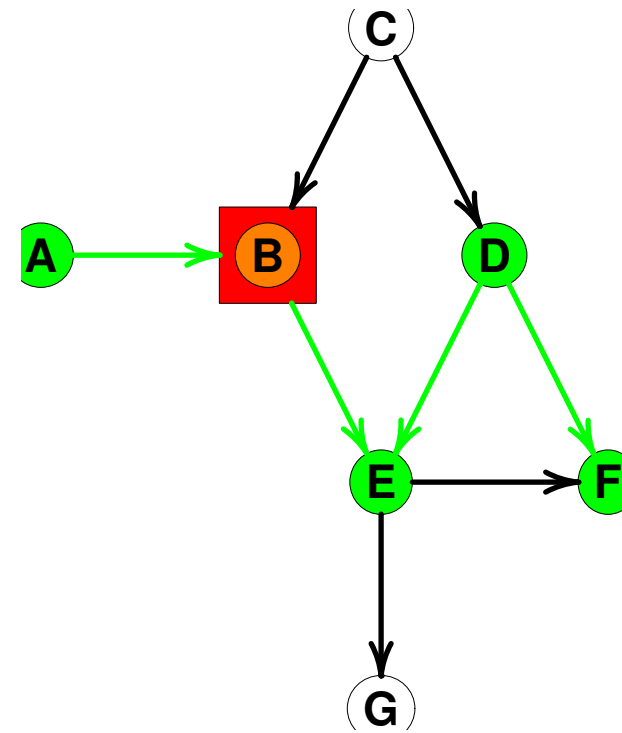


Figure 10: Chain  $(A, B, E, D, F)$  is blocked by  $Z = \{B\}$  at 2.



## Blocked chains / more examples

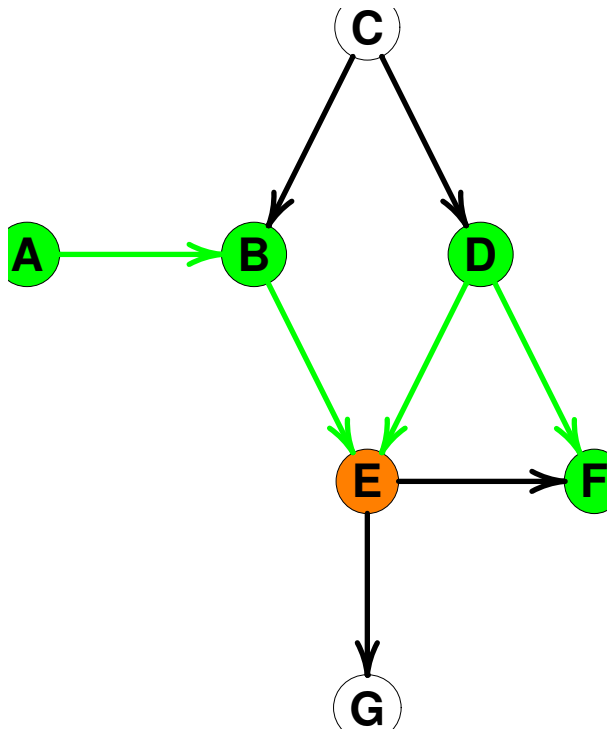


Figure 11: Chain  $(A, B, E, D, F)$  is blocked by  $Z = \emptyset$  at 3.

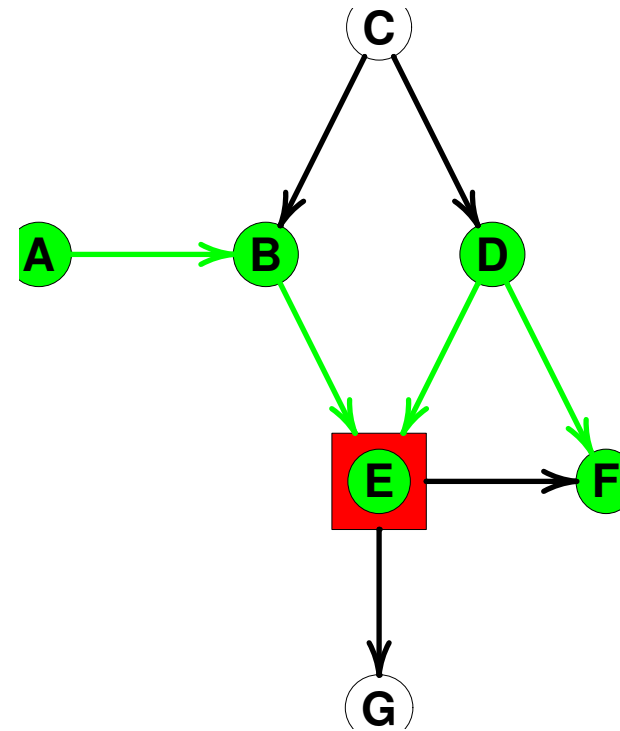


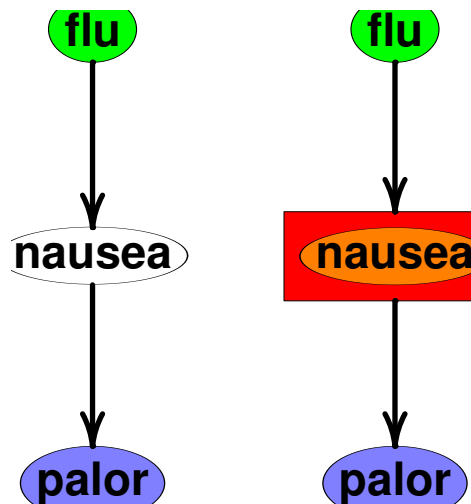
Figure 12: Chain  $(A, B, E, D, F)$  is **not** blocked by  $Z = \{E\}$  at 3.



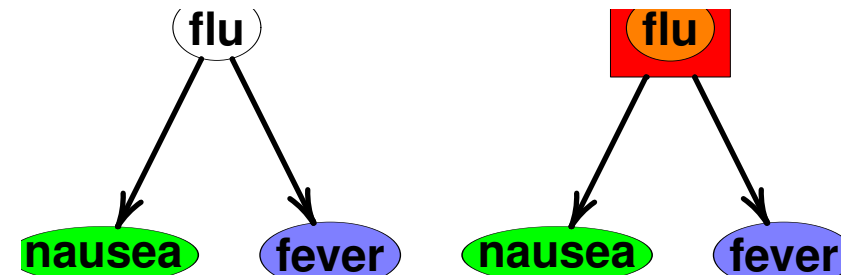
## Blocked chains / rationale

The notion of blocking is chosen in a way so that chains model "flow of causal influence" through a causal network where the states of the vertices  $Z$  are already known.

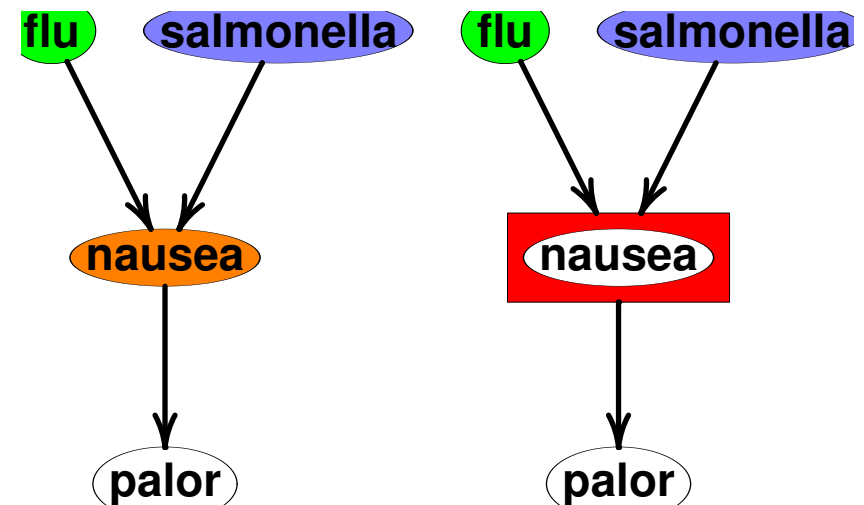
1) Serial connection / intermediate cause:



2) Diverging connection / common cause:



3) Converging connection / common effect:



Models "discounting" [Nea03, p. 51].

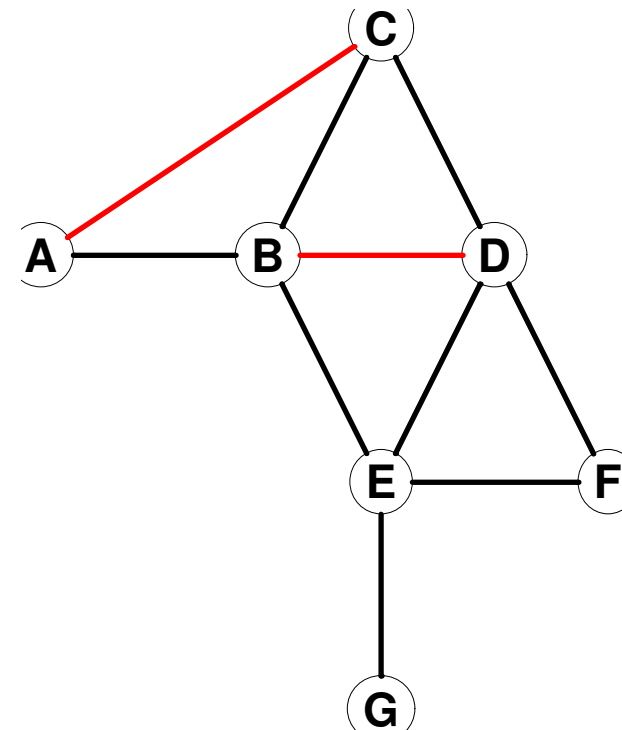
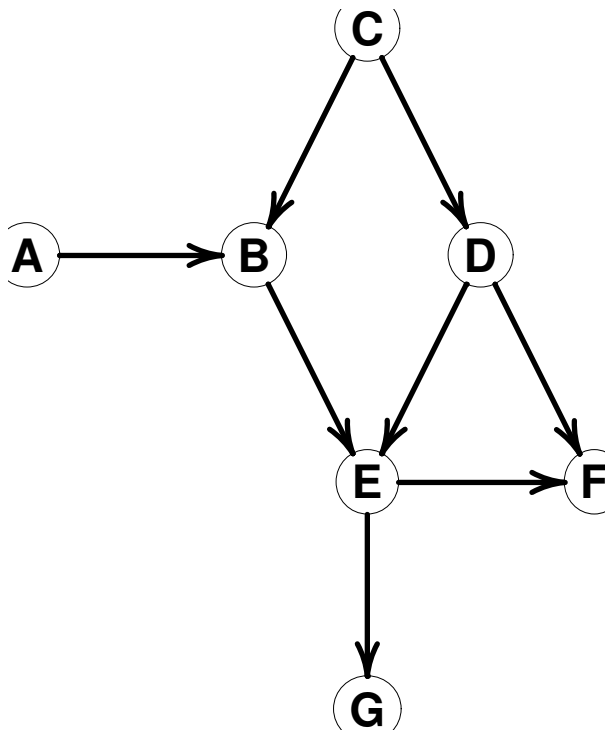


## The moral graph

**Definition 8.** Let  $G := (V, E)$  be a DAG.

As the **moral graph** of  $G$  we denote the undirected skeleton graph of  $G$  plus additional edges between each two parents of a vertex, i.e.  $\text{moral}(G) := (V, E')$  with

$$E' := u(E) \cup \{\{x, y\} \mid \exists z \in V : x, y \in \text{pa}(z)\}$$





## Separation in DAGs (d-separation)

Let  $G := (V, E)$  be a DAG.

Let  $X, Y, Z \subseteq V$  be three disjoint subsets of vertices. We say, the vertices  $X$  and  $Y$  are **separated by  $Z$  in  $G$** , if

- (i) every chain from any vertex from  $X$  to any vertex from  $Y$  is blocked by  $Z$  *or equivalently*
- (ii)  $X$  and  $Y$  are u-separated by  $Z$  in the moral graph of the ancestral hull of  $X \cup Y \cup Z$ .

We write  $I_G(X, Y|Z)$  for the statement, that  $X$  and  $Y$  are separated by  $Z$  in  $G$ .

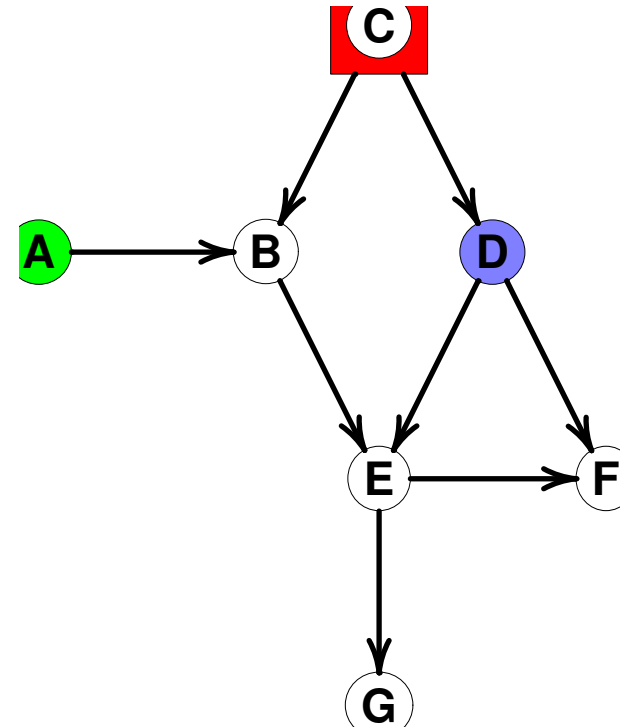
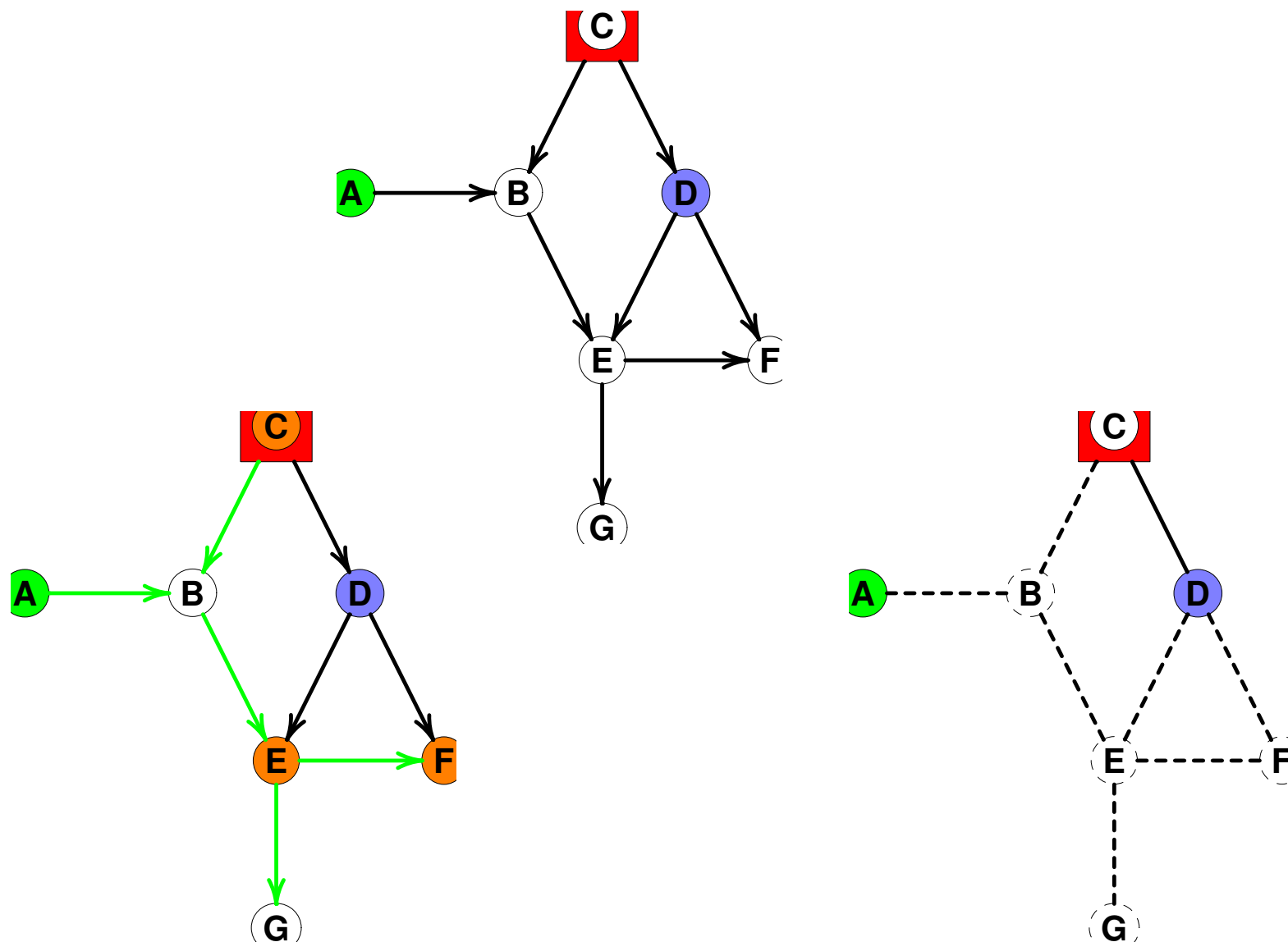


Figure 15: Are the vertices  $A$  and  $D$  separated by  $C$  in  $G$ ?

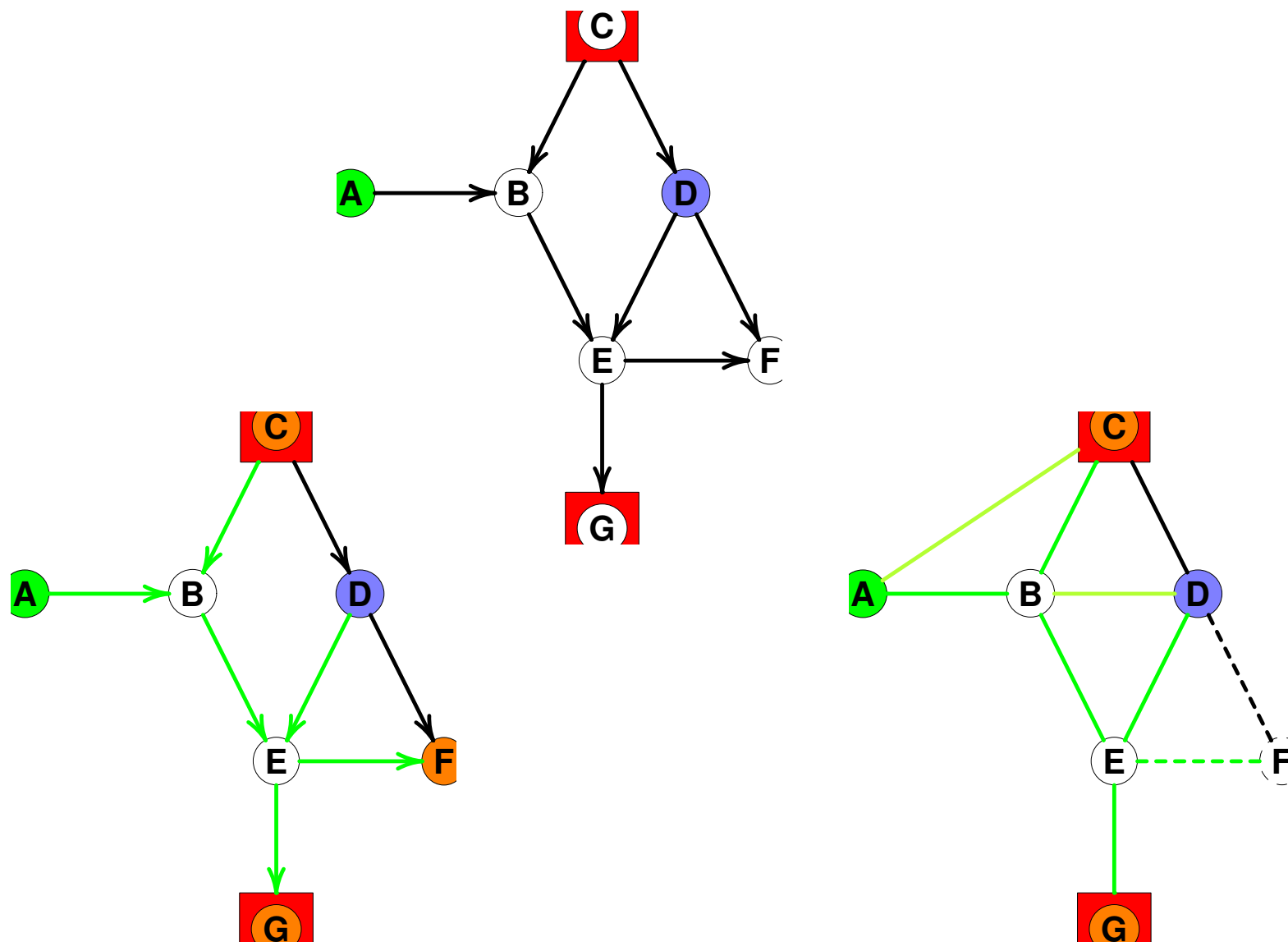


## Separation in DAGs (d-separation) / examples

Figure 16:  $A$  and  $D$  are separated by  $C$  in  $G$ .



## Separation in DAGs (d-separation) / more examples

Figure 17:  $A$  and  $D$  are not separated by  $\{C, G\}$  in  $G$ .



## Checking d-separation

To test, if for a given graph  $G = (V, E)$  two given sets  $X, Y \subseteq V$  of vertices are d-separated by a third given set  $Z \subseteq V$  of vertices, we may

- build the moral graph of the ancestral hull and
- apply the u-separation criterion.

```
1 check-d-separation( $G, X, Y, Z$ ) :  
2  $G' := \text{moral}(\text{anc}_G(X \cup Y \cup Z))$   
3 return check-u-separation( $G', X, Y, Z$ )
```

Figure 18: Algorithm for checking d-separation via u-separation in the moral graph.

A drawback of this algorithm is that we have to rebuild the moral graph of the ancestral hull whenever  $X$  or  $Y$  changes.



## Checking d-separation

Instead of constructing a moral graph, we can modify a **breadth-first search for chains** to find all vertices not d-separated from  $X$  by  $Z$  in  $G$ .

The breadth-first search must not hop over head-to-head meetings with the middle vertex not in  $Z$  nor having an descendent in  $Z$ .

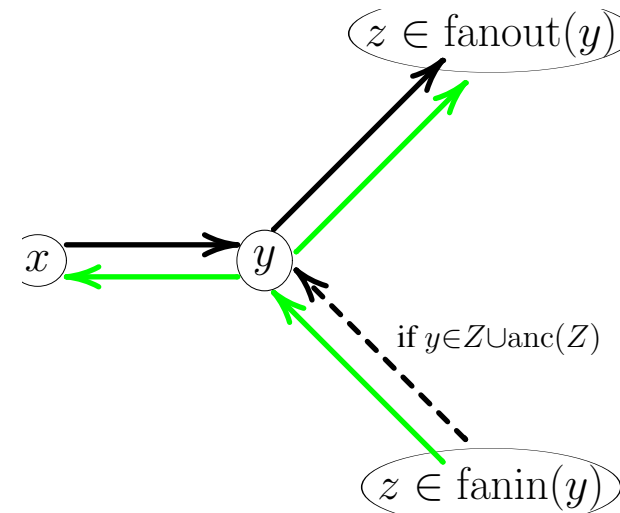


Figure 19: Restricted breadth-first search of non-blocked chains.

```

1 enumerate-d-separation( $G = (V, E), X, Z$ ) :
2    $borderForward := \emptyset$ 
3    $borderBackward := X \setminus Z$ 
4    $reached := \emptyset$ 
5   while  $borderForward \neq \emptyset$  or  $borderBackward \neq \emptyset$  do
6      $reached := reached \cup (borderForward \setminus Z) \cup borderBackward$ 
7      $borderForward := fanout_G(borderBackward \cup (borderForward \setminus Z)) \setminus reached$ 
8      $borderBackward := fanin_G(borderBackward \cup (borderForward \cap (Z \cup anc(Z)))) \setminus Z \setminus reached$ 
9   od
10  return  $V \setminus reached$ 

```

Figure 20: Algorithm for enumerating all vertices d-separated from  $X$  by  $Z$  in  $G$  via restricted breadth-first search (see [Nea03, p. 80–86] for another formulation).



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## Complete graphs, orderings

**Definition 9.** An undirected graph  $G := (V, E)$  is called **complete**, if it contains all possible edges (i.e. if  $E = \mathcal{P}^2(V)$ ).

**Definition 10.** Let  $G := (V, E)$  be a directed graph.

A bijective map

$$\sigma : \{1, \dots, |V|\} \rightarrow V$$

is called an **ordering of (the vertices of)  $G$** .

We can write an ordering as enumeration of  $V$ , i.e. as  $v_1, v_2, \dots, v_n$  with  $V = \{v_1, \dots, v_n\}$  and  $v_i \neq v_j$  for  $i \neq j$ .

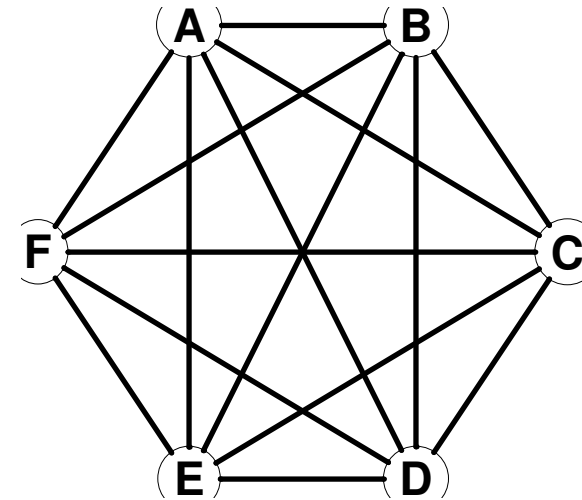


Figure 21: Undirected complete graph with 6 vertices.



## Topological orderings (1/2)

**Definition 11.** An ordering  $\sigma = (v_1, \dots, v_n)$  is called **topological ordering** if

(i) all parents of a vertex have smaller numbers, i.e.

$$\text{fanin}(v_i) \subseteq \{v_1, \dots, v_{i-1}\}, \quad \forall i = 1, \dots, n$$

or equivalently

(ii) all edges point from smaller to larger numbers

$$(v, w) \in E \Rightarrow \sigma^{-1}(v) < \sigma^{-1}(w), \quad \forall v, w \in V$$

The reverse of a topological ordering – e.g. got by using the `fanout` instead of the `fanin` – is called **ancestral numbering**.

In general there are several topological orderings of a DAG.

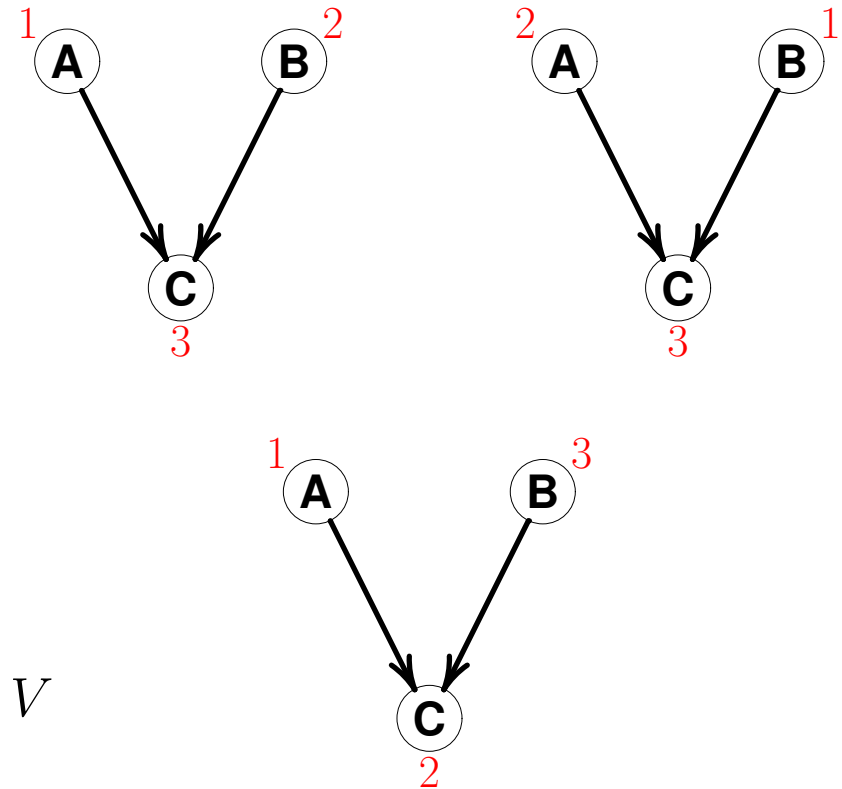


Figure 22: DAG with different topological orderings:  $\sigma_1 = (A, B, C)$  and  $\sigma_2 = (B, A, C)$ . The ordering  $\sigma_3 = (A, C, B)$  is not topological.



## Topological orderings (2/2)

**Lemma 2.** *Let  $G$  be a directed graph. Then*

*$G$  is acyclic (a DAG)  $\Leftrightarrow G$  has a topological ordering*

```
1 topological-ordering( $G = (V, E)$ ) :  
2 choose  $v \in V$  with  $\text{fanout}(v) = \emptyset$   
3  $\sigma(|V|) := v$   
4  $\sigma|_{\{1, \dots, |V|-1\}} := \text{topological-ordering}(G \setminus \{v\})$   
5 return  $\sigma$ 
```

Figure 23: Algorithm to compute a topological ordering of a DAG.

Exercise: write an algorithm for checking if a given directed graph is a acyclic.



## Complete DAGs

**Definition 12.** A DAG  $G := (V, E)$  is called complete, if

- (i) it has a topological ordering  $\sigma = (v_1, \dots, v_n)$  with  
 $\text{fanin}(v_i) = \{v_1, \dots, v_{i-1}\}, \quad \forall i = 1, \dots, n$   
*or equivalently*
- (ii) it has exactly one topological ordering  
*or equivalently*
- (iii) every additional edge introduces a cycle.

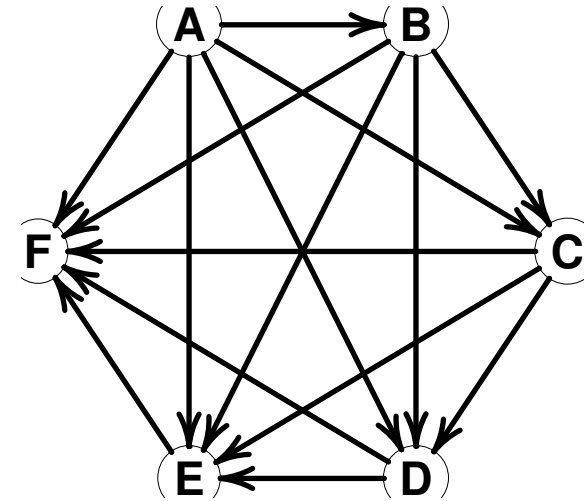


Figure 24: Complete DAG with 6 vertices. Its topological ordering is  $\sigma = (A, B, C, D, E, F)$ .



Graph representations of ternary relations on  $\mathcal{P}(V)$ 

**Definition 13.** Let  $V$  be a set and  $I$  a ternary relation on  $\mathcal{P}(V)$  (i.e.  $I \subseteq \mathcal{P}(V)^3$ ). In our context  $I$  is often called an **independency model**.

Let  $G$  be a graph on  $V$  (undirected or DAG).

$G$  is called a **representation of  $I$** , if

$$I_G(X, Y|Z) \Rightarrow I(X, Y|Z) \quad \forall X, Y, Z \subseteq V$$

A representation  $G$  of  $I$  is called **faithful**, if

$$I_G(X, Y|Z) \Leftrightarrow I(X, Y|Z) \quad \forall X, Y, Z \subseteq V$$

Representations are also called **independency maps of  $I$**  or **markov w.r.t.  $I$** , faithful representations are also called **perfect maps of  $I$** .

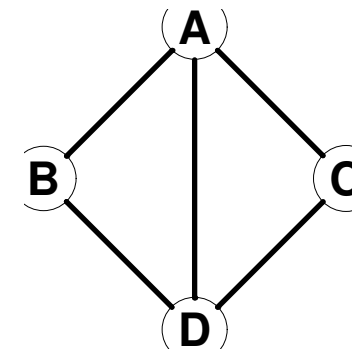


Figure 25: Non-faithful representation of

$$I := \{(A, B|\{C, D\}), (B, C|\{A, D\}), (B, A|\{C, D\}), (C, B|\{A, D\})\}$$

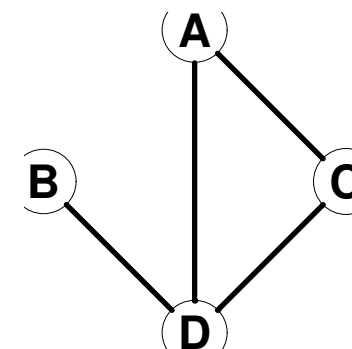


Figure 26: Faithful representation of  $I$ . Which  $I$ ?



## Faithful representations

In  $G$  also holds

$$I_G(B, \{A, C\} | D), I_G(B, A | D), I_G(B, C | D), \dots$$

so  $G$  is not a representation of

$$I := \{(A, B | \{C, D\}), (B, C | \{A, D\}), \\ (B, A | \{C, D\}), (C, B | \{A, D\})\}$$

at all. It is a representation of

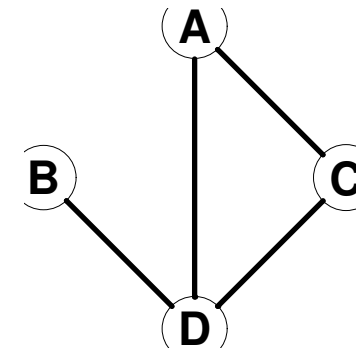


Figure 27: Faithful representation of  $J$ .

$$J := \{(A, B | \{C, D\}), (B, C | \{A, D\}), (B, \{A, C\} | D), (B, A | D), (B, C | D), \\ (B, A | \{C, D\}), (C, B | \{A, D\}), (\{A, C\}, B | D), (A, B | D), (C, B | D)\}$$

and as all independency statements of  $J$  hold in  $G$ , it is faithful.



## Trivial representations

For a complete undirected graph or a complete DAG  $G := (V, E)$  there is

$$I_G \equiv \text{false},$$

i.e. there are no triples  $X, Y, Z \subseteq V$  with  $I_G(X, Y|Z)$ . Therefore  $G$  represents any independency model  $I$  on  $V$  and is called **trivial representation**.

There are independency models without faithful representation.

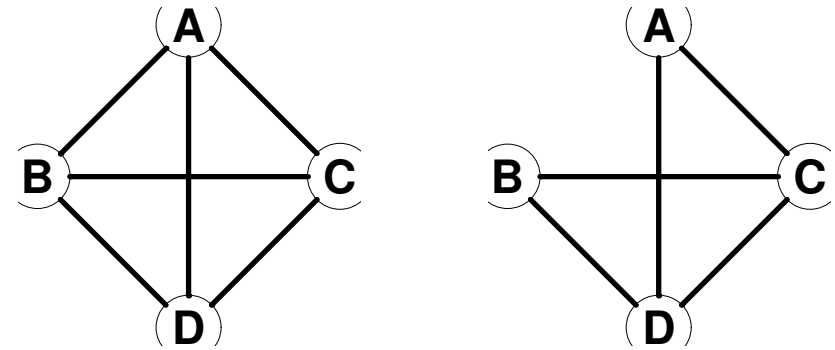


Figure 28: Independency model

$$I := \{(A, B|\{C, D\})\}$$

without faithful representation.



## Minimal representations

**Definition 14.** A representation  $G$  of  $I$  is called **minimal**, if none of its subgraphs omitting an edge is a representation of  $I$ .

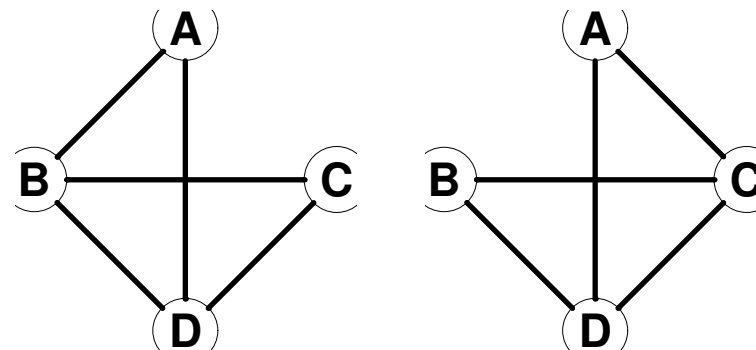


Figure 29: Different minimal undirected representations of the independence model

$$I := \{(A, B | \{C, D\}), (A, C | \{B, D\}), \\ (B, A | \{C, D\}), (C, A | \{B, D\})\}$$



## Minimal representations

**Lemma 3 (uniqueness of minimal undirected representation).**

*An independency model  $I$  has exactly one minimal undirected representation, if and only if it is*

*(i) symmetric:  $I(X, Y|Z) \Rightarrow I(Y, X|Z)$ .*

*(ii) decomposable:  $I(X, Y|Z) \Rightarrow I(X, Y'|Z)$  for any  $Y' \subseteq Y$*

*(iii) intersectable:  $I(X, Y|Y' \cup Z)$  and  $I(X, Y'|Y \cup Z) \Rightarrow I(X, Y \cup Y'|Z)$*

*Then this representation is  $G = (V, E)$  with*

$$E := \{\{x, y\} \in \mathcal{P}^2(V) \mid \text{not } I(x, y|V \setminus \{x, y\})\}$$



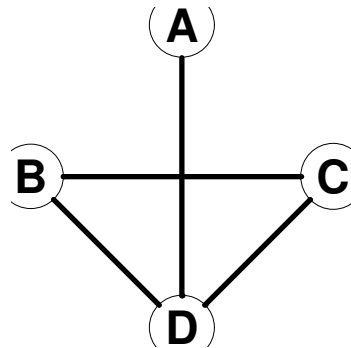
## Minimal representations (2/2)

**Example 5.**

$$I := \{(A, B|\{C, D\}), (A, C|\{B, D\}), (A, \{B, C\}|D), (A, B|D), (A, C|D), \\ (B, A|\{C, D\}), (C, A|\{B, D\}), (\{B, C\}, A|D), (B, A|D), (C, A|D)\}$$

is symmetric, decomposable and intersectable.

Its unique minimal undirected representation is



If a faithful representation exists, obviously it is the unique minimal representation, and thus can be constructed by the rule in lemma 3.



## Representation of conditional independency

**Definition 15.** We say, a graph **represents a JPD**  $p$ , if it represents the conditional independency relation  $I_p$  of  $p$ .

General JPDs may have several minimal undirected representations (as they may violate the intersection property).

Non-extreme JPDs have a unique minimal undirected representation.

To compute this representation we have to check  $I_p(X, Y | V \setminus \{X, Y\})$  for all pairs of variables  $X, Y \in V$ , i.e.

$$p \cdot p^{\downarrow V \setminus \{X, Y\}} = p^{\downarrow V \setminus \{X\}} \cdot p^{\downarrow V \setminus \{Y\}}$$

Then the minimal representation is the complete graph on  $V$  omitting the edges  $\{X, Y\}$  for that  $I_p(X, Y | V \setminus \{X, Y\})$  holds.



## Representation of conditional independency

**Example 6.** Let  $p$  be the JPD on  $V := \{X, Y, Z\}$  given by:

$Z$	$X$	$Y$	$p(X, Y, Z)$
0	0	0	0.024
0	0	1	0.056
0	1	0	0.036
0	1	1	0.084
1	0	0	0.096
1	0	1	0.144
1	1	0	0.224
1	1	1	0.336

Checking  $p \cdot p^{\downarrow V \setminus \{X, Y\}} = p^{\downarrow V \setminus \{X\}} \cdot p^{\downarrow V \setminus \{Y\}}$  one finds that the only independency relations of  $p$  are  $I_p(X, Y|Z)$  and  $I_p(Y, X|Z)$ .

Its marginals are:

$Z$	$X$	$p(X, Z)$
0	0	0.08
0	1	0.12
1	0	0.24
1	1	0.56

$Z$	$Y$	$p(Y, Z)$
0	0	0.06
0	1	0.14
1	0	0.32
1	1	0.48

$X$	$Y$	$p(X, Y)$
0	0	0.12
0	1	0.2
1	0	0.26
1	1	0.42

$X$	$p(X)$
0	0.32
1	0.68

$Y$	$p(Y)$
0	0.38
1	0.62

$Z$	$p(Z)$
0	0.2
1	0.8



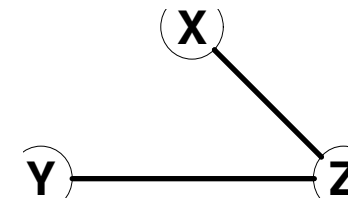
## Representation of conditional independency

## Example 6 (cont.).

$Z$	$X$	$Y$	$p(X, Y, Z)$
0	0	0	0.024
0	0	1	0.056
0	1	0	0.036
0	1	1	0.084
1	0	0	0.096
1	0	1	0.144
1	1	0	0.224
1	1	1	0.336

Checking  $p \cdot p^{\downarrow V \setminus \{X, Y\}} = p^{\downarrow V \setminus \{X\}} \cdot p^{\downarrow V \setminus \{Y\}}$  one finds that the only independency relations of  $p$  are  $I_p(X, Y|Z)$  and  $I_p(Y, X|Z)$ .

Thus, the graph



represents  $p$ , as its independency model is  $I_G := \{(X, Y|Z), (Y, X|Z)\}$ .

As for  $p$  only  $I_p(X, Y|Z)$  and  $I_p(Y, X|Z)$  hold,  $G$  is a faithful representation.



## Markov networks

**Definition 16.** A pair  $(G, (\psi_C)_{C \in \mathcal{C}_G})$  consisting of

- (i) an undirected graph  $G$  on a set of variables  $V$  and
- (ii) a set of potentials

$$\psi_C : \prod_{X \in C} \text{dom}(X) \rightarrow \mathbb{R}_0^+, \quad C \in \mathcal{C}_G$$

on the cliques<sup>1)</sup> of  $G$  (called **clique potentials**)

is called a **markov network**.

<sup>1)</sup> on the product of the domains of the variables of each clique.

Thus, a markov network encodes

- (i) a joint probability distribution factorized as

$$p = \left( \prod_{C \in \mathcal{C}_G} \psi_C \right)^{|\emptyset|}$$

and

- (ii) conditional independency statements

$$I_G(X, Y | Z) \Rightarrow I_p(X, Y | Z)$$

$G$  represents  $p$ , but not necessarily faithfully.



## Markov networks / examples

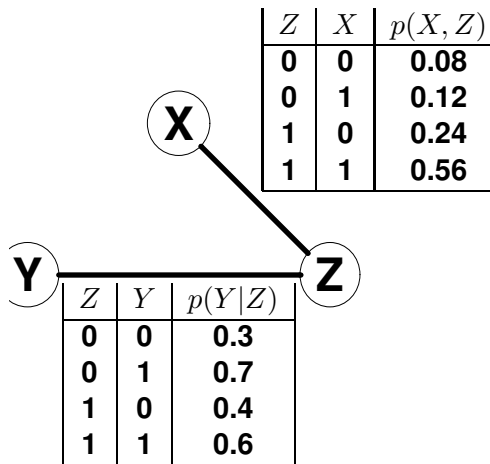


Figure 30: Example for a markov network.



- 1. Basic Probability Calculus**
- 2. Separation in undirected graphs**
- 3. Separation in directed graphs**
- 4. Markov networks**
- 5. Bayesian networks**



## DAG-representations

**Lemma 4 (criterion for DAG-representation).** *Let  $p$  be a joint probability distribution of the variables  $V$  and  $G$  be a graph on the vertices  $V$ . Then:*

*$G$  represents  $p \Leftrightarrow v$  and  $\text{nondesc}(v)$  are conditionally independent given  $\text{pa}(v)$  for all  $v \in V$ , i.e.,*

$$I_p(\{v\}, \text{nondesc}(v) \mid \text{pa}(v)), \quad \forall v \in V$$

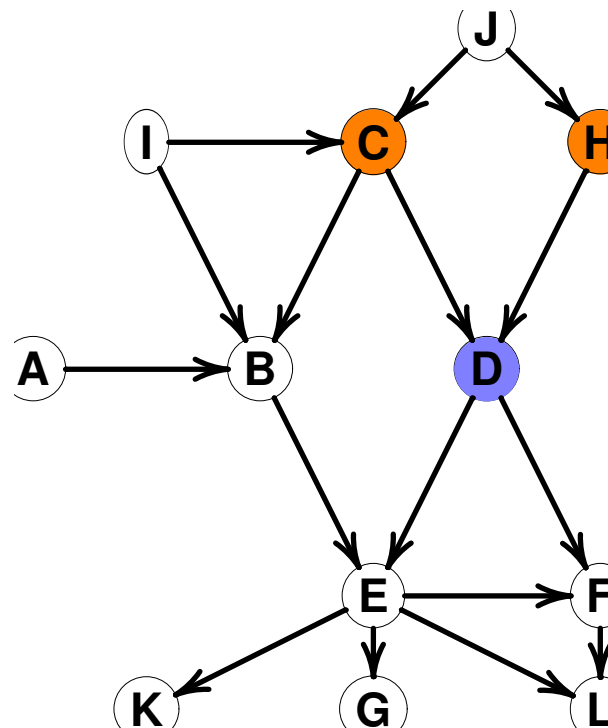


Figure 31: Parents of a vertex (orange).



## Example for a not faithfully DAG-representable independency model

Probability distributions may have no faithful DAG-representation.

### **Example 7.** The independency model

$$I := \{I(x, y|z), I(y, x|z), I(x, y|w), I(y, x|w)\}$$

does not have a faithful DAG-representation. [CGH97, p. 239]

Exercise: compute all minimal DAG-representations of  $I$  using lemma 5 and check if they are faithful.



## Minimal DAG-representations

**Lemma 5 (construction and uniqueness of minimal DAG-representation, [VP90])**

*Let  $I$  be an independence model of a JPD  $p$ . Then:*

- (i) A minimal DAG-representation can be constructed as follows: Choose an arbitrary ordering  $\sigma := (v_1, \dots, v_n)$  of  $V$ . Choose a minimal set  $\pi_i \subseteq \{v_1, \dots, v_{i-1}\}$  of  $\sigma$ -precursors of  $v_i$  with*

$$I(v_i, \{v_1, \dots, v_{i-1}\} \setminus \pi_i | \pi_i)$$

*Then  $G := (V, E)$  with*

$$E := \{(w, v_i) \mid i = 1, \dots, n, w \in \pi_i\}$$

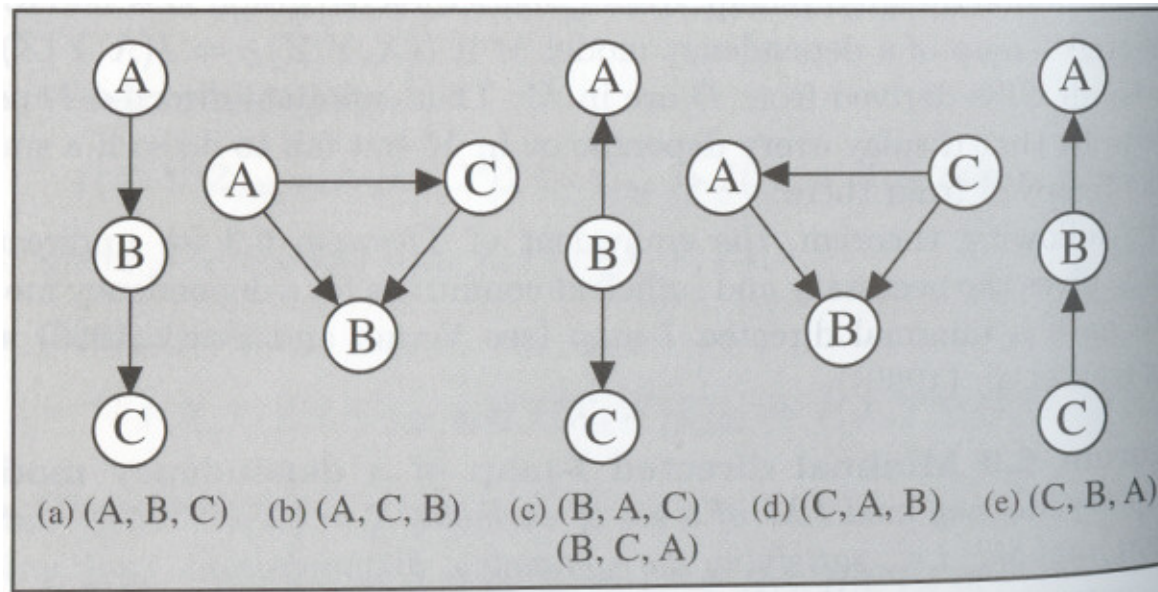
*is a minimal DAG-representation of  $p$ .*

- (ii) If  $p$  also is non-extreme, then the minimal representation  $G$  is unique up to ordering  $\sigma$ .*



## Minimal DAG-representations / example

$$I := \{(A, C|B), (C, A|B)\}$$

Figure 32: Minimal DAG-representations of  $I$  [CGH97, p. 240].



## Minimal representations / conclusion

Representations always exist (e.g., trivial).

Minimal representations always exist  
(e.g., start with trivial and drop edges successively).

	Markov network (undirected)		Bayesian network (directed)	
	minimal	faithful	minimal	faithful
general JPD	may not be unique	may not exist	may not be unique	may not exist
non-extreme JPD	unique	may not exist	unique up to ordering	may not exist



## Bayesian Network

**Definition 17.** A pair  $(G := (V, E), (p_v)_{v \in V})$  consisting of

- (i) a directed graph  $G$  on a set of variables  $V$  and
- (ii) a set of conditional probability distributions

$$p_X : \text{dom}(X) \times \prod_{Y \in \text{pa}(X)} \text{dom}(Y) \rightarrow \mathbb{R}_0^+$$

at the vertices  $X \in V$  conditioned on its parents (called **(conditional) vertex probability distributions**)

is called a **bayesian network**.

Thus, a bayesian network encodes

- (i) a joint probability distribution factorized as

$$p = \prod_{X \in V} p(X | \text{pa}(X))$$

and

- (ii) conditional independency statements

$$I_G(X, Y | Z) \Rightarrow I_p(X, Y | Z)$$

$G$  represents  $p$ , but not necessarily faithfully.

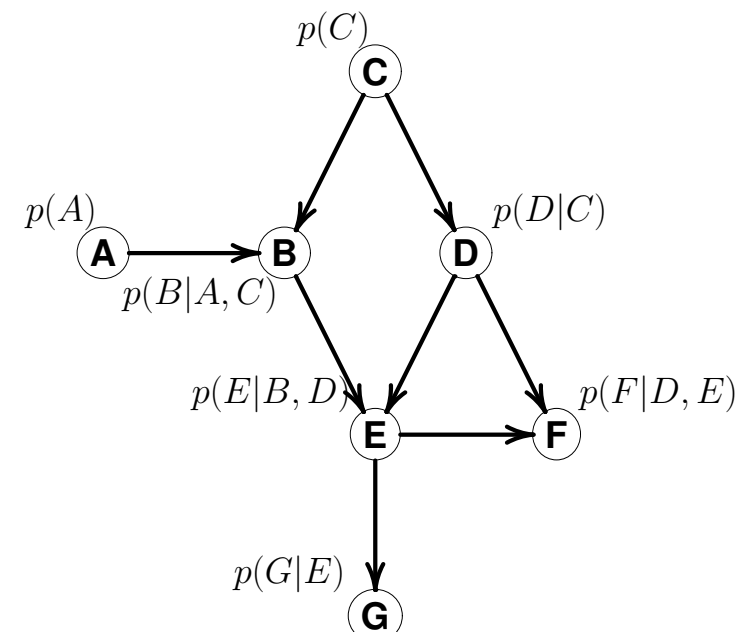


Figure 33: Example for a bayesian network.



## Types of probabilistic networks

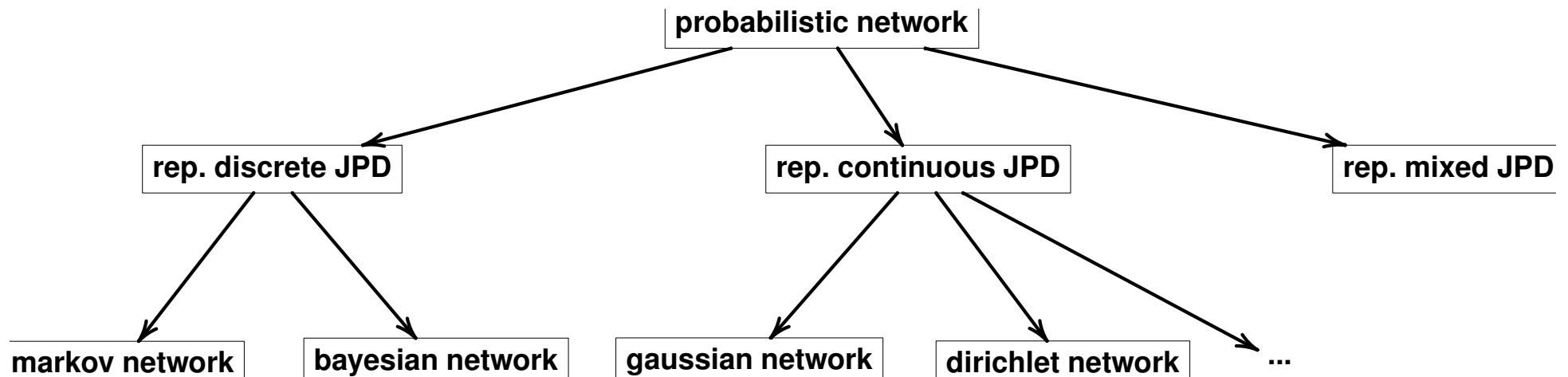


Figure 34: Types of probabilistic networks.



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