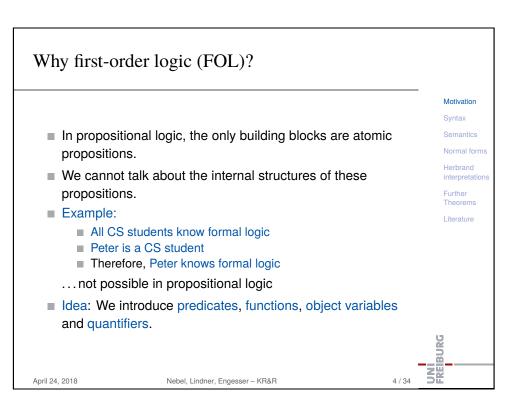
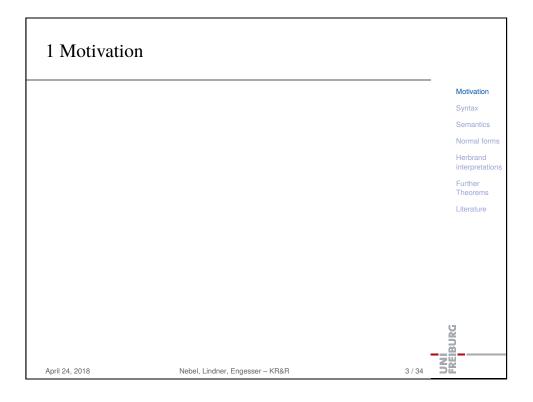
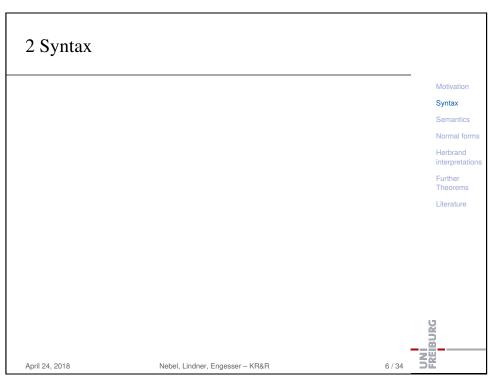
Principles of Knowledge Representation and Reasoning Predicate logic Bernhard Nebel, Felix Lindner, and Thorsten Engesser April 24, 2018







Syntax

- \blacksquare variable symbols: x, y, z, ... \blacksquare *n*-ary function symbols: f, g, ...
- constant symbols: a,b,c,...
- \blacksquare *n*-ary predicate symbols: P, Q, ...
- logical symbols: \forall , \exists , =, \neg , \wedge , ...

Terms

t := xvariable function application constant

Formulae

 $::= P(t_1,\ldots,t_n)$ atomic formulae t = t'identity formulae propositional connectives $\forall x \phi'$ universal quantification $\exists x \omega'$ existential quantification

Ground term, etc.: term, etc. without variable occurrences

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Semantics: idea

- In FOL, the universe of discourse consists of objects: we consider functions and relations over these objects.
- Function symbols are mapped to functions, predicate symbols are mapped to relations, and terms to objects.
- Notation: Instead of $\mathcal{I}(x)$ we write $x^{\mathcal{I}}$.
- Note: Usually one considers all possible non-empty universes. (However, sometimes the interpretations are restricted to particular domains, e.g. integers or real numbers.)
- Satisfiability and validity is then considered wrt. all these universes.

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Formal semantics: interpretations

Interpretations: $\mathcal{I} = \langle \mathcal{D}, \cdot^{\mathcal{I}} \rangle$ with \mathcal{D} being an arbitrary non-empty set and $\cdot^{\mathcal{I}}$ being a function which maps

- \blacksquare *n*-ary function symbols *f* to *n*-ary functions $f^{\mathcal{I}} \in [\mathcal{D}^n \to \mathcal{D}]$,
- \blacksquare constant symbols a to objects $a^{\mathcal{I}} \in \mathcal{D}$, and
- *n*-ary predicates *P* to *n*-ary relations $P^{\mathcal{I}} \subseteq \mathcal{D}^n$.

Interpretation of ground terms:

$$(f(t_1,\ldots,t_n))^{\mathcal{I}} = f^{\mathcal{I}}(t_1^{\mathcal{I}},\ldots,t_n^{\mathcal{I}}) \ (\in \mathcal{D})$$

Truth of ground atoms:

$$\mathcal{I} \models P(t_1, \dots, t_n) \quad \text{iff} \quad \langle t_1^{\mathcal{I}}, \dots, t_n^{\mathcal{I}} \rangle \in P^{\mathcal{I}}$$

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Examples

 $Cat^{\mathcal{I}} = \{d_1\}$

 $\mathcal{I} \models Red(b)$

Cat(b)

 $Red^{\mathcal{I}} = \mathcal{D}$

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 $\mathcal{I} \models \text{even}(\text{succ}(3))$

 $\operatorname{succ}^{\mathcal{I}} = \{(1 \mapsto 2), (2 \mapsto 3), \dots\}$

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 $even^{\mathcal{I}} = \{2,4,6,\ldots\}$

 $\mathcal{I} \not\models \text{even(3)}$

 $\mathcal{D} = \{d_1, ..., d_n\}, n \ge 2$ $\mathcal{D} = \{1, 2, 3, ...\}$

Formal semantics: variable assignments

V is the set of variables. Functions $\alpha: V \to \mathcal{D}$ are called variable assignments.

Notation: $\alpha[x/d]$ is identical to α except for x where $\alpha[x/d](x) = d.$

Interpretation of terms under \mathcal{I}, α :

$$x^{\mathcal{I},\alpha} = \alpha(x)$$

$$a^{\mathcal{I},\alpha} = a^{\mathcal{I}}$$

$$(f(t_1,\ldots,t_n))^{\mathcal{I},\alpha} = f^{\mathcal{I}}(t_1^{\mathcal{I},\alpha},\ldots,t_n^{\mathcal{I},\alpha})$$

Truth of atomic formulae:

$$\mathcal{I}, \alpha \models P(t_1, \dots, t_n) \quad \text{iff} \quad \langle t_1^{\mathcal{I}, \alpha}, \dots, t_n^{\mathcal{I}, \alpha} \rangle \in P^{\mathcal{I}}$$

Example (cont'd):

$$\alpha = \{x \mapsto d_1, y \mapsto d_2\}$$
 $\mathcal{I}, \alpha \models \text{Red}(x)$ $\mathcal{I}, \alpha[y/d_1] \models \text{Cat}(y)$

$$\mathcal{I}, \alpha \models \text{Red}(x)$$

$$\mathcal{I}, \alpha[y/d_1] \models \text{Cat}(y)$$

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Formal semantics: truth

Truth of φ under \mathcal{I} and α ($\mathcal{I}, \alpha \models \varphi$) is defined as follows.

$$\mathcal{I}, \alpha \models P(t_1, \dots, t_n) & \text{iff } \langle t_1^{\mathcal{I}, \alpha}, \dots, t_n^{\mathcal{I}, \alpha} \rangle \in P^{\mathcal{I}} \\
\mathcal{I}, \alpha \models t_1 = t_2 & \text{iff } t_1^{\mathcal{I}, \alpha} = t_2^{\mathcal{I}, \alpha} \\
\mathcal{I}, \alpha \models \neg \varphi & \text{iff } \mathcal{I}, \alpha \not\models \varphi \\
\mathcal{I}, \alpha \models \varphi \land \psi & \text{iff } \mathcal{I}, \alpha \models \varphi \text{ and } \mathcal{I}, \alpha \models \psi \\
\mathcal{I}, \alpha \models \varphi \lor \psi & \text{iff } \mathcal{I}, \alpha \models \varphi \text{ or } \mathcal{I}, \alpha \models \psi \\
\mathcal{I}, \alpha \models \varphi \to \psi & \text{iff } if \mathcal{I}, \alpha \models \varphi, \text{then } \mathcal{I}, \alpha \models \psi \\
\mathcal{I}, \alpha \models \varphi \leftrightarrow \psi & \text{iff } \mathcal{I}, \alpha \models \varphi \text{ iff } \mathcal{I}, \alpha \models \psi \\
\mathcal{I}, \alpha \models \forall x \varphi & \text{iff } \mathcal{I}, \alpha[x/d] \models \varphi \text{ for all } d \in \mathcal{D} \\
\mathcal{I}, \alpha \models \exists x \varphi & \text{iff } \mathcal{I}, \alpha[x/d] \models \varphi \text{ for some } d \in \mathcal{D}$$

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Questions:

$$\mathcal{D} = \{d_1, \dots, d_n\}, \ n > 1$$

$$a^{\mathcal{I}} = d_1$$

$$b^{\mathcal{I}} = d_1$$

$$\operatorname{Cat}^{\mathcal{I}} = \{d_1\}$$

$$\operatorname{Red}^{\mathcal{I}} = \mathcal{D}$$

$$\alpha = \{(x \mapsto d_1), (y \mapsto d_2)\}$$

$$\mathcal{I},$$

$$\mathcal{I}, \alpha \models \mathsf{Cat}(b) \lor \neg \mathsf{Cat}(b)$$
?
Yes
 $\mathcal{I}, \alpha \models \mathsf{Cat}(x) \rightarrow$

$$\mathcal{I}, \alpha \models \operatorname{Cat}(x) \rightarrow \operatorname{Cat}(x) \vee \operatorname{Cat}(y)$$
?

$$\mathcal{I}, \alpha \models \mathsf{Cat}(x) \to \mathsf{Cat}(y)$$
?

 $\mathcal{I}, \alpha \models \mathsf{Cat}(a) \land \mathsf{Cat}(b)$?

$$\mathcal{I}, \alpha \models \forall x (Cat(x) \rightarrow Red(x))$$
?

$$Red(x)$$
?

$$\mathcal{I}, \alpha \models \Theta$$
? Yes

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 \mathcal{I}, α is a model of φ iff

$$\mathcal{I}, \alpha \models \varphi$$
.

A formula can be satisfiable, unsatisfiable, falsifiable, valid, ... Formulae φ and ψ are logically equivalent (symb.: $\varphi \equiv \psi$) iff for all \mathcal{I}, α :

$$\mathcal{I}, \alpha \models \varphi \text{ iff } \mathcal{I}, \alpha \models \psi.$$

Note: $P(x) \not\equiv P(y)!$

Logical implication is also analogous to propositional logic:

$$\Theta \models \varphi$$
 iff for all \mathcal{I}, α s.t. $\mathcal{I}, \alpha \models \Theta$ also $\mathcal{I}, \alpha \models \varphi$.

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Open & closed formulae

- Formulae without free variables are called closed formulae or sentences. Formulae with free variables are called open formulae.
- Closed formulae are all we need when we want to state something about the world. Open formulae (and variable assignments) are only necessary for technical reasons (semantics of \forall and \exists).
- Note that logical equivalence, satisfiability, and entailment are independent from variable assignments if we consider only closed formulae.
- For closed formulae, we omit α in connection with \models :

Free and bound variables

Variables can be free or bound (by a quantifier) in a formula:

 $free(x) = \{x\}$

 $free(f(t_1, ..., t_n)) = free(t_1) \cup \cdots \cup free(t_n)$

 $free(t_1 = t_2) = free(t_1) \cup free(t_2)$

 $free(P(t_1, ..., t_n)) = free(t_1) \cup \cdots \cup free(t_n)$

 $free(\neg \phi) = free(\phi)$

 $free(\phi * \psi) = free(\phi) \cup free(\psi), for * = \lor, \land, \rightarrow, \leftrightarrow$

 $free(Qx\varphi) = free(\varphi) \setminus \{x\}, for Q = \forall, \exists$

Example: $\forall x (R(y,z) \land \exists y (\neg P(y,x) \lor R(y,z)))$

Which occurrences are free, which are not free?

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 $\mathcal{I} \models \varphi$.

Prenex Normal Form

The prenex normal form of a FOL formula has the following form:

quantifier prefix + (quantifier free) matrix

Generate prenex normal form:

- II Eliminate \rightarrow and \leftrightarrow .
- 2 Move ¬ inside.
- Moving quantifiers out (using a number of equivalences).

Theorem

For each FOL formula, an equivalent formula in prenex normal form exists and can be effectively computed.

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Skolemization

We can further simplify formulae by eliminating existential quantifiers using fresh function symbols (Skolem functions).

Theorem (Skolem normal form)

Let φ be a closed formula in prenex normal form with all variables pairwise distinct of the form $\varphi = \forall x_1 \dots \forall x_i \exists y \psi$. Let g_i be an i-ary function symbols not appearing in φ . Then φ is satisfiable iff

$$\varphi' = \forall x_1 \dots \forall x_i \psi[y/g_i(x_1, \dots, x_i)]$$

is satisfiable.

Proof idea.

For each assignment to $x_1 ldots x_i$, there is a value of $y[=g(x_1, ldots, x_i)]$ and vice versa.

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Skolem normal form

Skolem Normal Form

Prenex normal form without existential quantifiers.

Notation: ϕ^* is SNF of ϕ

Theorem

For each closed formula ϕ , a corresponding SNF ϕ^* can be effectively computed.

Example

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 $\exists x ((\forall x p(x)) \land \neg q(x))$ $\exists y ((\forall x p(x)) \land \neg q(y))$ $\exists y (\forall x (p(x) \land \neg q(y)))$ $\forall x (p(x) \land \neg q(g_0))$

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Note: SNF is not unique

Reducing FOL satisfiability to propositional satisfiability ...

Idea 1: We use one particular interpretation which has as the universe of discourse all possible ground terms – and we add one constant if we do not have already one \rightsquigarrow Herbrand universe

Example:
$$\forall x \forall y (\neg P(x,y) \lor R(g_2(x,y),x))$$

 $\mathcal{D}^H = \{a_0, g_2(a_0, a_0), g_2(a_0, g_2(a_0, a_0)), \dots \}$

Idea 2: Function symbols are interpreted syntactically, predicate symbols are interpreted arbitrarily over this universe (each ground atom gets a truth value):

Herbrand interpretation

$$a^{\mathcal{I}} = a$$

 $(f(t_1, \dots, t_n))^{\mathcal{I}} = f(t_1, \dots, t_n)$

 \mathcal{I} could then be defined such that, e.g., $\mathcal{I} \not\models P(a_0, a_0)$, $\mathcal{I} \not\models P(a_0, g_2(a_0, a_0))$, etc.

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A reduction to a satisfiability problem with infinitely many formulae

- Note that the Herbrand universe can be infinite, therefore $E(\varphi)$ can be infinite!
- $\hfill\blacksquare$ If the Herbrand base is finite there is no problem (well, $\ldots)$

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- lacktriangle Use $E(\phi)$ in a "lazy" way, expand only as needed
- Semi-decision method for unsatisfiability
- In fact, unsatisfiability (and validity) in FOL is only semi-decidable (use e.g. PCP to prove)!

Herbrand models and Herbrand expansions

Theorem

A formula φ has a model iff it has a Herbrand model.

Idea 3: We expand each SNF-formula by substituting all variables by all possible terms \rightsquigarrow Herbrand expansion $(E(\varphi))$

Example: $\neg P(a_0, a_0) \lor R(g_2(a_0, a_0), a_0), \neg P(a_0, g_2(a_0, a_0)) \lor R(g_2(a_0, g_2(a_0, a_0)), a_0), \dots$

Theorem

A formula φ is satisfiable if $E(\varphi)$ is satisfiable.

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Further theorems

Some corollaries from the previous theorems:

Theorem (Compactness)

Let $\Phi \cup \{\psi\}$ be a set of closed formulae.

- (a) $\Phi \models \psi$ iff there exists a finite subset $\Phi' \subseteq \Phi$ s. t. $\Phi' \models \psi$.
- (b) Φ is satisfiable iff each finite subset $\Phi' \subseteq \Phi$ is satisfiable.

Theorem (Löwenheim-Skolem)

Each countable set of closed formulae that is satisfiable is satisfiable on a countable domain.

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