Why Logic?
Why logic?

- Logic is one of the best developed systems for representing knowledge.
- Can be used for analysis, design and specification.
- Understanding formal logic is a prerequisite for understanding most research papers in KR&R.
The right logic…

- Logics of different orders (1st, 2nd, ...)
- Modal logics
  - epistemic
  - temporal
  - dynamic (program)
  - multi-modal logics
  - ...
- Many-valued logics
- Nonmonotonic logics
- Intuitionistic logics
- ...

Why Logic?
Propositional Logic
Syntax
Semantics
Terminology
Decision Problems and Resolution
The logical approach

- Define a **formal language**: logical & non-logical symbols, syntax rules
The logical approach

- Define a **formal language**: logical & non-logical symbols, syntax rules
- Provide language with **compositional semantics**:
  - Fix **universe** of discourse
  - Specify how the non-logical symbols can be **interpreted**: interpretation
  - Rules how to **combine** interpretation of single symbols
  - **Satisfying interpretation** = model
  - Semantics often entails concept of **logical implication** / entailment
The logical approach

- Define a **formal language**: logical & non-logical symbols, syntax rules
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  - Rules how to combine interpretation of single symbols
  - Satisfying interpretation = model
  - Semantics often entails concept of logical implication / entailment
- Specify a **calculus** that allows to derive new formulae from old ones – according to the entailment relation
Propositional Logic
Propositional logic: main ideas

- **Non-logical symbols**: propositional variables or atoms
  - representing propositions which cannot be decomposed
  - which can be true or false (for example: “Snow is white”, “It rains”)
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Propositional logic: main ideas

- **Non-logical symbols**: propositional *variables* or *atoms*
  - representing *propositions* which cannot be decomposed
  - which can be *true* or *false* (for example: “Snow is white”, “It rains”)
- **Logical symbols**: propositional connectives such as: and ($\land$), or ($\lor$), and not ($\neg$)
- **Formulae**: built out of atoms and connectives
- **Universe of discourse**: truth values
Syntax
Countable alphabet $\Sigma$ of propositional variables: $a, b, c, \ldots$

Propositional formulae are built according to the following rule:

$$\varphi ::= a \quad \text{atomic formula}$$

$$\quad \downarrow \quad \text{falsity}$$

$$\quad \top \quad \text{truth}$$

$$\neg \varphi' \quad \text{negation}$$

$$(\varphi' \land \varphi'') \quad \text{conjunction}$$

$$(\varphi' \lor \varphi'') \quad \text{disjunction}$$

$$(\varphi' \rightarrow \varphi'') \quad \text{implication}$$

$$(\varphi' \leftrightarrow \varphi'') \quad \text{equivalence}$$

Parentheses can be omitted if no ambiguity arises.

Operator precedence: $\neg > \land > \lor > \rightarrow = \leftrightarrow$. 
### Syntax

Countable alphabet $\Sigma$ of **propositional variables**: $a, b, c, \ldots$

**Propositional formulae** are built according to the following **rule**:

<table>
<thead>
<tr>
<th>$\varphi$</th>
<th>$a$</th>
<th>atomic formula</th>
</tr>
</thead>
<tbody>
<tr>
<td>$\bot$</td>
<td>falsity</td>
<td></td>
</tr>
<tr>
<td>$\top$</td>
<td>truth</td>
<td></td>
</tr>
<tr>
<td>$\neg \varphi'$</td>
<td>negation</td>
<td></td>
</tr>
<tr>
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**Operator precedence**: $\neg > \land > \lor > \rightarrow = \leftrightarrow$. 
Language and meta-language

- $(a \lor b)$ is an expression of the language of propositional logic.
- $\varphi ::= a \mid \ldots \mid ((\varphi' \leftrightarrow \varphi''))$ is a statement about how expressions in the language of propositional logic can be formed. It is stated using meta-language.
- In order to describe how expressions (in this case formulae) can be formed, we use meta-language.
- When we describe how to interpret formulae, we use meta-language expressions.
Semantics
Semantics: idea

- Atomic propositions can be true (1, T) or false (0, F).
- Provided the truth values of the atoms have been fixed (truth assignment or interpretation), the truth value of a formula can be computed from the truth values of the atoms and the connectives.

Example:

\[(a \lor b) \land c\]

is true iff c is true and, additionally, a or b is true.
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**Example:**

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is true iff \( c \) is true and, additionally, \( a \) or \( b \) is true.

Logical implication can then be defined as follows:

- \( \phi \) is implied by a set of formulae \( \Theta \) iff \( \phi \) is true for all truth assignments (world states) that make all formulae in \( \Theta \) true.
Formal semantics

An interpretation (or truth assignment) over $\Sigma$ is a function:

$$I: \Sigma \rightarrow \{T, F\}.$$
Formal semantics

An interpretation (or truth assignment) over $\Sigma$ is a function:

$$\mathcal{I} : \Sigma \rightarrow \{T, F\}.$$ 

A formula $\psi$ is true under $\mathcal{I}$ or is satisfied by $\mathcal{I}$ (symb. $\mathcal{I} \models \psi$):

- $\mathcal{I} \models a$ iff $\mathcal{I}(a) = T$
- $\mathcal{I} \models \top$
- $\mathcal{I} \not\models \bot$
- $\mathcal{I} \models \neg \phi$ iff $\mathcal{I} \not\models \phi$
- $\mathcal{I} \models \phi \land \phi'$ iff $\mathcal{I} \models \phi$ and $\mathcal{I} \models \phi'$
- $\mathcal{I} \models \phi \lor \phi'$ iff $\mathcal{I} \models \phi$ or $\mathcal{I} \models \phi'$
- $\mathcal{I} \models \phi \rightarrow \phi'$ iff if $\mathcal{I} \models \phi$ then $\mathcal{I} \models \phi'$
- $\mathcal{I} \models \phi \leftrightarrow \phi'$ iff $\mathcal{I} \models \phi$ if and only if $\mathcal{I} \models \phi'$
Example

Given

\[ I : a \mapsto T, \; b \mapsto F, \; c \mapsto F, \; d \mapsto T, \]

Is \((a \lor b) \leftrightarrow (c \lor d)) \land (\neg(a \land c) \lor (c \land \neg d))\) true or false?
Example

Given

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Terminology
Terminology

An interpretation $\mathcal{I}$ is a **model** of $\varphi$ iff $\mathcal{I} \models \varphi$.

A formula $\varphi$ is

- **satisfiable** if there is an $\mathcal{I}$ such that $\mathcal{I} \models \varphi$;
- **unsatisfiable**, otherwise; and
- **valid** if $\mathcal{I} \models \varphi$ for each $\mathcal{I}$ (or **tautology**);
- **falsifiable**, otherwise.
An interpretation $\mathcal{I}$ is a **model** of $\varphi$ iff $\mathcal{I} \models \varphi$.

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Formulae $\varphi$ and $\psi$ are **logically equivalent** (symb. $\varphi \equiv \psi$) if for all interpretations $\mathcal{I}$,

$$\mathcal{I} \models \varphi \text{ iff } \mathcal{I} \models \psi.$$
Examples

Satisfiable, unsatisfiable, falsifiable, valid?

\((a \lor b \lor \neg c) \land (\neg a \lor \neg b \lor d) \land (\neg a \lor b \lor \neg d)\)
Examples

Satisfiable, unsatisfiable, falsifiable, valid?

\((a \lor b \lor \neg c) \land (\neg a \lor \neg b \lor d) \land (\neg a \lor b \lor \neg d)\)

\(\Rightarrow\) satisfiable: \(a \mapsto T, b \mapsto F, d \mapsto F, \ldots\)
Examples

Satisfiable, unsatisfiable, falsifiable, valid?

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⇒ satisfiable: \(a \mapsto T, b \mapsto F, d \mapsto F, \ldots\)

⇒ falsifiable: \(a \mapsto F, b \mapsto F, c \mapsto T, \ldots\)
Examples

Satisfiable, unsatisfiable, falsifiable, valid?

\[(a \lor b \lor \neg c) \land (\neg a \lor \neg b \lor d) \land (\neg a \lor b \lor \neg d)\]

\[\implies \text{satisfiable: } a \mapsto T, b \mapsto F, d \mapsto F, \ldots\]

\[\implies \text{falsifiable: } a \mapsto F, b \mapsto F, c \mapsto T, \ldots\]

\[((\neg a \rightarrow \neg b) \rightarrow (b \rightarrow a))\]
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Equivalence? \(\neg (a \lor b) \equiv \neg a \land \neg b\)
Examples

Satisfiable, unsatisfiable, falsifiable, valid?

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\(\Rightarrow\) valid: Consider all interpretations or argue about falsifying ones.

Equivalence? \(\neg(a \lor b) \equiv \neg a \land \neg b\)

\(\Rightarrow\) Of course, equivalent (de Morgan).
Some obvious consequences

Proposition

φ is valid iff ¬φ is unsatisfiable.
φ is satisfiable iff ¬φ is falsifiable.
Some obvious consequences

**Proposition**

\( \phi \) is valid iff \( \neg \phi \) is unsatisfiable.

\( \phi \) is satisfiable iff \( \neg \phi \) is falsifiable.

**Proposition**

\( \phi \equiv \psi \) iff \( \phi \leftrightarrow \psi \) is valid.
Some obvious consequences

**Proposition**

\[ \phi \text{ is valid iff } \neg \phi \text{ is unsatisfiable.} \]
\[ \phi \text{ is satisfiable iff } \neg \phi \text{ is falsifiable.} \]

**Proposition**

\[ \phi \equiv \psi \text{ iff } \phi \leftrightarrow \psi \text{ is valid.} \]

**Theorem**

*If* \( \phi \equiv \psi \), *and* \( \chi' \) *results from substituting* \( \phi \) *by* \( \psi \) *in* \( \chi \), *then* \( \chi' \equiv \chi \).
### Some equivalences

<table>
<thead>
<tr>
<th>Property</th>
<th>Expression</th>
<th>Simplified Expression</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Simplifications</strong></td>
<td>$\varphi \rightarrow \psi$</td>
<td>$\neg \varphi \lor \psi$</td>
</tr>
<tr>
<td></td>
<td>$\varphi \leftrightarrow \psi$</td>
<td>$(\varphi \rightarrow \psi) \land (\psi \rightarrow \varphi)$</td>
</tr>
<tr>
<td><strong>Idempotency</strong></td>
<td>$\varphi \lor \varphi$</td>
<td>$\varphi$</td>
</tr>
<tr>
<td></td>
<td>$\varphi \land \varphi$</td>
<td>$\varphi$</td>
</tr>
<tr>
<td><strong>Commutativity</strong></td>
<td>$\varphi \lor \psi$</td>
<td>$\psi \lor \varphi$</td>
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<tr>
<td></td>
<td>$\varphi \land \psi$</td>
<td>$\psi \land \varphi$</td>
</tr>
<tr>
<td><strong>Associativity</strong></td>
<td>$(\varphi \lor \psi) \lor \chi$</td>
<td>$\varphi \lor (\psi \lor \chi)$</td>
</tr>
<tr>
<td></td>
<td>$(\varphi \land \psi) \land \chi$</td>
<td>$\varphi \land (\psi \land \chi)$</td>
</tr>
<tr>
<td><strong>Absorption</strong></td>
<td>$\varphi \lor (\varphi \land \psi)$</td>
<td>$\varphi$</td>
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<tr>
<td></td>
<td>$\varphi \land (\varphi \lor \psi)$</td>
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<tr>
<td><strong>Distributivity</strong></td>
<td>$\varphi \land (\psi \lor \chi)$</td>
<td>$(\varphi \land \psi) \lor (\varphi \land \chi)$</td>
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</tr>
<tr>
<td><strong>Double Negation</strong></td>
<td>$\neg \neg \varphi$</td>
<td>$\varphi$</td>
</tr>
<tr>
<td>** Constants**</td>
<td>$\neg \top$</td>
<td>$\bot$</td>
</tr>
<tr>
<td></td>
<td>$\neg \bot$</td>
<td>$\top$</td>
</tr>
<tr>
<td><strong>De Morgan</strong></td>
<td>$\neg (\varphi \lor \psi)$</td>
<td>$\neg \varphi \land \neg \psi$</td>
</tr>
<tr>
<td></td>
<td>$\neg (\varphi \land \psi)$</td>
<td>$\neg \varphi \lor \neg \psi$</td>
</tr>
<tr>
<td><strong>Truth</strong></td>
<td>$\varphi \lor \top$</td>
<td>$\top$</td>
</tr>
<tr>
<td></td>
<td>$\varphi \land \top$</td>
<td>$\varphi$</td>
</tr>
<tr>
<td><strong>Falsity</strong></td>
<td>$\varphi \lor \bot$</td>
<td>$\varphi$</td>
</tr>
<tr>
<td></td>
<td>$\varphi \land \bot$</td>
<td>$\bot$</td>
</tr>
<tr>
<td><strong>Taut./Contrad.</strong></td>
<td>$\varphi \lor \neg \varphi$</td>
<td>$\top$</td>
</tr>
<tr>
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<td>$\varphi \land \neg \varphi$</td>
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</table>
How many different formulae are there …

… for a given finite alphabet $\Sigma$?
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- Infinitely many: $a, a \lor a, a \land a, a \lor a \lor a, \ldots$
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  - A formula can be characterized by its set of models (if two formulae are not logically equivalent, then their sets of models differ).
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  - A formula can be characterized by its set of models (if two formulae are not logically equivalent, then their sets of models differ).
  - For \( \Sigma \) with \( n = |\Sigma| \), there are \( 2^n \) different interpretations.
  - There are \( 2^{(2^n)} \) different sets of interpretations.
  - There are \( 2^{(2^n)} \) (logical) equivalence classes of formulae.
Logical implication

- Extension of the relation $\models$ to sets $\Theta$ of formulae:

$$\mathcal{I} \models \Theta \iff \mathcal{I} \models \varphi \text{ for all } \varphi \in \Theta.$$
Logical implication

- Extension of the relation $\models$ to sets $\Theta$ of formulae:

  $$\mathcal{I} \models \Theta \iff \mathcal{I} \models \varphi \text{ for all } \varphi \in \Theta.$$ 

- $\varphi$ is logically implied by $\Theta$ (symbolically $\Theta \models \varphi$) iff $\varphi$ is true in all models of $\Theta$:

  $$\Theta \models \varphi \iff \mathcal{I} \models \varphi \text{ for all } \mathcal{I} \text{ such that } \mathcal{I} \models \Theta.$$
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Some consequences:
- **Deduction theorem:** $\Theta \cup \{ \varphi \} \models \psi$ iff $\Theta \models \varphi \rightarrow \psi$
Logical implication

- Extension of the relation $|$ to sets $\Theta$ of formulae:
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- **Deduction theorem**: $\Theta \cup \{\varphi\} \models \psi$ iff $\Theta \models \varphi \rightarrow \psi$
- **Contraposition**: $\Theta \cup \{\varphi\} \models \neg \psi$ iff $\Theta \cup \{\psi\} \models \neg \varphi$
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- Extension of the relation $\models$ to sets $\Theta$ of formulae:

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- **Deduction theorem:** $\Theta \cup \{\phi\} \models \psi$ iff $\Theta \models \phi \rightarrow \psi$  
- **Contraposition:** $\Theta \cup \{\phi\} \models \neg \psi$ iff $\Theta \cup \{\psi\} \models \neg \phi$  
- **Contradiction:** $\Theta \cup \{\phi\}$ is unsatisfiable iff $\Theta \models \neg \phi$
Normal forms

**Terminology:**

- Atomic formulae $a$, negated atomic formulae $\neg a$, truth $\top$ and falsity $\bot$ are **literals**.
- A disjunction of literals is a **clause**.
Normal forms

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- Atomic formulae \( a \), negated atomic formulae \( \neg a \), truth \( \top \) and falsity \( \bot \) are **literals**.
- A disjunction of literals is a **clause**.
- If \( \neg \) only occurs in front of an atom and there are no \( \rightarrow \) and \( \leftrightarrow \), the formula is in **negation normal form** (NNF).

**Example:** \((\neg a \lor \neg b) \land c\), but **not**: \(\neg(a \land b) \land c\)
Normal forms

Terminology:

- Atomic formulae $a$, negated atomic formulae $\neg a$, truth $\top$ and falsity $\bot$ are literals.
- A disjunction of literals is a clause.
- If $\neg$ only occurs in front of an atom and there are no $\rightarrow$ and $\leftrightarrow$, the formula is in negation normal form (NNF).
  Example: $(\neg a \lor \neg b) \land c$, but not: $\neg (a \land b) \land c$
- A conjunction of clauses is in conjunctive normal form (CNF).
  Example: $(a \lor b) \land (\neg a \lor c)$
Normal forms

Terminology:

- Atomic formulae \( a \), negated atomic formulae \( \neg a \), truth \( \top \) and falsity \( \bot \) are literals.
- A disjunction of literals is a clause.
- If \( \neg \) only occurs in front of an atom and there are no \( \rightarrow \) and \( \leftrightarrow \), the formula is in negation normal form (NNF).
  Example: \((\neg a \lor \neg b) \land c\), but not: \(\neg(a \land b) \land c\)
- A conjunction of clauses is in conjunctive normal form (CNF).
  Example: \((a \lor b) \land (\neg a \lor c)\)
- The dual form (disjunction of conjunctions of literals) is in disjunctive normal form (DNF).
  Example: \((a \land b) \lor (\neg a \land c)\)
Negation normal form

Theorem

*For each propositional formula there is a logically equivalent formula in NNF.*
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Proof.

First eliminate $\rightarrow$ and $\leftrightarrow$ by the appropriate equivalences.
Negation normal form

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Base case: Claim is true for \( a, \neg a, \top, \bot \).
**Theorem**

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**Proof.**

First eliminate $\rightarrow$ and $\leftrightarrow$ by the appropriate equivalences.

Base case: Claim is true for $a$, $\neg a$, $\top$, $\bot$.

Inductive case: Assume claim is true for all formulae $\varphi$ (up to a certain number of connectives) and call its NNF $\text{nnf}(\varphi)$. 

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- $\text{nnf}(\varphi \land \psi) = (\text{nnf}(\varphi) \land \text{nnf}(\psi))$
### Theorem

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- $\text{nnf}(\varphi \land \psi) = (\text{nnf}(\varphi) \land \text{nnf}(\psi))$
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- $\text{nnf}(\neg\neg \varphi) = \text{nnf}(\varphi)$
For each propositional formula there exist logically equivalent formulae in CNF and DNF, respectively.
Conjunctive normal form

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Proof.

The claim is true for $a$, $\neg a$, $\top$, $\bot$. 
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Let us assume it is true for all formulae $\varphi$ (up to a certain number of connectives) and call its CNF $\text{cnf}(\varphi)$ (and its DNF $\text{dnf}(\varphi)$).
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- Assume $\text{cnf}(\varphi) = \bigwedge_i \chi_i$ and $\text{cnf}(\psi) = \bigwedge_j \rho_j$ with $\chi_i, \rho_j$ being clauses.
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Then $\text{cnf}(\varphi \lor \psi) = \text{cnf}((\bigwedge_i \chi_i) \lor (\bigwedge_j \rho_j)) = \bigwedge_i \bigwedge_j (\chi_i \lor \rho_j)$ (by distributivity)
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Similar for $\text{dnf}(\varphi)$.  

$\square$
Decision Problems and Resolution
How do we decide whether a formula is satisfiable, unsatisfiable, valid, or falsifiable?

**Note**: Satisfiability and falsifiability are **NP-complete**. Validity and unsatisfiability are **co-NP-complete**.
How to decide properties of formulae

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- A CNF formula is valid iff all clauses contain two complementary literals or $\top$.
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- A DNF formula is satisfiable iff one disjunct does not contain $\bot$ or two complementary literals.
- However, transformation to CNF or DNF may take exponential time (and space!).
- One can try out all truth assignments.
- One can test systematically for satisfying truth assignments (backtracking) $\leadsto$ Davis-Putnam-Logemann-Loveland.
Deciding entailment

We want to decide \( \Theta \models \varphi \).
Deciding entailment

- We want to decide $\Theta \models \varphi$.
- Use deduction theorem and reduce to validity:
  
  \[
  \Theta \models \varphi \text{ iff } \bigwedge \Theta \rightarrow \varphi \text{ is valid.}
  \]
- Now negate and test for unsatisfiability using DPLL.
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- Different approach: Try to derive $\varphi$ from $\Theta$ – find a proof of $\varphi$ from $\Theta$. 

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Why Logic?
Propositional Logic
Syntax
Semantics
Terminology
Decision Problems and Resolution
Completeness
Resolution Strategies
Horn Clauses
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- Use inference rules to derive new formulae from $\Theta$. Continue to deduce new formulae until $\varphi$ can be deduced.
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- Now negate and test for unsatisfiability using DPLL.
- Different approach: Try to derive $\varphi$ from $\Theta$ – find a proof of $\varphi$ from $\Theta$.
- Use inference rules to derive new formulae from $\Theta$. Continue to deduce new formulae until $\varphi$ can be deduced.
- One particular calculus: resolution.
Resolution: representation

- We assume that all formulae are in CNF.
  - Can be generated using the described method.
  - Often formulae are already close to CNF.
  - There is a “cheap” conversion from arbitrary formulae to CNF that preserves satisfiability – which is enough as we will see.
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- More convenient representation:
  - CNF formula is represented as a set.
  - Each clause is a set of literals.
  - \((a \lor \neg b) \land (\neg a \lor c) \leadsto \{\{a, \neg b\}, \{\neg a, c\}\}\)
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- Empty clause (symbolically \(\Box\)) and empty set of clauses (symbolically \(\emptyset\)) are different!
Resolution: the inference rule

Let \( l \) be a literal and \( \bar{l} \) its complement.

The resolution rule

\[
\frac{C_1 \cup \{l\}, C_2 \cup \{\bar{l}\}}{C_1 \cup C_2}
\]
Resolution: the inference rule

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\( C_1 \cup C_2 \) is the resolvent of the parent clauses \( C_1 \cup \{l\} \) and \( C_2 \cup \{\overline{l}\} \). \( l \) and \( \overline{l} \) are the resolution literals.
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Example: $\{a, b, \neg c\}$ resolves with $\{a, d, c\}$ to $\{a, b, d\}$.
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Note: The resolvent is not logically equivalent to the set of parent clauses!
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Example: $\{a, b, \neg c\}$ resolves with $\{a, d, c\}$ to $\{a, b, d\}$.

Note: The resolvent is not logically equivalent to the set of parent clauses!

Notation:

\[
R(\Delta) = \{ C \mid C \text{ is resolvent of two clauses in } \Delta \}
\]
Resolution: derivations

\( D \) can be derived from \( \Delta \) by resolution (symbolically \( \Delta \vdash D \)) if there is a sequence \( C_1, \ldots, C_n \) of clauses such that

1. \( C_n = D \) and \( C_i \in R(\Delta \cup \{ C_1, \ldots, C_{i-1} \}) \), for all \( i \in \{1, \ldots, n\} \).

Define \( R^*(\Delta) = \{ D \mid \Delta \vdash D \} \).

**Theorem (Soundness of resolution)**

Let \( D \) be a clause. If \( \Delta \vdash D \) then \( \Delta \models D \).
Resolution: derivations

*D* can be derived from *Δ* by resolution (symbolically *Δ ⊢ D*) if there is a sequence *C_1, . . . , C_n* of clauses such that

\[ C_n = D \text{ and } C_i \in R(\Delta \cup \{C_1, . . . , C_{i-1}\}), \text{ for all } i \in \{1, . . . , n\}. \]

Define \( R^*(\Delta) = \{D \mid \Delta \vdash D\} \).

**Theorem (Soundness of resolution)**

Let *D* be a clause. If *Δ ⊢ D* then *Δ |= D*.

**Proof idea.**

Show \( \Delta \models D \) if *D \in R(\Delta)* and use induction on proof length.
Resolution: derivations

D can be derived from Δ by resolution (symbolically Δ ⊢ D) if there is a sequence C₁, . . . , Cₙ of clauses such that

1. Cₙ = D and Cᵢ ∈ R(Δ ∪ {C₁, . . . , Cᵢ₋₁}), for all i ∈ {1, . . . , n}.

Define R*(Δ) = {D | Δ ⊢ D}.

Theorem (Soundness of resolution)

Let D be a clause. If Δ ⊢ D then Δ |= D.

Proof idea.

Show Δ |= D if D ∈ R(Δ) and use induction on proof length.
Let C₁ ∪ {l} and C₂ ∪ {l} be the parent clauses of D = C₁ ∪ C₂.
Resolution: derivations

*D* can be derived from *Δ* by resolution (symbolically *Δ ⊢ D*) if there is a sequence *C₁, . . . , Cₙ* of clauses such that

1. *Cₙ = D* and *Cᵢ ∈ R(Δ ∪ {C₁, . . . , Cᵢ−₁})*, for all *i ∈ {1, . . . , n}*. Define *R*(Δ) = \{*D* | Δ ⊢ D\}.

Theorem (Soundness of resolution)

*Let D be a clause. If* Δ ⊢ D *then* Δ |= D.

Proof idea.

Show Δ |= D if D ∈ R(Δ) and use induction on proof length.

Let *C₁ ∪ \{l\} and C₂ ∪ \{l\}* be the parent clauses of *D = C₁ ∪ C₂*. Assume *I |= Δ*, we have to show *I |= D.*
**Resolution: derivations**

* D can be derived from \( \Delta \) by resolution (symbolically \( \Delta \models D \)) if there is a sequence \( C_1, \ldots, C_n \) of clauses such that

\[
C_n = D \quad \text{and} \quad C_i \in R(\Delta \cup \{C_1, \ldots, C_{i-1}\}), \quad \text{for all } i \in \{1, \ldots, n\}.
\]

Define \( R^*(\Delta) = \{D \mid \Delta \models D\} \).

**Theorem (Soundness of resolution)**

* Let \( D \) be a clause. If \( \Delta \models D \) then \( \Delta \models D \).

**Proof idea.**

Show \( \Delta \models D \) if \( D \in R(\Delta) \) and use induction on proof length.

Let \( C_1 \cup \{l\} \) and \( C_2 \cup \{l\} \) be the parent clauses of \( D = C_1 \cup C_2 \).

Assume \( \mathcal{I} \models \Delta \), we have to show \( \mathcal{I} \models D \).

Case 1: \( \mathcal{I} \models l \) then \( \exists m \in C_2 \) s.t. \( \mathcal{I} \models m \). This implies \( \mathcal{I} \models D \).
Resolution: derivations

$D$ can be derived from $\Delta$ by resolution (symbolically $\Delta \vdash D$) if there is a sequence $C_1, \ldots, C_n$ of clauses such that

$C_n = D$ and $C_i \in R(\Delta \cup \{C_1, \ldots, C_{i-1}\})$, for all $i \in \{1, \ldots, n\}$.

Define $R^*(\Delta) = \{ D \mid \Delta \vdash D \}$.

Theorem (Soundness of resolution)

Let $D$ be a clause. If $\Delta \vdash D$ then $\Delta \models D$.

Proof idea.

Show $\Delta \models D$ if $D \in R(\Delta)$ and use induction on proof length.
Let $C_1 \cup \{l\}$ and $C_2 \cup \{\bar{l}\}$ be the parent clauses of $D = C_1 \cup C_2$.
Assume $\mathcal{I} \models \Delta$, we have to show $\mathcal{I} \models D$.
Case 1: $\mathcal{I} \models l$ then $\exists m \in C_2$ s.t. $\mathcal{I} \models m$. This implies $\mathcal{I} \models D$.
Case 2: $\mathcal{I} \models \bar{l}$ similarly, $\exists m \in C_1$ s.t. $\mathcal{I} \models m$. 
Resolution: derivations

$D$ can be derived from $\Delta$ by resolution (symbolically $\Delta \vdash D$) if there is a sequence $C_1, \ldots, C_n$ of clauses such that

1. $C_n = D$ and $C_i \in R(\Delta \cup \{C_1, \ldots, C_{i-1}\})$, for all $i \in \{1, \ldots, n\}$.

Define $R^*(\Delta) = \{D \mid \Delta \vdash D\}$.

Theorem (Soundness of resolution)

*Let $D$ be a clause. If $\Delta \vdash D$ then $\Delta \vdash D$.*

Proof idea.

Show $\Delta \vdash D$ if $D \in R(\Delta)$ and use induction on proof length.
Let $C_1 \cup \{l\}$ and $C_2 \cup \{\lnot l\}$ be the parent clauses of $D = C_1 \cup C_2$.
Assume $\mathcal{I} \vdash \Delta$, we have to show $\mathcal{I} \vdash D$.
Case 1: $\mathcal{I} \vdash l$ then $\exists m \in C_2$ s.t. $\mathcal{I} \vdash m$. This implies $\mathcal{I} \vdash D$.
Case 2: $\mathcal{I} \vdash \lnot l$ similarly, $\exists m \in C_1$ s.t. $\mathcal{I} \vdash m$.
This means that each model $\mathcal{I}$ of $\Delta$ also satisfies $D$, i.e., $\Delta \vdash D$. 

April 17, 2018 Nebel, Lindner, Engesser – KR&R
Resolution: completeness?

Do we have

$$\Delta \models \varphi \text{ implies } \Delta \vdash \varphi ?$$

Of course, could only hold for CNF.

However:

$$\{ \{a, b\}, \{\neg b, c\}\} \models \{a, b, c\} \nvdash \{a, b, c\}.$$
Resolution: completeness?

Do we have

$$\Delta \models \varphi$$ implies $$\Delta \vdash \varphi$$?

Of course, could only hold for CNF.

However:

$$\left\{ \{a,b\}, \{\neg b, c\} \right\} \models \{a,b,c\}$$

$$\not\vdash \{a,b,c\}$$
Resolution: completeness?

Do we have

\[ \Delta \models \phi \text{ implies } \Delta \vdash \phi? \]

Of course, could only hold for CNF.
However:

\[
\begin{align*}
\{ \{a, b\}, \{\neg b, c\} \} & \models \{a, b, c\} \\
\not\vdash \{a, b, c\}
\end{align*}
\]

However, one can show that resolution is refutation-complete:

\[ \Delta \text{ is unsatisfiable iff } \Delta \vdash \square. \]
Resolution: completeness?

Do we have

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Of course, could only hold for CNF.

However:

\[ \{a, b\}, \{\neg b, c\} \models \{a, b, c\} \]

\[ \not\vdash \{a, b, c\} \]

However, one can show that resolution is refutation-complete:

\[ \Delta \text{ is unsatisfiable iff } \Delta \vdash \Box. \]

Entailment: Reduce to unsatisfiability testing and decide by resolution.
Resolution strategies

- Trying out all different resolutions can be very costly, and might not be necessary.

- There are different resolution strategies.

- Examples:
  - **Input resolution** \((R_I(\cdot))\): In each resolution step, one of the parent clauses must be a clause of the input set.
  - **Unit resolution** \((R_U(\cdot))\): In each resolution step, one of the parent clauses must be a unit clause.
  - Not all strategies are (refutation) completeness preserving. Neither input nor unit resolution is. However, there are others.
Horn clauses & resolution

**Horn clauses:** Clauses with at most one positive literal

**Example:** \((a \lor \neg b \lor \neg c), (\neg b \lor \neg c)\)
Horn clauses & resolution

Horn clauses: Clauses with at most one positive literal

Example: \((a \lor \neg b \lor \neg c), (\neg b \lor \neg c)\)

Proposition

Unit resolution is refutation-complete for Horn clauses.

Proof idea.

Consider \(R^*_U(\Delta)\) of Horn clause set \(\Delta\). We have to show that if \(\square \not\in R^*_U(\Delta)\), then \(\Delta(\equiv R^*_U(\Delta))\) is satisfiable.
Horn clauses & resolution

Horn clauses: Clauses with at most one positive literal
Example: \((a \lor \neg b \lor \neg c), (\neg b \lor \neg c)\)

Proposition

Unit resolution is refutation-complete for Horn clauses.

Proof idea.

Consider \(R_u^*(\Delta)\) of Horn clause set \(\Delta\). We have to show that if \(\square \notin R_u^*(\Delta)\), then \(\Delta(\equiv R_u^*(\Delta))\) is satisfiable.

- Assign true to all unit clauses in \(R_u^*(\Delta)\).
- Those clauses that do not contain a literal \(l\) such that \(\{l\}\) is one of the unit clauses have at least one negative literal.
- Assign true to these literals.
- Results in satisfying truth assignment for \(R_u^*(\Delta)\) (and \(\Delta \subseteq R_u^*(\Delta)\)).
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