2.1 Introduction

Agent Architectures
Symbolic reasoning agents

What are agent architectures?

An agent architecture is a software design for an agent.

The last lecture introduced the top-level decomposition into:

- perception
- state
- decision
- action

An agent architecture defines:
- Key data structures
- operations on data structures
- control flow between operations
Agent Architectures

“[A] particular methodology for building [agents]. It specifies how ... the agent can be decomposed into the construction of a set of component modules and how these modules should be made to interact. The total set of modules and their interactions has to provide an answer to the question of how the sensor data and the current internal state of the agent determine the actions ... and future internal state of the agent. An architecture encompasses techniques and algorithms that support this methodology.” (Pattie Maes, 1991)

“A specific collection of software (or hardware) modules, typically designated by boxes with arrows indicating the data and control flow among the modules. A more abstract view of an architecture is as a general methodology for designing particular modular decomposition for particular tasks.” (Leslie Kaelbing, 1991)

Example: “MoveIt” component for ROS

“The high-level system architecture for the primary node provided by MoveIt! called move_group, pulling all the individual components together to provide a set of ROS actions and services for users to use.”

Types of Agents

- 1956 – present: Symbolic Reasoning Agents
  Its purest expression, proposes that agents use explicit logical reasoning in order to decide what to do

- 1985 – present: Reactive Agents
  Problems with symbolic reasoning led to a reaction against this – led to the reactive agents movement

- 1990 – present: Hybrid Agents
  Hybrid architectures attempt to combine the best of symbolic and reactive architectures

Symbolic Reasoning Agents

The classical approach to building agents:
- Agents as a particular type of knowledge-based system
- Make use of associated methodologies
- Paradigm known as symbolic AI

Definition 16: Deliberative Agent (Architecture)
A deliberative agent or agent architecture is one that:
- contains an explicitly represented, symbolic model of the world, and
- makes decisions (e.g. about actions to perform) via symbolic reasoning.
Representing the environment symbolically

The transduction problem:
- how to translate the real world into an accurate, adequate symbolic description
- in time for that description to be useful ⇒ vision, speech understanding, etc.

Problems with Symbolic Approaches

The representation/reasoning problem:
- how to symbolically represent information about complex real-world entities and processes
- how to let agents reason with this information in time for the results to be useful ⇒ knowledge representation, automated reasoning, planning

In general:
- Real-world problems (apart from games like chess) are very hard to be solved this way
- Underlying problem is the complexity of symbol manipulation algorithms in general, e.g. intractability of of search-based symbol manipulation algorithms
- These problems let to alternative approaches discussed later...

2.2 Deductive Reasoning Agents

Main assumptions:
- Agents use symbolic representations of the world around them
- They reason about the world by syntactically manipulating symbols
- Assumed sufficient to achieve intelligent behavior according to the "symbol system hypothesis"

Deductive reasoning ⇒ specific kind of symbolic approach where representations are logical formulae and syntactic manipulation is achieved by logical deduction (theorem proving)
Agents as theorem provers – background

Simple model of “deliberate” agents:
- Internal state is a database of first-order logic formulae
- Corresponds to the “beliefs” of the agent (may be erroneous, out of date, etc.)
- Let \( L \) be the set of sentences of first-order logic, \( D = \mathcal{P}(L) \) be the set of all \( L \)-databases (i.e., set of internal agent states)
- Write \( \Delta \vdash \rho \psi \) if \( \psi \) can be proved from DB \( \Delta \in D \) using (only) deduction rules \( \rho \)

Modify abstract agent architecture specification:

\[
\begin{align*}
\text{see:} & \quad E \rightarrow \text{Per} & (1) \\
\text{action:} & \quad D \rightarrow \text{Ac} & (2) \\
\text{next:} & \quad D \times \text{Per} \rightarrow D & (3)
\end{align*}
\]

Example: the vacuum world

A small robot to help with housework:
- Perceptions: dirt sensor, orientation (N, S, E, W)
- Actions: suck up dirt, step forward, turn right (90°)
- Starting point \((0,0)\), robot cannot exit the room

Goal: traverse room continually, search for dirt and remove it

Action selection as theorem proving
- Assume special predicate \( \text{Do}(\alpha) \) for action description \( \alpha \)
- If \( \text{Do}(\alpha) \) can be derived, \( \alpha \) is the best action to perform

Control loop:

\[
\begin{align*}
\text{function } \text{action}(\Delta : D) & : \text{return an action } \alpha \in \text{Ac} \\
\text{for each } & \quad \alpha \in \text{Ac} \\
\text{if } & \quad \Delta \vdash \text{Do}(\alpha) \\
\text{return } & \quad \alpha; \\
\text{end} \\
\text{for each } & \quad \alpha \in \text{Ac} \\
\text{if } & \quad \Delta \not\vdash \text{Do}(\alpha) \\
\text{return } & \quad \alpha; \\
\text{end} \\
\text{return } & \quad \text{null};
\end{align*}
\]

If no “good” action \( \Rightarrow \) search for consistent action instead.

Example: the agent for the vacuum world

A sketch of the agent:
Example: Logical formulation
Formulate this problem in logical terms:
- Percept is either dirt or null
- Actions are forward, suck, and turn
- Domain predicates are $In(x,y)$, $Dirt(x,y)$, $Facing(d)$

next() function must update internal (belief) state of agent:

\[
\text{old}(\Delta) := \{P(t_1 \ldots t_n) | P \in \{In, Dirt, Facing\} \land P(t_1 \ldots t_n) \in \Delta\}
\]

Assume new : $D \times Per \rightarrow D$ adds new predicates to database, then

\[
\text{next}(\Delta, p) = (\Delta/\text{old}(\Delta)) \cup \text{new}(\Delta, p)
\]

Agent behavior specified by hardwired rules, e.g.:

\[
\begin{align*}
In(x, y) \land Dirt(x, y) & \Rightarrow \text{Do(suck)} \\
In(0,0) \land Facing(N) \land \neg\text{Dirt(0,0)} & \Rightarrow \text{Do(forward)} \\
In(0,1) \land Facing(N) \land \neg\text{Dirt(0,1)} & \Rightarrow \text{Do(forward)} \\
In(0,2) \land Facing(N) \land \neg\text{Dirt(0,2)} & \Rightarrow \text{Do(turn)} \\
In(0,2) \land Facing(E) & \Rightarrow \text{Do(forward)}
\end{align*}
\]

Critique of the symbolic approach
How useful is this kind of agent design in practice?
- Naive implementation certainly won't work!
- What if world changes after optimal action was chosen?
  ⇒ notion of calculative rationality, i.e. decision of system optimal, when decision making began
- In case of first-order logic, not even termination can be guaranteed (undecidability) . . . let alone real time behavior
- Formalization of real-world environments often difficult and counter-intuitive
- Clear advantage: elegant semantics, declarative flavor, simplicity

Agent oriented programming
Agent oriented programming (AOP)
Based on Shoham’s (1993) idea of bringing societal view into agent programming (AGENT0 programming language). Programming agents using mentalistic notions (beliefs, desires, intentions).
Agent specified in terms of:
- set of capabilities
- set of initial beliefs
- set of initial commitments
- set of commitment rules

Key component commitment rules:
- composed of message condition, mental condition, and action (private or communicative)
- rule matching determines whether rule should fire
- message types are requests, unrequests (change commitments), and info rm messages (change beliefs)

AOP – commitment rules in AGENT0
Suppose we want to describe commitment rule:

"If I receive a message from agent requesting me to do action at time and I believe that (a) agent is a friend, (b) I can do the action and (c) at time I am not committed to doing any other action then commit to action at time"

This can be expressed in AGENT0 like so:

\[
\text{COMMIT(agent,REQUEST,DO(time,action)} \\
(B, [now,Friend agent] AND CAN(self,action) AND NOT [time,CMT(self,anyaction)]), \\
self, DO(time,action))
\]

Top-level control loop used to describe AGENT0 operation:
- Read all messages, update beliefs and commitments
- Execute all commitments with satisfied capability condition
- Loop.
Concurrent MetateM

Features of Concurrent MetateM:

- Based on direct execution of logical formulae
- Concurrently executing agents communicate via asynchronous broadcast message passing
- Two-part agent specification:
  - interface defines how agent interacts with other agents
  - computational engine defines how agent will act
- Agent interface consists of:
  - unique agent identifier
  - ‘environment propositions’, i.e. a set of symbols specifying which messages the agent accepts
  - ‘component propositions’, i.e. a set of symbols specifying messages the agent will send
- Example: interface definition of “stack”

\[
\Rightarrow \text{stack}(\text{pop}, \text{push})[\text{popped}, \text{full}]
\]

Propositional MetateM logic

Propositional logic with temporal operators:

- \( \bigcirc \varphi \): \( \varphi \) is true tomorrow
- \( \bigodot \varphi \): \( \varphi \) was true yesterday
- \( \bigtriangledown \varphi \): \( \varphi \) now or at some point in the future
- \( \square \varphi \): \( \varphi \) now and at all points in the future
- \( \lozenge \varphi \): \( \varphi \) was true sometimes in the past
- \( \blacklozenge \varphi \): \( \varphi \) was always true in the past
- \( \varphi \psi \): \( \psi \) some time in the future, \( \varphi \) until then
- \( \varphi S \psi \): \( \psi \) some time in the past, \( \varphi \) since then (but not now)
- \( \varphi W \psi \): \( \psi \) was true unless \( \varphi \) was true in the past
- \( \varphi Z \psi \): like “\( S \)” but \( \varphi \) may have never become true

Beginning of time: special nullary operator (start) satisfied only at the beginning

Concurrent MetateM – program rules

Computational engine of Concurrent MetateM is based on MetateM, which is based on program rules, which are temporal logic formulae of the form:

\[ \text{antecedent about past} \Rightarrow \text{consequent about present and future} \]

“Declarative about past ⇒ consequent about present and future” paradigm (Gabbay, 1989)

Agents try to make present and future true, given the past:

- Collect constraints with old commitments
- These taken together form current constraints
- Next state is constructed by trying to fulfil these
- Disjunctive formula ⇒ choices
- Unsatisfied commitments are carried over to the next cycle

Agent execution

Some examples:

- \( \square \text{important}(\text{agents}) \): “now and for all times agents are important”
- \( \Diamond \text{important}(\text{agents}) \): “agents will be important at some point”
- \( \lnot \text{friends}(\text{us}) \cup \text{apologize}(\text{you}) \): “not friends until you apologize”
- \( \bigcirc \text{apologize}(\text{you}) \): “you will apologize tomorrow”

Agent execution:

Attempt to match past-time antecedents of rules against history and execute consequents of rules that fire.

More precisely:

1. Update history with received messages (environment propositions)
2. Check which rules fire by comparing antecedents with history
3. Jointly execute fired rule consequents together with commitments carried over from previous cycles
4. Loop.
Specification of an example system

Consider the following definition of a system:

\[
rp(\text{ask}^1, \text{ask}^2)(\text{give}^1, \text{give}^2) :
\]

\[
\circ \text{ask}^1 \Rightarrow \lozenge \text{give}^1
\]

\[
\circ \text{ask}^2 \Rightarrow \lozenge \text{give}^2
\]

\[
\text{start} \Rightarrow \Box \neg (\text{give}^1 \land \text{give}^2)
\]

\[
rc^1(\text{give}^1)(\text{ask}^1) :
\]

\[
\text{start} \Rightarrow \text{ask}^1
\]

\[
\circ \text{ask}^1 \Rightarrow \text{ask}^1
\]

\[
rc^2(\text{ask}^1, \text{give}^2)(\text{ask}^2) :
\]

\[
\circ (\text{ask}^1 \land \neg \text{ask}^2) \Rightarrow \text{ask}^2
\]

Example run

\(rp\) is [r]esource [p]roducer, cannot [g]ive to both agents simultaneously, but will give eventually to any agent that asks.

\(rc^1\) and \(rc^2\) are resource consumers:

\(rc^1\) will ask in every cycle

\(rc^2\) only asks if it has not asked previously and \(rc^1\) has asked

Example run:

\[
\begin{array}{|c|c|c|c|}
\hline
\text{time} & \text{rp} & rc^1 & rc^2 \\
\hline
0 & \text{ask}^1 & & \\
1 & \text{ask}^1 & \text{ask}^1 & \text{ask}^2 \\
2 & \text{ask}^1, \text{ask}^2, \text{give}^1 & \text{ask}^1 & \\
3 & \text{ask}^1, \text{give}^2 & \text{ask}^1, \text{give}^1 & \text{ask}^2 \\
4 & \text{ask}^1, \text{ask}^2, \text{give}^1, \text{give}^2 & \text{ask}^1 & \text{give}^2 \\
5 & \ldots & \ldots & \ldots \\
\hline
\end{array}
\]

2.3 Summary

- Agent architectures / MoveIt (ROS)
- Symbolic Reasoning Agents
- Agents as theorem provers
- General architecture, vacuum world example
- Agent-oriented programming (AGENT0): first approach to use mentalistic concepts in programming (but not a true programming language)
- Concurrent MetateM & temporal logic: powerful and expressive but somewhat specific

⇒ Next time: Practical reasoning agents

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